

Some logical aspects of topos theory and examples from algebraic geometry (Part II)

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Recap of Part I

- sheaves $\mathcal{F} \in \text{Sh}(X) \rightsquigarrow \mathcal{F}(U)$ sets
- $\text{Sh}(X)$ is a topos
- internal language:

$$\begin{array}{c} \mathcal{F}(U) \text{ sets} \\ \downarrow \\ \mathcal{F}(V) \quad V \subseteq U \\ \hline f: \mathcal{F} \rightarrow \mathcal{G} \\ R \subseteq \mathcal{F} \times \mathcal{G} \end{array}$$

$$U \models s = s' \quad :\Leftrightarrow \quad s = s' \quad s, s' \in \mathcal{F}(U)$$

$$U \models \phi \vee \psi \quad :\Leftrightarrow \quad \text{there is } U = U_i, U_i \text{ and } U_i \models \phi \text{ or } U_i \models \psi$$

$$U \models \perp \quad :\Leftrightarrow \quad U = \emptyset$$

$$U \models \phi \Rightarrow \psi \quad :\Leftrightarrow \quad (V \models \phi) \Rightarrow (V \models \psi) \text{ f.o. } V \subseteq U$$

$$\neg \phi \equiv \underline{\phi \Rightarrow \perp}$$

$$\underbrace{(U \cup U) \models \phi}_{U \models \phi} = [\phi]$$

$$U \models \phi \Rightarrow \perp \quad \text{iff} \quad \underbrace{(V \models \phi) \Rightarrow V = \emptyset}_{[\phi] \cap U = \emptyset}$$

$$U \models (\phi \Rightarrow \perp) \Rightarrow \perp$$

$$\text{iff} \quad \underbrace{V \models (\phi \Rightarrow \perp)}_{[\phi] \cap V = \emptyset} \Rightarrow V = \emptyset$$

$$\text{iff} \quad [\phi] \text{ is } \underbrace{\text{dense}}_{\text{open}} \text{ in } U$$

Interpreting IZF in $\text{Sh}(X)$

problem: $x \in y$ only makes sense for
 $x: \mathcal{F}, y: \mathcal{P}(\mathcal{F})$

$$\bigcup_{\alpha} \mathcal{P}^{\alpha}(\emptyset) =: V$$

$$"e" \in V \times V$$

$$V_{\alpha} := \mathcal{P}^{\alpha}(\emptyset)$$

$$"e" \in V_{\alpha} \times V_{\alpha+1}$$

Classifying toposes

Definition

A *geometric formula* is a first order formula ϕ that only uses the logical connectives

$=, \top, \wedge, \perp, \vee, \exists, \forall.$

~~\neg~~ ~~\Rightarrow~~ ~~\Leftrightarrow~~

A *geometric theory* is a first-order theory with axioms

$\phi \vdash_{\vec{x}:\vec{A}} \psi$

$((\forall x_1:A_1 \dots \forall x_n:A_n.$

where ϕ and ψ are geometric formulas.

$\phi \Rightarrow \psi)$

$\top \vdash_{x,y,z:A} x \cdot (y \cdot z) = (x \cdot y) \cdot z$

$x=0 \vee \exists y:A. xy=1$

$x \neq 0 \Rightarrow \dots$

Why geometric logic?

Every continuous map $f : X \rightarrow Y$ induces

$$\mathrm{Sh}(X) \xleftarrow{f^*} \mathrm{Sh}(Y).$$

$$\begin{array}{ccccc} \phi & \wedge & \psi & & \\ S_1 & \wedge & S_1 & \hookrightarrow & S_1 \\ \downarrow & & \downarrow & & \downarrow \\ S_2 & \hookrightarrow & S & & \end{array}$$

This f^* preserves finite limits and all colimits.

$\Rightarrow f^*$ preserves models of geometric theories.

In fact, every *geometric morphism* $f : \mathcal{E}' \rightarrow \mathcal{E}$ has such a “pull-back part”

$$\mathcal{E}' \xleftarrow{f^*} \mathcal{E}.$$

Definition

A geometric theory \mathbb{T} is *classified* by a topos $\mathcal{E}_{\mathbb{T}}$ if there is a *universal model* $U_{\mathbb{T}} \in \text{Mod}_{\mathcal{E}_{\mathbb{T}}}(\mathbb{T})$, that is:

$$\begin{array}{ccc} \text{Geom}(\mathcal{E}', \mathcal{E}_{\mathbb{T}}) & \cong & \text{Mod}_{\mathcal{E}'}(\mathbb{T}) \\ f & \mapsto & f^* U_{\mathbb{T}} \end{array}$$

is an equivalence of categories for every topos \mathcal{E}' .

think: bijection

Aside:

A *point* of \mathcal{E} is a geometric morphism $\text{Set} \rightarrow \mathcal{E}$.

$$\text{points of } \mathcal{E}_{\mathbb{T}} = \text{Geom}(\text{Set}, \mathcal{E}_{\mathbb{T}}) = \text{Mod}_{\text{Set}}(\mathbb{T})$$

Example

What is the classifying topos of sub-singletons?

- ▶ one sort A
- ▶ no function/relation symbols
- ▶ one axiom: $\top \vdash_{x,y:A} x = y$

$$\left\{ \begin{array}{l} p \vee q \\ \neg(p \wedge q) \end{array} \right.$$

$$p \vee \neg p$$

$$\text{Cont}(X, Y) = \{ \text{opens } U \subseteq X \}$$

$$Y = \text{O} \circledast$$

Grothendieck topos

Theorem

Every geometric theory has a unique (up to equivalence) classifying topos.

Every topos classifies some geometric theory.

Theorem

A geometric sequent $\phi \vdash_{\vec{x}:\vec{A}} \psi$ is fulfilled for $U_{\mathbb{T}}$ iff it is provable modulo \mathbb{T} .

Examples from algebraic geometry

Definition

A *site* is a category together with a “notion of covering” (a *Grothendieck topology*).

	site for $\text{Sh}(X)$	site for Zar
objects	$U \subseteq X$ open	finitely presented ring A
morphisms	$V \subseteq U$	ring homomorphism $B \leftarrow A$
covers	$U = \bigcup_i U_i$	$(A[a_i^{-1}] \leftarrow A)_{i=1 \dots n}$ whenever $a_1 + \dots + a_n = 1$

Theorem

\mathbb{A}^1 is the universal local ring.
(So Zar classifies local rings.)

Proof sketch.

Infinitesimals? No problem!

$$\varepsilon \neq 0, \quad \underline{\varepsilon^2 = 0}$$

$$\mathbb{C}[X]/(X^2) \ni X$$

$$\text{Spec } \mathbb{C} = \cdot$$

$$\text{Spec } \mathbb{C}[X]/(X^2) = \bullet$$

Theorem

The big infinitesimal topos classifies the theory of infinitesimally thickened local rings.

~~Proof sketch:~~

• theory of rings: A, \dots

• A local

• $\mathfrak{m} \subseteq A$ ideal

with

$x \in \mathfrak{m}$

$$\vdash_{x \in A} \bigvee_{n \in \mathbb{N}} x^n = 0$$

