

Co-analytic families of functions

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Joint work with Julia Millhouse

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Combinatorial and descriptive set theory

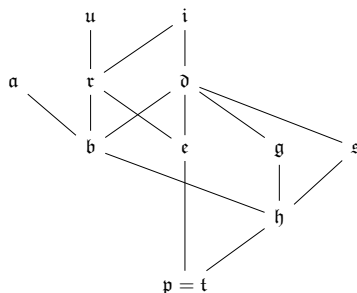
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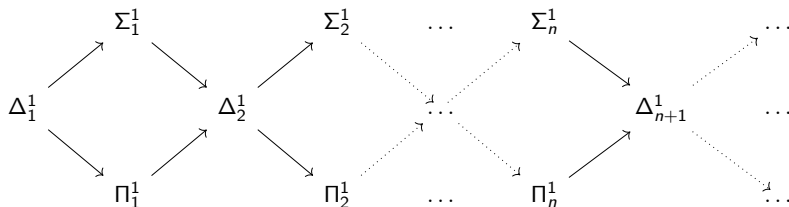
Combinatorial and descriptive set theory

- 1 In combinatorial set theory the possible sizes of certain special families of real numbers are studied.
- 2 Usually special means maximal with respect to some combinatorial or topological property.
- 3 The minimal sizes of such special families are called cardinal characteristics and their relations give rise to a very rich and complicated theory:



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- ① What is the minimal complexity of various special families of reals?
- ② In models separating cardinal characteristics $\mathfrak{x} < \mathfrak{y}$, can we additionally have witnesses for \mathfrak{x} and \mathfrak{y} of minimal complexity?
- ③ Given a special family of some complexity, can we construct a special family of lower complexity from it?

Example: Definability of mad families

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There are no analytic (i.e. Σ_1^1) mad families.

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Theorem (Miller, 1989, [10])

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This result was improved in many papers to obtain various co-analytic forcing indestructible mad families. In particular we have

Theorem (Bergfalk, Fischer, Switzer, 2022, [2])

In L there is a co-analytic tight mad family.

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For example this implies the following for the Miller model over L :

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Consistently, $\aleph_1 = \mathfrak{a} < \mathfrak{d} = \aleph_2$ and there is a Π_1^1 witness for $\mathfrak{a} = \aleph_1$.

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However, Törnquist discovered the following shortcut:

Theorem (Törnquist, 2013 [13])

If there is a Σ_2^1 mad family, then there is Π_1^1 mad family of the same size.

In particular, a Σ_2^1 witness for \mathfrak{a} immediately implies a Π_1^1 witness for \mathfrak{a} .

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In particular, a Σ_2^1 witness for \mathfrak{a} immediately implies a Π_1^1 witness for \mathfrak{a} . This gives an easier proof for the existence of a co-analytic mad family in L , but also for forcing extensions over L ...

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Let \mathbb{P} be a forcing, for which we know how to construct a \mathbb{P} -indestructible mad family (Sacks, Miller, Laver, random, etc.).

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Hence, in \mathbb{P} -extensions over L there is a co-analytic witness for $\mathfrak{a} = \aleph_1$.

Shortcuts for other families

Analogous shortcuts of the form $\Sigma_2^1 \implies \Pi_1^1$ are now known for others types of families, e.g.

- ① maximal independent families - i
(Brendle, Fischer, Khomskii, 2018, [7]),
- ② maximal eventually different families of functions - α_e
(Fischer, Schritterser, 2021, [4]),
- ③ towers - \mathfrak{t}
(Fischer, Schilhan, 2021, [3]),
- ④ Hausdorff gaps
(Millhouse, 2024, [11])

Our goal was to prove similar shortcuts for various relatives of α_e and the ideal-independence number \mathfrak{s}_{mm} .

How does the shortcut work?

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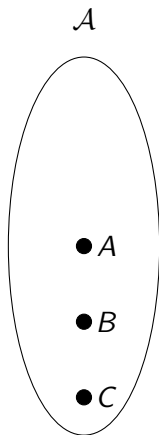
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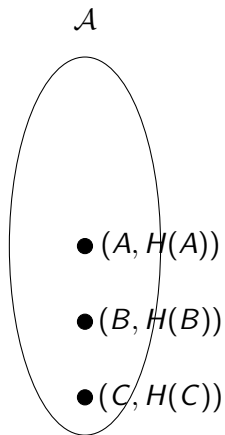
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- 5 Basically, for every $A \in [\omega]^\omega$ and $c \in {}^\omega 2$ we need to be able to definably reconstruct A and c from $\Phi(A, c)$.
- 6 Finally, the Spector-Gandy-Theorem shows that \mathcal{B} is in fact Π_1^1 :

$$B \in \mathcal{B} \iff \exists (A, c) \in \Delta_1^1(B) ((A, c) \in H \text{ and } \Phi(A, c) = B)$$

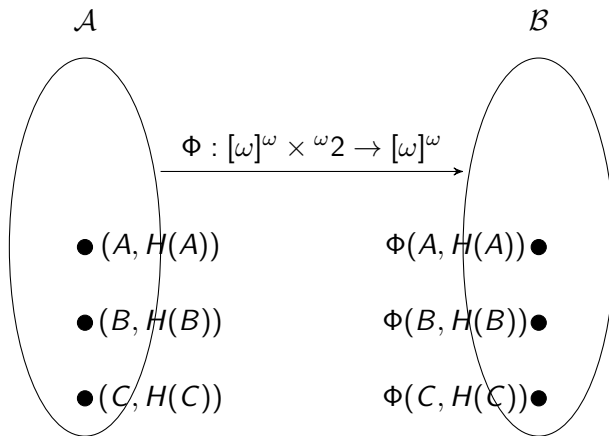
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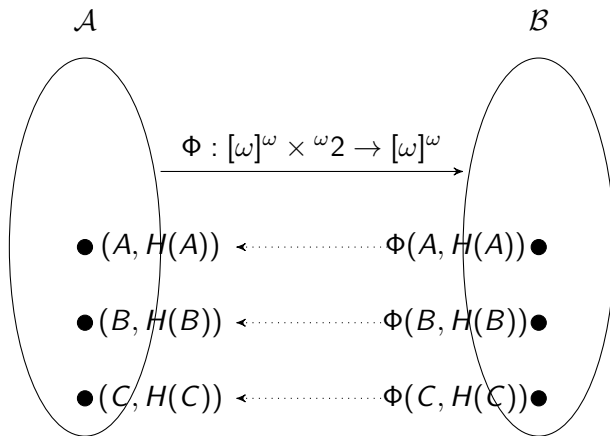
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Definition

$$\mathfrak{a}_e := \min \{ |F| \mid F \text{ is a maximal eventually different family} \}$$

Further, a maximal eventually different family of size \mathfrak{a}_e is said to be a witness for that cardinal characteristic.

Eventually different families of permutations

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Again, we define the cardinal invariant:

$$\mathfrak{a}_p := \min\{|F| \mid F \text{ is a maximal eventually different family of permutations}\}$$

Definability of eventually different families of permutations

Contrary to mad families there always is a maximal eventually different family of permutations of low complexity.

Theorem (Horowitz, Shelah, 2016, [6])

There is a Borel maximal eventually different family of permutations.

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The goal of this talk is to prove the following theorem:

Theorem (S., Millhouse, 2025)

If there is a Σ_2^1 maximal eventually different family of permutations, then there is a Π_1^1 maximal eventually different family of permutations of the same size.

Proof of the Theorem

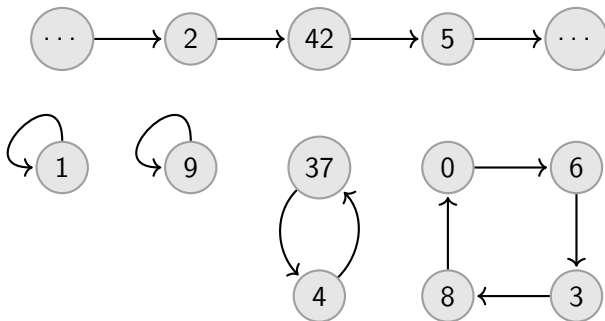
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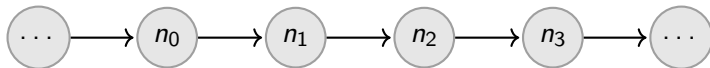
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- 3 We may visualize a permutation f as the union of chains and cycles:



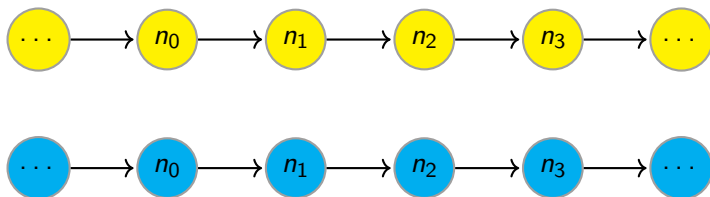
Visualization of the function Φ

For simplicity we just consider one chain



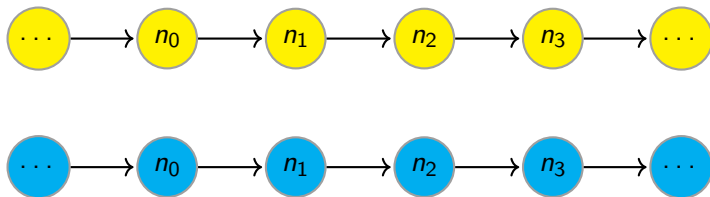
Visualization of the function ϕ

First we duplicate both the domain and range of the function, but keep the structure of f on both copies. (Thus, the new domain and range is $\omega \times 2$)



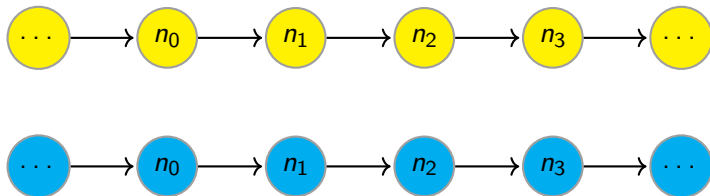
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Next, we need to adapt this permutation, so that it codes a real $c \in 2^\omega$, but still is a permutation and keeps the structure of f :



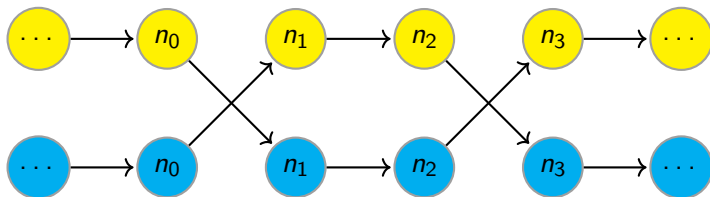
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Key Idea: At place n flip both functions values iff $c(n) = 1$:



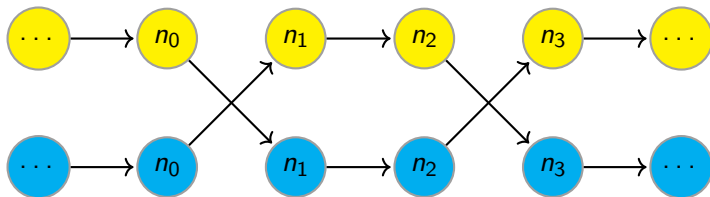
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For example, if $c(n_0) = 1$, $c(n_1) = 0$ and $c(n_2) = 1$ we get:



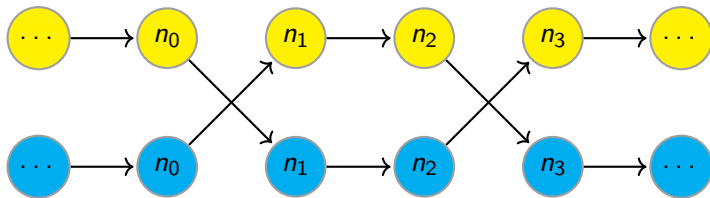
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This gives us $\Phi(f, c)$ from which we can clearly decode f and c again. It is also easy to see that $\Phi[H]$ will still be eventually different.



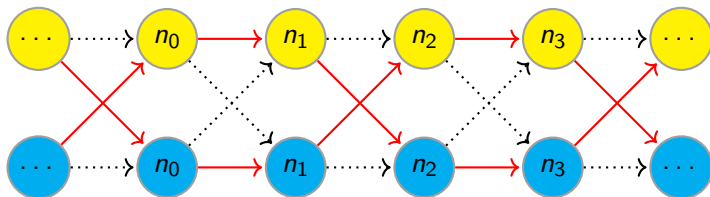
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The more difficult part of the proof is the maximality of $\Phi[H]$. In fact, $\Phi[H]$ will NOT be maximal!



Visualization of the function Φ

Note that $\Phi(f, 1 - c)$ is everywhere different from $\Phi(f, c)$:



Definition of G

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We will define a certain permutation \tilde{h} on ω from h .

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By maximality of F choose $f \in F$ with $\tilde{h} =^\infty f$.

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$$\begin{array}{ccccc} \mathbf{G} \text{ (on } \omega \times 2) & & h & & \Phi(f, H(f)) \text{ or } \Phi^*(f, H(f)) \\ & & \downarrow & & \uparrow \\ & & \tilde{h} & \longrightarrow & f \\ \mathbf{F} \text{ (on } \omega) & & & & \end{array}$$

Finally, we will either have $h =^\infty \Phi(f, H(f))$ or $h =^\infty \Phi^*(f, H(f))$.

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- 5 To this end, note that the function $g : \omega \times 2 \rightarrow \omega$ defined by $g(n, i) := p_0(h(n, i))$ is 2-to-1, i.e. every $n \in \omega$ has exactly two preimages.

Crucial lemma

Lemma

Assume that $g : \omega \times 2 \rightarrow \omega$ is 2-to-1. Then there is a function $i : \omega \rightarrow 2$, such that $\tilde{h}(n) := g(n, i(n))$ is a permutation.

Proof.

We define a bipartite graph with possible multi-edges:

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We define a bipartite graph with possible multi-edges:

- 1 We have countably many left $\{L_n \mid n \in \omega\}$ and right $\{R_n \mid n \in \omega\}$ nodes.

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- 1 We have countably many left $\{L_n \mid n \in \omega\}$ and right $\{R_n \mid n \in \omega\}$ nodes.
- 2 For each $n \in \omega$ and $i \in 2$ we have an edge $e_{n,i}$ between L_n and $R_{g(n,i)}$.

Lemma

Assume that $g : \omega \times 2 \rightarrow \omega$ is 2-to-1. Then there is a function $i : \omega \rightarrow 2$, such that $\tilde{h}(n) := g(n, i(n))$ is a permutation.

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By construction, every L_n has degree 2.

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By construction, every L_n has degree 2. Since g is 2-to-1, also every R_n has degree 2, i.e. the graph is 2-regular.

Crucial lemma

Reminder: For $n \in \omega$ and $i \in 2$ we have an edge $e_{n,i}$ between L_n and $R_{g(n,i)}$.

Proof.

Choose a perfect matching P of our graph and define

$$i(n) := i \quad \text{iff} \quad e_{n,i} \in P.$$

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So assume $\tilde{h}(n) = g(n, i(n)) = g(m, i(m)) = \tilde{h}(m)$.

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Then, by definition both $e_{n,i(n)}$ and $e_{m,i(m)}$ are in P . But they are both incident to $R_{g(n,i(n))} = R_{g(m,i(m))}$ and thus have to be equal.

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Now, let $m \in \omega$. Since, P is perfect there is an edge $e_{n,i}$ incident to R_m in P . But $e_{n,i}$ is incident $R_{g(n,i)}$, so we have $g(n, i(n)) = m$.

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Now, let $m \in \omega$. Since, P is perfect there is an edge $e_{n,i}$ incident to R_m in P . But $e_{n,i}$ is incident $R_{g(n,i)}$, so we have $g(n, i(n)) = m$. Hence, $\tilde{h}(n) = m$ and \tilde{h} is surjective. □

- 1 Reminder: We started with a permutation $h : \omega \times 2 \rightarrow \omega \times 2$.

Continuation of maximality

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- ⑤ Let $n \in A$, then there is a $j \in 2$ with

$$h(n, i(n)) = (p_0(h(n, i(n))), j) = (\tilde{h}(n), j) = (f(n), j).$$

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- ⑥ Hence, either

$$h(n, i(n)) = \Phi(f, c)(n, i(n)) \quad \text{or} \quad h(n, i(n)) = \Phi^*(f, c)(n, i(n)).$$

Reduction theorem for cofinitary groups

This finishes the proof of our theorem:

Theorem (S., Millhouse, 2025)

If there is a Σ_2^1 maximal eventually different family of permutations, then there is a Π_1^1 maximal eventually different family of permutations of the same size.

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We also came up with other ideas to obtain similar results for

- 1 Van Douwen families,
- 2 free generating sets of (strongly) maximal cofinitary groups,
- 3 maximal ideal independent families.

Definition

An eventually different family of functions $F \subseteq {}^\omega\omega$ is called **Van Douwen** iff for every infinite partial function $h : \omega \xrightarrow{\text{partial}} \omega$ there is an $f \in F$ with $h =^\infty f$.

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If there is a Σ_2^1 Van Douwen family, then there is a Π_1^1 Van Douwen family of the same size.

Cofinitary groups

Let $G \subseteq S_\omega$ be an eventually different family of permutations. Then, G need not be a group. It may fail to be closed under inverses or concatenation.

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If G is additionally a group (with concatenation) we say that G is a **cofinitary group**. G is called **maximal** if G is maximal with respect to inclusion.

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$$a_g := \min \{|G| \mid G \text{ is a maximal cofinitary group}\}$$

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Definition

$$a_g := \min \{|G| \mid G \text{ is a maximal cofinitary group}\}$$

Observation

A group $G \subseteq S_\omega$ is cofinitary iff every element is either the identity or only has finitely many fixpoints.

Definition

Let $G \subseteq S_\omega$ be a group and $G_0 \subseteq G$. Then G is a **generating set** if $G = \langle G_0 \rangle$, where $\langle G_0 \rangle$, is the group generated by G_0 .

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For cofinitary groups its often easier to provide a definable generating set instead of the whole group. If G_0 is definable, also G is definable from G_0 , however its complexity may be higher.

Definability of cofinitary groups

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For cofinitary groups its often easier to provide a definable generating set instead of the whole group. If G_0 is definable, also G is definable from G_0 , however its complexity may be higher.

Observation

Recursively, it is easy to construct a Σ_2^1 generating set for a maximal cofinitary group in L using the Δ_2^1 well-order.

Theorem (Gao, Zhang, 2008 [5])

In \mathcal{L} there is a Π_1^1 generating set for a maximal cofinitary group.

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Theorem (Horowitz, Shelah, 2016 [6])

In ZF there is a Borel maximal cofinitary group (and a Borel maximal eventually different family).

Reduction theorem for cofinitary groups

Theorem (S., Millhouse, 2025)

If there is a Σ_2^1 generating set for a free maximal cofinitary group, which is also maximal as an eventually different family of permutations, then there also is a Π_1^1 generating set for a maximal cofinitary group of the same size, which is also maximal as an eventually different family of permutations.

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Corollary

If there is a Σ_2^1 generating set for a free tight cofinitary group, then there is a Π_1^1 generating set for a maximal cofinitary group of the same size.

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Question

If there is a Σ_2^1 maximal cofinitary group, is there is a Π_1^1 maximal cofinitary group of the same size?

Proof sketch

- 1 Fix a Σ_2^1 generating set F for a free maximal cofinitary group $G = \langle F \rangle$, which is also maximal as an eventually different family of permutations and a Π_1^1 partial function $H \subseteq {}^\omega\omega \times {}^\omega 2$ as before.

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- 2 Again, we define $\hat{F} := \Phi[H] \cup \Phi^*[H]$ and let $\hat{G} = \langle \hat{F} \rangle$.

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$$\begin{array}{ccccc}
 \hat{G} \text{ (on } \omega \times 2) & & h & & w[\Phi(f, \vec{H}(f))] \text{ or } w[\Phi^*(f, \vec{H}(f))] \\
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 \mathbf{G} \text{ (on } \omega) & &
 \end{array}$$

Crucially, we can find a natural surjective group homomorphism $\Psi : \hat{G} \rightarrow G$ and compute its kernel as $\{\text{id}, \tau\}$, where $\tau(n, i) := (n, 1 - i)$ is the flip map. This also implies that $\hat{G} \cong G \times \mathbb{Z}/2$.

Definition

A family $\mathcal{I} \subseteq [\omega]^\omega$ is called ideal independent iff for every $A \in \mathcal{I}$ and $\mathcal{I}_0 \in [\mathcal{I}]^{<\omega}$ we have that $A \not\subseteq^* \bigcup \mathcal{I}_0$.

Ideal independent families

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Definition

Again, we define a cardinal invariant:

$$\mathfrak{s}_{\text{mm}} := \min \{ |\mathcal{I}| \mid \mathcal{I} \text{ is a maximal ideal independent family} \}$$

Observation

Mad families and maximal independent families are ideal independent, but neither have to be maximal as an ideal independent family.

Reduction theorem for ideal independent families

Theorem (Bardyla, Cancino, Fischer, Switzer, 2025, [1])

(CH) There is a maximal ideal independent family indestructible by any proper, ω_ω -bounding, p -point preserving forcing.

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Corollary

Consistently, $\aleph_1 = \mathfrak{s}_{mm} < \mathfrak{c}$ with a Π_1^1 witness for \mathfrak{s}_{mm} .

- 1 Fix a Σ_2^1 maximal ideal independent family \mathcal{I} and a Π_1^1 partial function $H \subseteq [\omega]^\omega \times [\omega]^\omega$ as before.

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- 2 Define a function $\Psi : [\omega]^\omega \times [\omega]^\omega \rightarrow [\omega]^\omega$ and $z \in [\omega]^\omega$ by

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- 1 Define $\mathcal{J} := \Psi[H] \cup \{z\}$.
- 2 It is not too difficult to show that \mathcal{J} is maximal ideal independent and Π_1^1 .

Summary

Invariant	Borel	$\Pi_1^1 + \neg\mathbf{CH}$	$\Sigma_2^1 \Rightarrow \Pi_1^1$
a	No	Yes	Yes
a_v	No	Yes	?
a_e	Yes	Yes	Yes
a_p	Yes	Yes	?
a_g	Yes	Yes	?
s_{mm}	?	?	?

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a_p	Yes	Yes	Yes
a_g	Yes	Yes	Yes*
s_{mm}	?	Yes	Yes

*under extra assumptions and only for the generators

References I



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Thank you for your attention!