

The Constructible Universe

VI

13 May 2024

§M Definition of the constructible universe

REGAP

Definition If X is a set,
then let

$$\mathcal{D}(X) := \{A \subseteq X; A \text{ is definable with parameters over } X\}$$

$$= \{A \subseteq X; \exists \Delta(A, X)\}$$

Sometimes called "the definable power set operation".

universe.

Recursion

$$V_0 := \emptyset$$

$$V_{\alpha+1} := \mathcal{P}(V_\alpha)$$

$$V_\lambda := \bigcup_{\alpha < \lambda} V_\alpha$$

von Neumann hierarchy

[Family that satisfies N produces N since $\exists \Delta$ Foundation.]

"definable power set" is a slight misnomer

The constructible hierarchy

$$L_0 := \emptyset$$

$$L_{\alpha+1} := \mathcal{D}(L_\alpha)$$

$$L_\lambda := \bigcup_{\alpha < \lambda} L_\alpha$$

$$\mathcal{D}(X) \subseteq \mathcal{P}(X)$$

History of letter L unknown.

Note that the \mathcal{D} -operator was shown to be absolute in §10.

Main Theorem of §10 said that transfinite recursion preserves absoluteness.

From lecture V

Corollary

The statement

$$x \in L_\alpha$$

is absolute for transitive models of set theory.

Definition x is called constructible if there is $\alpha \in \text{Ord}$ s.t. $x \in L_\alpha$.CorollaryIf $M \subseteq N$ and $\text{Ord} \cap M = \text{Ord} \cap N$, then "x is constructible" is absolute between M & N .Thus if I write $L := \bigcup_{\alpha \in \text{Ord}} L_\alpha$, then L (the constructible sets) is contained in every model of ZF.Thusif we can show that $L \models \text{ZF}$, then L is the **MINIMAL MODEL** of ZF.Main goal for today: $L \models \text{ZF}$

Then: it is the minimal transitive model.

Consequences of our definition

Lecture V

Observe ① For all X , $\mathcal{D}(X) \subseteq \mathcal{P}(X)$.
② If X is transitive, then so is $\mathcal{D}(X)$.
[if $a \in b \in \mathcal{D}(X)$. But every element of X is definable over X with parameters.
 $\Rightarrow b \subseteq X$
 $\Rightarrow a \in X$
 $\Rightarrow a \in \mathcal{D}(X)$.]

\Rightarrow By induction, we can prove

$$L_\alpha \subseteq V_\alpha$$

\Rightarrow By induction, each L_α is transitive.
Therefore, the class of constructible sets is a transitive class.

From the von Neumann hierarchy

$$V_\alpha \cap \text{Ord} = \alpha$$

$$L_\alpha \cap \text{Ord} \subseteq \alpha$$

I claim $L_\alpha \cap \text{Ord} = \alpha$.

Proof (by induction) True for $\alpha = 0$.

The limit case is trivial. Suppose $L_\alpha \cap \text{Ord} = \alpha$ and show that $\alpha \in L_{\alpha+1}$.

[Then $\alpha \in L_{\alpha+1} \cap \text{Ord} \subseteq V_{\alpha+1} \cap \text{Ord} = \alpha+1$ which proves the claim.]

Note that α is the set of ordinals in L_α
by IH, so

$$\{x \in L_\alpha; L_\alpha \models x \text{ is an ordinal}\} = \alpha,$$

$$\text{so } \alpha \in \mathcal{D}(L_\alpha) = L_{\alpha+1}.$$

q.e.d.

Corollary The constructible sets form a proper class,
containing all ordinals.

Also For all α , $L_\alpha \in L_{\alpha+1}$.

Because

$$L_\alpha = \{x \in L_\alpha; L_\alpha \models x = x\},$$

$$\text{so } L_\alpha \in \mathcal{D}(L_\alpha) = L_{\alpha+1}.$$

Consider the relationship between the L - and
the V -hierarchy:

$$V_0 = L_0 = \emptyset.$$

$$V_n = L_n \text{ for } n \in \mathbb{N}$$

[since finite sets can always be
defined with parameters by
listing their elements]

$$L_\omega = \bigcup_{\text{new}} L_n = \bigcup_{\text{new}} V_n = V_\omega.$$

What happens at $\alpha = \omega + 1$.

$V_{\omega+1} = \mathcal{P}(V_\omega)$ is uncountable; cardinality 2^{\aleph_0} .

$L_{\omega+1} = \mathcal{D}(L_\omega)$ is countable

There is a surjection from $\text{Func} \times L_\omega^{<\omega}$ onto $\mathcal{D}(L_\omega)$:

$$(\varphi, \vec{p}) \mapsto \{x \in L_\omega; L_\omega \models \varphi(x, \vec{p})\}$$

Thus: $V_{\omega+1} \neq L_{\omega+1}$.

Now that we know that $|L_{\omega+1}| = \aleph_0$, we get surjection from $\text{Func} \times L_{\omega+1}^{<\omega}$ onto $\mathcal{D}(L_{\omega+1}) = L_{\omega+2}$.

Thus for all countable $\alpha < \omega_1$, $|L_\alpha| = \aleph_0$ by induction.

Since $\omega_1 \subseteq L_{\omega_1}$, this is not true anymore for L_{ω_1} .
Remark: $L_{\omega_1} = \bigcup_{\alpha < \omega_1} L_\alpha$, so $|L_{\omega_1}| = \aleph_1$.

Consequence The hierarchies are different in the sense that there are α s.t. $L_\alpha \neq V_\alpha$, but that not mean that there are non-constructible sets.

Def. If x is constructible, let

$\rho_L(x) := \min\{\alpha; x \in L_{\alpha+1}\}$
be the constructible rank (the L -rank)
of x .

If all sets were constructible, lots of sets in $V_{\omega+1}$
must have higher L -rank than Minimumoff rank.

Remark The fact that we understand the size of
 L_α much better than the size of V_α ;
in particular, we know at least one L -rank,
viz., L_{ω_1} , with cardinality \aleph_1
will be relevant later when we get
to CH.

§ 12 Simple axioms of ZF in the constructible universe

ZF

- Ext
- Found
- Pair
- Un
- Pow
- Sep
- Repl

true in all transitive models, so we get them in the constructible universe

true since $\forall \alpha \in L_{\omega+1}$ and the formula defining it is absolute

WORKING IN SOME $M \models ZF$, BUILDING $L = \bigcup_{\alpha \text{ ord}} L_{\alpha} \subseteq M$ TRANSITIVE SUBCLASS.

AC [left out & discussed in Lecture VII.]

(a) Pairing

Suppose $x, y \in L$, $x \in L_{\alpha}$, $y \in L_{\beta}$
 wlog $\alpha \geq \beta$, so $x, y \in L_{\alpha}$.

Clearly, since $M \supseteq L$ is a model of ZF, $\{x, y\}$ exists and $\{x, y\} \subseteq L_{\alpha}$.

The formula $\varphi(z) := (z=x \vee z=y)$ defines $\{x, y\}$ over L_{α} , so

$$\{x, y\} \in \mathcal{D}(L_{\alpha}) = L_{\alpha+1}.$$

(b) Union.

Fix $x \in L_\alpha$. We have $\cup x$ is a set in M
by ZF in M .

$$z \in \cup x \iff \exists y \underbrace{(y \in x \wedge z \in y)}_{\varphi(z, x)}$$

* implies (since L_α is transitive) $\cup x \subseteq L_\alpha$.

φ defines $\cup x$ over L_α , by absoluteness
for transitive models, absolutely,

so $\cup x \in \mathcal{D}(L_\alpha) = L_{\alpha+1}$.

(c) Power set

Fix $x \in L_\alpha$. By ZF in M , there is $p(x) \in M$.
The canonical candidate for the L -power set
of x would be

$$P := p(x) \cap L$$

If $P \in L$, then absoluteness of the formula
" $y \subseteq x$ " gives us that P is the L -power set
of x .

If P is very large compared to α , then $L_{\alpha+1}$
may not be large enough to contain all
of P .

For each $p \in P$, consider $f_L(p)$ and

$$\{f_L(p); p \in \underbrace{p(x) \cap L}_{\in M}\} =: R$$

This is a set in M by Replacement in M .

So, find γ s.t. $R \subseteq \gamma$

i.e. $p(x) \cap L \subseteq L_\gamma$.

This is defined by the absolute formula $y \subseteq x$.

$$\{y \in L_\gamma; L_\gamma \models y \subseteq x\} \in \mathcal{D}(L_\gamma) = L_{\gamma+1}$$

Remark

We have NO IDEA (at the moment) what γ is. It could be $\alpha+1$ or α^+ or α^{++} or ... It depends on the relationship between L -rank & Mirimanoff rank (which we don't understand properly yet).

Separation

This simple technique won't work
& here's the reason:

Suppose $x \in L_\alpha$, φ formula, $\vec{p} \in L_\alpha^{<\omega}$

Want $\{y \in x; L \models \varphi(y, \vec{p})\} = S$
to be in L .

What we get is unfortunately only

$\{y \in x; L_\alpha \models \varphi(y, \vec{p})\} = S'$

If φ is not absolute, then S' might be
very different from S and so the
fact that $S' \in \mathcal{D}(L_\alpha)$ doesn't
mean much.

hope There is no reason to assume that φ
is absolute between L_α and L , but
maybe we can find $\gamma > \alpha$ s.t.
 φ is absolute between L_γ and L .

→ Lévy Reflection Theorem

§13 The Lévy Reflection Theorem

Def. An assignment $\alpha \mapsto Z_\alpha$ is called a **HIERARCHY** if

- ① Z_α is a transitive set
- ② $\text{Ord} \cap Z_\alpha = \alpha$
- ③ $\alpha < \beta \implies Z_\alpha \subseteq Z_\beta$
- ④ λ limit $\implies Z_\lambda = \bigcup_{\alpha < \lambda} Z_\alpha$

If we have a hierarchy, can define a proper class $Z = \bigcup_{\alpha \in \text{Ord}} Z_\alpha$ [by ② $\text{Ord} \subseteq Z$] and

$$j_Z(x) := \min\{\alpha; x \in Z_\alpha\}$$

for any $x \in Z$.

Lévy Reflection Theorem

If $M \models ZF$, $\alpha \mapsto Z_\alpha$ is a hierarchy, and φ is a formula, there are unboundedly many α s.t. φ is absolute between Z_α and Z .

Remark We make no assumption that $Z \models ZF$ or even a fragment thereof. In fact $Z_\alpha := \alpha$ is a hierarchy in our sense.

The proof of the LRT is essentially an elaboration of

- (a) Tarski-Vaught Lemma
- (b) the use of TVL in constructing elementary submodels.

Lemma (Improved Tarski-Vaught Lemma)

Let $M \subseteq N$ be any models and Φ be a collection of formulas closed under subformulas. Then TFAE

(a) all formulas in Φ are absolute between M & N

(b) for any $\varphi \in \Phi$ s.t. $\varphi = \exists x \psi$, we have

$$\forall \vec{y} \in M \left(\exists a \in N \ N \models \psi(a, \vec{y}) \implies \exists a \in M \ N \models \psi(a, \vec{y}) \right).$$

Tarski-Vaught condition

Proof.

Note that the original TVL is just the special case of $\Phi = \text{True}$.

The proof of TVL proves this improved version; the closure of Φ under subformulas is needed to apply the IH to ψ which is a subformula of φ .

q.e.d.

Proof of LRT

Fix φ and let Φ be the closure of $\{\varphi\}$ under subformulas. Note that Φ is a finite set of formulas.

By the improved TVL, we only need to find \mathcal{Q} s.t. $Z_{\mathcal{Q}}$ contains witnesses ($n \geq$) to all existential formulas in Φ .

Fix some α — need to find $\mathcal{Q} > \alpha$ s.t. φ absolute between $Z_{\mathcal{Q}}$ and Z .

If $\exists x \psi \in \Phi$, check whether

$$Z \models \exists x \psi(x, \vec{p}) \quad \text{for some } \vec{p} \in Z_{\alpha}^{\omega}$$

If not, set $F(\alpha, \psi, \vec{p}) := 0$.

O/w, set $F(\alpha, \psi, \vec{p}) := \min \{ f_Z(w); Z \models \psi(w, \vec{p}) \}$

If $\psi \in \Phi$, define

$$F_\psi(\alpha) := \sup \{ F(\alpha, \psi, \vec{p}) ; \vec{p} \in Z_\alpha^{<\omega} \}$$

$\{ F(\alpha, \psi, \vec{p}) ; \vec{p} \in Z_\alpha^{<\omega} \}$ is a set because of Replacement in M

So, the supremum exists.

Define $F(\alpha) := \max \{ F_\psi(\alpha) ; \psi \in \Phi \} \cup \{ \alpha \}$

$$\mathcal{D}_0 := \alpha$$

$$\mathcal{D}_{i+1} := F(\mathcal{D}_i)$$

$$\mathcal{D} := \sup \{ \mathcal{D}_i ; i \in \omega \}$$

to make it strictly increasing

Since $\mathcal{D}_0 < \mathcal{D}_1 < \mathcal{D}_2 < \dots$, \mathcal{D} is a limit ordinal.

Claim φ is absolute between $Z_\mathcal{D}$ and Z :

[Show that all formulas in Φ are abs. between $Z_\mathcal{D}$ & Z]

By improved TVL, only need to show the TV condition.

Suppose $Z \models \exists x \psi(x, \vec{p})$ where $\vec{p} \in Z_\mathcal{D}^{<\omega}$
Every $p_i \in Z_\alpha$ with $\alpha < \mathcal{D}$
So there is $\bar{\alpha} < \mathcal{D}$ s.t.
 $\vec{p} \in Z_{\bar{\alpha}}^{<\omega}$.

There is some i s.t. $\mathcal{D}_i \leq \bar{\alpha} < \mathcal{D}_{i+1}$

Thus $\vec{p} \in Z_{\mathcal{D}_{i+1}}^{<\omega}$

Therefore by definition of \mathcal{V}_{i+2} , there is
 $a \in \mathcal{Z}_{\mathcal{V}_{i+2}} \subseteq \mathcal{Z}_{\mathcal{Q}}$ s.t.

$$\mathcal{Z} \models \psi(a, \vec{p}).$$

That's what TVL needs us to show.

q.e.d.

Summary

For every hierarchy and every
single formula φ , I can find
a sufficiently large \mathcal{Q} s.t.

$$\{y \in x; \mathcal{Z} \models \varphi(y, \vec{p})\}$$

$$= \{y \in x; \mathcal{Z}_{\mathcal{Q}} \models \varphi(y, \vec{p})\}.$$

MONDAY

27 May 2024

due to Pentecost

In lecture VII, we'll use this to fix the
gap in our attempt to prove Separation
in L and also prove Replacement.

[Also discuss $L \models AC$ which is a
relative consistency proof of ZFC
from ZF.]