

The Constructible Universe

Lecture I

8 April 2024

Lectures I-III (8-22 April)
Lectures IV-X (29 April - 24 June)
Lectures XI & XII (1 & 8 July)

IN PERSON

ONLINE

IN PERSON

Exams:

Oral Exam

probably late July 2024

GOAL

Prove that AC and CH are consistent (relative to ZF).

Axiom of Choice

Continuum Hypothesis

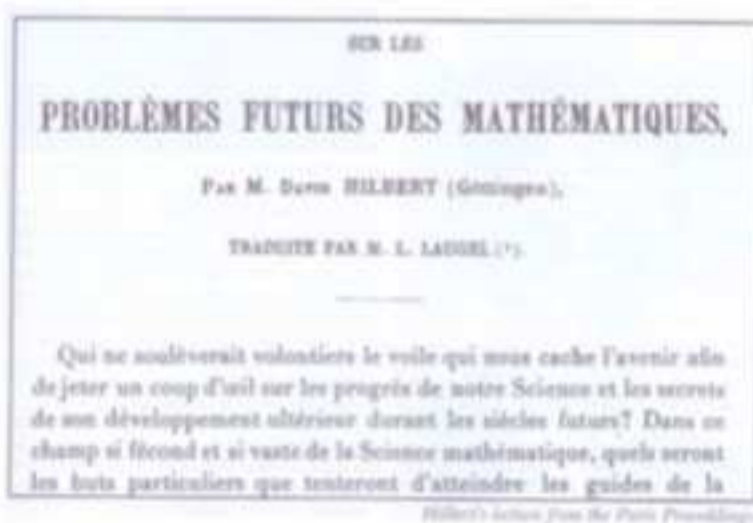
§ 1 Historical background & motivation



David Hilbert (1862-1943)

Hilbert's Problems

23 problems



Problem #1

The Continuum Problem

Q:

$$2^{\aleph_0} = \aleph_1 ?$$

Note: this unique ordinal only exists due to AC.

Cardinal: initial ordinal

α s.t. $\forall \beta < \alpha \beta \neq \alpha$

$[X \sim Y : \Leftrightarrow$

there is a bij. betw. X & $Y]$

2^{\aleph_k} : unique initial ordinal w/ bij. w/ \aleph_k .

Historical note

Hilbert noted in 1900 that the fam. $2^{\aleph_0} = \aleph_1$ implies some fragment of choice.

$$2^{\aleph_0} = \aleph_1$$

[This presupposes that \aleph_1 can be wellordered.]

Cantor's theorem:

does not need AC

No surjection from \mathbb{N} onto $\mathcal{P}(\mathbb{N})$.
 Thus: no injection from $\mathcal{P}(\mathbb{N})$ into \mathbb{N} .

With AC, this implies $\aleph_1 \leq 2^{\aleph_0}$.
 [An injection from \aleph_1 into $\mathcal{P}(\mathbb{N})$ requires AC.]

Without AC, we get:

Prop. There is a surj. from $\mathcal{P}(\mathbb{N})$ onto \aleph_1 .

Proof. By Cantor's zigzag: $\mathbb{N} \sim \mathbb{N} \times \mathbb{N}$.
 Thus: $\mathcal{P}(\mathbb{N}) \sim \mathcal{P}(\mathbb{N} \times \mathbb{N})$.

I am going to show: $f: \mathcal{P}(\mathbb{N} \times \mathbb{N}) \rightarrow \{\alpha; \omega \leq \alpha < \omega_1\}$ surjection.

Take $R \subseteq \mathbb{N} \times \mathbb{N}$. Consider (\mathbb{N}, R) . Define

$$f(R) := \begin{cases} \alpha & \text{if } (\mathbb{N}, R) \cong (\alpha, \varepsilon) \\ 0 & \text{o/w} \end{cases}$$

Claim: f is a surjection, if $\omega \leq \alpha < \omega_1$,
 there is a bij. $g: \omega \rightarrow \alpha$. Define

$$(u, v) \in R_g \iff g(u) \in g(v). \quad (\alpha, \varepsilon) \cong (\mathbb{N}, R_g). \quad \text{Q.e.d.}$$

Hilbert #1

$$2^{\aleph_0} = \aleph_1 ?$$

CH

Solution to Hilbert #1:

Two parts:

- ① CH is consistent
- ② \neg CH is consistent

Gödel 1938
(will be proved in this course)

Cohen 1962
(will not be proved in this course)

INDEPENDENCE

(the truth value of CH is not determined by ZF).

GRENZEN DES NATURERKENNENS

IGNORABIMUS

Hilbert Radio Address 1930

Wir müssen wissen,
wir werden wissen.

Kongress deutscher Naturforscher & Ärzte
Königsberg 1930: Hilbert was made
Honorary Citizen of Königsberg &
gave a national radio address.

Emil du Bois-Reymond



Born	Emil Heinrich du Bois-Reymond 7 November 1818 Berlin, Kingdom of Prussia
Died	26 December 1896 (aged 78) Berlin, Germany
Nationality	German

Emil: older
brother of
Paul

[matrilineal
who first defined
Cantor's
Ziffern]



Paul David Gustav du Bois-Reymond.

1931:

Gödel's Incompleteness
Theorem shows that
these vesicles be
an IGNORABIMUS.



CH:

CH is consistent
(Gödel 1938)

Main goal of this course
(7 or 8 lectures)

INNER MODELS

Then if $(M, \varepsilon) \models ZF$ then there is $N \subseteq M$ s.t. $(N, \varepsilon) \models ZF + AC + CH$.

In fact, this N is going to be the "minimal" inner model.

$\neg CH$ is consistent
(Cohen 1962)

OUTER MODELS
("Forcing")

Then if $(M, \varepsilon) \models ZFC$ which is countable & transitive, then there is $N \supseteq M$ such that $N \models ZFC + \neg CH$.

This is the topic of a different course.

Kurt Gödel



Gödel around 1926

Born Kurt Friedrich Gödel
April 28, 1906
Brünn, Austria-Hungary
Died January 14, 1978 (aged 71)
Princeton, New Jersey, U.S.

Paul Cohen <

American mathematician



Paul Joseph Cohen was an American mathematician. He is best known for his proofs that the continuum hypothesis and the axiom of choice are independent from Zermelo–Fraenkel set theory, for which he was awarded a Fields Medal. [Wikipedia](#)

Born: 2 April 1934, Long Branch, New Jersey, United States

Died: 23 March 2007, Stanford, California, United States

Known for: Cohen forcing; Continuum hypothesis

Fields: Mathematics

§ 2 Substructures

\mathcal{L} first-order language

R relation symbol

f function symbol

c constant symbols

(M, α)

\mathcal{L} -structure

where $\alpha(R)$ is a relation
 $\alpha(f)$ is a function
 $\alpha(c) \in M$

$N \subseteq M$ substructure if

(a) for every c , $\alpha(c) \in N$

(b) for every f , n -ary, $x_1, \dots, x_n \in N$
 $\alpha(f)(x_1, \dots, x_n) \in N$

More precisely,

(N, α')

$$\alpha'(c) := \alpha(c)$$

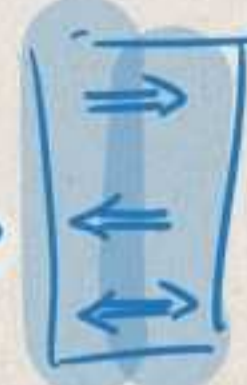
$$\alpha'(f) := \alpha(f) \cap N^n \times N$$

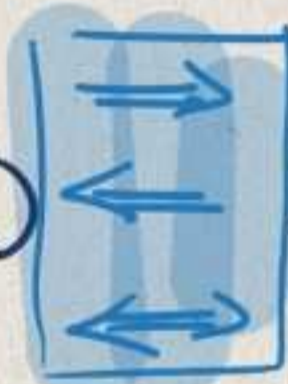
$$\alpha'(R) := \alpha(R) \cap N^n$$

Usually, we write (N, α) for this.

This is literally the notion of substructure that we know from other parts of math.

Def. Let φ be a formula, (M, α) an α -structure and $N \subseteq M$, we say that

n free variable φ is
 { upwards absolute
 downwards absolute
 absolute }  between N & M

if $\forall x_1, \dots, x_n \in N$
 $(N, \alpha) \models \varphi(x_1, \dots, x_n) \iff (M, \alpha) \models \varphi(x_1, \dots, x_n)$ 

Def.

Atomic formulas

Propositional formulas

$t = t'$
 or $R(t_1, \dots, t_n)$

Closure of atomic under $\wedge, \vee, \neg, \rightarrow,$

\leftrightarrow

if $\exists x \psi(x)$ ψ is prop.

$\forall x \psi(x)$ ψ is prop.

$\exists x \forall y \psi(x, y)$ ψ is prop.

$\forall x \exists y \psi(x, y)$ ψ is prop.

\exists -formula

\forall -formula

$\exists \forall$ -formula

$\forall \exists$ -formula

Proposition All atomic & propositional formulas are absolute between N & M

if N is a substructure.

Not so for \exists, \forall -formulas:

\mathcal{L} language of group theory $+, 0$

$\exists x (x \neq 0) \iff \varphi_{NT}$

$\underbrace{\quad}_{\text{propositional}}$
 $\underbrace{\quad}_{\exists\text{-formula}}$

Take $\{0\} \leq \mathbb{Z}$

have $\{0\} \not\models \varphi_{NT}$ & $\mathbb{Z} \models \varphi_{NT}$.

Similarly: $\forall x (x = 0) \iff \varphi_T$

$\{0\} \models \varphi_T$ & $\mathbb{Z} \not\models \varphi_T$.

\forall -formula

Proposition (1) If φ is upwards absolute between $N \& M$, then so is $\exists x \varphi$.

(2) If φ is downwards absolute betw. $N \& M$, then so is $\forall x \varphi$.

So \exists -formulas are upwards & \forall -formulas are downwards abs. for substructures.

$\forall\exists$ -formulas

$$\forall x \exists y (y + y = x) = \varphi_2$$

Everything is divisible by two.

$$\mathbb{Z} \leq \mathbb{Q} \leq \mathbb{Q} \times \mathbb{Z}$$

$$\begin{matrix} \pi & \pi & \pi \\ \neg\varphi_2 & \varphi_2 & \neg\varphi_2 \end{matrix}$$

Thus φ_2 is neither upwards nor downwards absolute.

Def.

A substructure N is called an elementary substructure if all formulas are absolute.

$$N \preceq M$$

Reminder

The isomorphism theorem states:
 If $f: (N, \alpha) \cong (M, \alpha)$ is an isomorphism,
 then $\{ \varphi; \varphi \text{ is sentence } \& N \models \varphi \} =$
 $\{ \varphi; \varphi \text{ is sentence } \& M \models \varphi \}$
 $\text{Th}(N) = \text{Th}(M)$

We say that N & M are elementarily equivalent.

In symbols: $N \equiv M$.

The proof is a proof by induction
& proves

$$(N, \alpha') \models \varphi(x_1, \dots, x_n) \iff (M, \alpha) \models \varphi(f(x_1), \dots, f(x_n))$$

This property is called:

f is an elementary embedding

Thus: $N \preceq M \iff \text{id}: N \rightarrow M$ is an elementary embedding.

§3 The language of set theory

$$L_{\in} = \{ \in \}$$

So, no constant or function symbols. That means: no non-trivial atomic formulas and thus almost nothing is absolute.

$$x=0 \iff x = \emptyset \iff \forall z (z \neq x) \iff \bar{\Phi}_0(x)$$

V-formula

So, not necessarily absolute.

$$x=1 \iff x = \{ \emptyset \} \iff \forall z (z \in x \iff \bar{\Phi}_0(z)) \iff \bar{\Phi}_1(x)$$

Example Let (M, E) be any \mathcal{L}_E -structure

with

$$(M, E) \models \bar{\Phi}_0(m_0)$$

$$(M, E) \models \bar{\Phi}_1(m_1)$$

Form $M' := M \setminus \{m_0\}$.

$$(M', E) \models \bar{\Phi}_0(m_1)$$

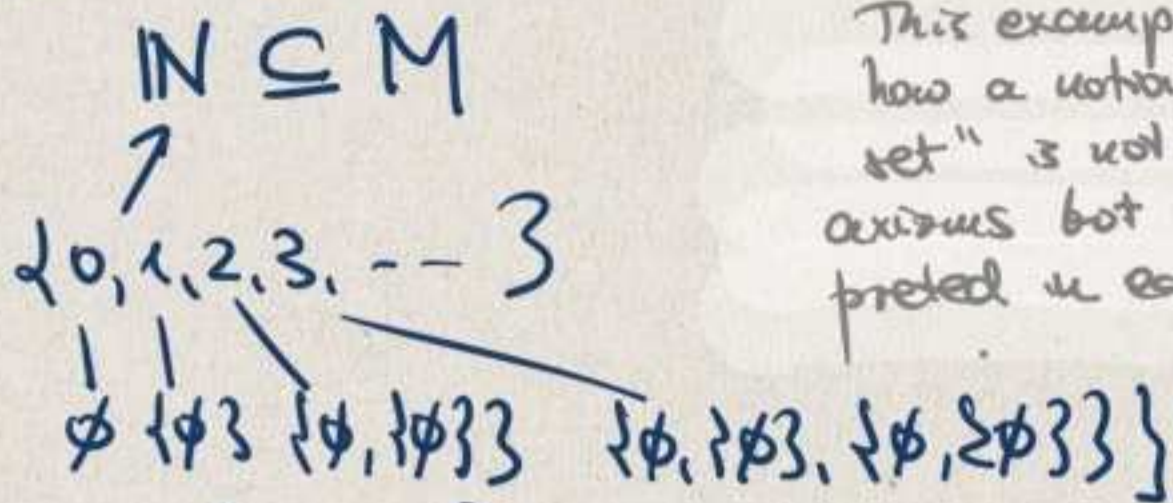
$$\forall z (z \neq m_1)$$

$$(M', E) \models \neg \bar{\Phi}_1(m_1)$$

This: Very basic properties of set theory
are not preserved by \mathcal{L}_E -sub-
structures.

Message Don't expect a formula to mean
what you think it means.

Example $(M, \varepsilon) \models ZFC$



This example illustrates how a notion like "power set" is not fixed by the axioms but will be re-interpreted in each model.

Consider (N, ε) . Which axioms of ZF does (N, ε) satisfy? For instance

$$\exists x \Phi_0(x)$$

$$\forall x \forall y \exists z \forall z' (z \varepsilon z' \leftrightarrow z \varepsilon x \vee z \varepsilon y)$$

What about power set?

$$\forall x \exists p \forall z (z \in p \leftrightarrow \forall w (w \varepsilon z \Rightarrow w \in x))$$

If $k, n \in \mathbb{N}$, when is $k \subseteq n$?

$$\text{iff } k \leq n.$$

So the "(\mathbb{N})-power set" of n is

$$\{k; k \leq n\} = n + 1.$$

So $(N, \varepsilon) \models$ Power set

and powersets do not always have card. 2^k .

Def. If (M, E) is an \mathcal{L}_E -structure, we say that $N \subseteq M$ is a **transitive substructure** if for all $x \in N$ & $y \in M$, if $y \in Ex$, then $y \in N$.

[That's just a model-theoretic rephrasing of the standard notion of "transitive set".]

If $N \subseteq M$ is transitive, then Φ_0 is absolute between N & M .

Proof. Since Φ_0 was \forall -formula, it is downwards absolute. Thus only need to show upwards.

Suppose not:

$$(N, E) \models \Phi_0(x) \quad \& \quad (M, E) \models \neg \Phi_0(x)$$

there is $y \in M$ s.t.
 $y \in Ex$

$y \in N$

$$\implies (N, E) \models \neg \Phi_0(x).$$

by transitivity Contradiction! q.e.d.