

TREE-MINIMAL GRAPHS ARE ALMOST REGULAR

D. DELLAMONICA JR., P. HAXELL, T. ŁUCZAK, D. MUBAYI, B. NAGLE,
Y. PERSON, V. RÖDL, AND M. SCHACHT

ABSTRACT. For all fixed trees T and any graph G we derive a counting formula for the number $\hat{N}_T(G)$ of homomorphisms from T to G in terms of the degree sequence of G .

As a consequence we obtain that any n -vertex graph G with edge density $p = p(n) \gg n^{-1/(t-2)}$, which contains at most $(1 + o(1))p^{t-1}n^t$ copies of some fixed tree T with $t \geq 3$ vertices must be almost regular, i.e., $\sum_{v \in V(G)} |\deg(v) - pn| = o(pn^2)$.

1. INTRODUCTION

For graphs F and G , let $N_F(G)$ denote the number of labeled copies of F in G . A well-known conjecture, due to Erdős and Simonovits (see [8]) and Sidorenko [6, 7], asserts that for every bipartite graph F and every $p > 0$, if an n -vertex graph G contains at least $p\binom{n}{2}$ edges, then

$$N_F(G) \geq (1 - o(1))p^{e_F}n^{v_F}. \quad (1)$$

This conjecture is known to hold for several classes of graphs F , including forests, even cycles, and complete bipartite graphs [7], Boolean cubes [5] and bipartite graphs F which contain a vertex that is connected to every vertex in the other vertex class [3].

The bound in (1) is asymptotically best possible, as for example the random graph $G(n, p)$ and, more generally, *quasi-random graphs* of density p show (see, e.g. [2, 10]). Skokan and Thoma [9] asked to what extent the converse holds, and posed the following problem: *if F is a bipartite graph, but not a forest, and if $N_F(G) \leq (1 + o(1))p^{e_F}n^{v_F}$, does G have to be quasi-random, in the sense of [2]?* It is believed that the answer to this question is positive, and this is known as the *forcing conjecture for quasi-random graphs*.

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The exclusion of forests F in the forcing conjecture is clearly necessary, as there are examples of regular graphs G which (therefore) minimize $N_F(G)$ but which are not quasi-random. We study this case and address the following problem :

*What structure on G is enforced when $N_T(G) \leq (1+o(1))p^{t-1}n^t$
for a forest T with t vertices?*

Since $N_T(G) = (1 + o(1))N_{T'}(G) \cdot N_{T''}(G)$, when T is the (vertex) disjoint union of T' and T'' , it suffices to consider connected graphs T only. Moreover, the case when T consists of only one edge is rather uninteresting, as in this case $N_T(G) = 2|E(G)|$. Consequently, from now on we restrict ourselves to trees T with at least three vertices.

For an n -vertex graph $G = (V, E)$ and $v \in V$, let d_v denote the degree of v in G and let $d = d(G) = \frac{1}{n} \sum_{v \in V} d_v$ denote the average degree of G . Clearly, if $|E| \geq p \binom{n}{2}$, then $d(G) \geq p(n-1)$. It is easy to see that for every forest T , and every $p = p(n) \gg 1/n$, the following holds: if $G = (V, E)$ is an *almost regular* n -vertex graph, that is, $\sum_{v \in V} |d_v - d| \leq o(pn^2)$ and $|E| = p \binom{n}{2}$, then $N_T(G) \leq (1 + o(1))p^{t-1}n^t$. Our first result provides a converse.

Theorem 1. *For all $\varepsilon \in (0, 1]$ and integers $t \geq 3$, there exists $\delta > 0$ so that whenever $p \gg n^{-1/(t-2)}$, the following statement holds. Let T be a tree with t vertices, and let $G = (V, E)$ be a graph with n vertices and $|E| = p \binom{n}{2}$ edges. If $N_T(G) \leq (1 + \delta)p^{t-1}n^t$, then $\sum_{v \in V} |d_v - d| \leq \varepsilon pn^2$.*

For the proof of Theorem 1 we consider not only copies of T , but more generally, homomorphisms (edge-preserving maps) $h: V(T) \rightarrow V(G)$. Let $\hat{N}_T(G)$ denote the number of such homomorphisms. Note that

$$N_T(G) \leq \hat{N}_T(G) \leq N_T(G) + O(pn^{t-1}). \quad (2)$$

To see the error term $O(pn^{t-1})$ above, consider a non-injective homomorphic image S of T in G . The image S contains at least one edge (of which there are $p \binom{n}{2}$) and at most $t - 3$ other vertices.

The bound in (2) restricts the proof of Theorem 1 to $p \gg n^{-1/(t-2)}$. We believe Theorem 1 is valid for much smaller values of p .

Problem 2. *Does Theorem 1 remain true as long as $p \gg 1/n$?*

Theorem 1 is a consequence of the main result of the paper, which extends a result of Alon, Hoory and Linial [1] from paths to arbitrary trees.

Theorem 3. *Let T be a tree with $t \geq 3$ vertices, and let $G = (V, E)$ be a graph with n vertices and average degree d . Then,*

$$\hat{N}_T(G) \geq nd \prod_{v \in V} d_v^{\frac{(t-2)d_v}{nd}}.$$

We prove Theorem 3 in Section 3 using ideas from [1].

We believe that the counting formula for homomorphisms of T in G given in Theorem 3 remains more or less valid for injective homomorphisms of T

in G (i.e. labeled copies of T in G), provided the minimum degree $\delta(G)$ of G is sufficiently large.

Problem 4. *Prove a lower bound for $N_T(G)$, the number of labeled copies of a tree T in a graph G with sufficiently large $\delta(G)$, which would be analogous to that proven for $\hat{N}_T(G)$ in Theorem 3.*

In fact, the proof of Theorem 3 can be altered to give such a result when T is the path P_3 with three edges. We will sketch the proof of Theorem 5 below in Section 5.

Theorem 5. *Let P_3 denote the path with 3 edges and let G be a graph with n vertices, average degree d and minimum degree at least 3. Then*

$$N_{P_3}(G) \geq nd \prod_{v \in V} (d_v - 2)^{\frac{2d_v}{nd}}.$$

Problem 4 and Theorem 5 are related to a result of Erdős and Simonovits [4]. They proved that the number of walks of length ℓ that are not paths in G is a negligible proportion of the total number of walks of length ℓ as the average degree d of G goes to infinity. In our notation,

$$\hat{N}_{P_\ell}(G) - N_{P_\ell}(G) = o(\hat{N}_{P_\ell}(G)) \quad \text{as } d \rightarrow \infty.$$

Together with the special case of Theorem 3 when $T = P_\ell$, which was earlier proved in [1], the Erdős-Simonovits result implies that

$$N_{P_\ell}(G) \geq (1 - o(1)) nd \prod_{v \in V(G)} d_v^{\frac{(\ell-1)d_v}{nd}} \quad \text{as } d \rightarrow \infty. \quad (3)$$

The dependence of the $o(1)$ term above on d was not very good (and not even really made explicit in [4]). Theorem 5 can thus be viewed as an improvement of (3) in the case $t = 3$.

2. PROOF OF THEOREM 1

In this section, we deduce Theorem 1 from Theorem 3. For that, we will also use the following consequence of Jensen's inequality.

Lemma 6. *Let $G = (V, E)$ be a graph with $|V| = n$ vertices and average degree $d = d(G)$. Suppose*

$$\frac{1}{n} \sum_{v \in V} d_v \log d_v < \frac{\gamma^2}{\gamma + 2} d + d \log d \quad (4)$$

for some $\gamma > 0$. Then

$$\frac{1}{n} \sum_{v \in V} \left| \frac{d_v}{d} - 1 \right| < \gamma.$$

For completeness, we include a proof of Lemma 6 at the end of the section and first we deduce Theorem 1 from Theorem 3 and Lemma 6.

Proof of Theorem 1. Let $\varepsilon \in (0, 1]$ and $t \geq 3$ be given. With $\gamma = \varepsilon$, let $\delta = \gamma^2/(3(\gamma + 2))$, and let $p \gg n^{-1/(t-2)}$. Let T be a tree with t vertices and let $G = (V, E)$ be a graph with $|V| = n$ vertices, with $|E| = p\binom{n}{2}$ edges, and average degree $d = p(n-1)$. Suppose

$$N_T(G) \leq (1 + \delta)p^{t-1}n^t.$$

From Theorem 3 and (2), we infer

$$\prod_{v \in V} d_v^{\frac{(t-2)d_v}{nd}} \leq \frac{\hat{N}_T(G)}{nd} \leq \frac{1}{nd} ((1 + \delta)p^{t-1}n^t + O(pn^{t-1})).$$

Since $p \gg n^{-1/(t-2)}$, we have $p^{t-1}n^t \gg pn^{t-1}$, and therefore,

$$\prod_{v \in V} d_v^{\frac{(t-2)d_v}{nd}} \leq \frac{(1 + 2\delta)}{nd} p^{t-1}n^t \leq (1 + 3\delta)d^{t-2}.$$

Taking natural logarithms then yields

$$\frac{t-2}{nd} \sum_{v \in V} d_v \log d_v \leq 3\delta + (t-2) \log d,$$

which implies,

$$\frac{1}{n} \sum_{v \in V} d_v \log d_v \leq \frac{3\delta}{t-2} d + d \log d \leq \frac{\gamma^2}{\gamma+2} d + d \log d. \quad (5)$$

Therefore, Lemma 6 and our choice of constants yield

$$\sum_{v \in V} |d_v - d| \leq \gamma dn \leq \varepsilon pn^2,$$

which proves Theorem 1. \square

Proof of Lemma 6. We proceed by contradiction. In fact, since the RHS of (4) is increasing in γ , we may assume

$$\frac{1}{n} \sum_{v \in V} \left| \frac{d_v}{d} - 1 \right| = \gamma. \quad (6)$$

Observe that (4) may be rewritten as

$$\frac{1}{n} \sum_{v \in V} \frac{d_v}{d} \log \frac{d_v}{d} < \frac{\gamma^2}{\gamma+2}.$$

We show, however, that the assumption in (6) implies

$$\frac{1}{n} \sum_{v \in V} \frac{d_v}{d} \log \frac{d_v}{d} \geq \frac{\gamma^2}{\gamma+2}, \quad (7)$$

a contradiction. It remains to prove (7).

Set

$$W = \{v \in V : d_v > d\},$$

and write, as usual, $W^C = V \setminus W$ for the complement. Since $\sum_{v \in V} (d_v - d) = 0$, it follows that

$$\sum_{v \in W} \left(\frac{d_v}{d} - 1 \right) = \sum_{v \in W^C} \left(1 - \frac{d_v}{d} \right).$$

Consequently, (6) implies that

$$\sum_{v \in W} \left(\frac{d_v}{d} - 1 \right) = \frac{\gamma n}{2} = \sum_{v \in W^C} \left(1 - \frac{d_v}{d} \right). \quad (8)$$

Write $\alpha_1 = |W|/n$ and $\alpha_2 = 1 - \alpha_1 = |W^C|/n$.

Setting $\phi(x) = x \log x$ (a convex function), Jensen's inequality implies that

$$\frac{1}{|W|} \sum_{v \in W} \phi\left(\frac{d_v}{d}\right) \geq \phi\left(\frac{1}{|W|} \sum_{v \in W} \frac{d_v}{d}\right) \stackrel{(8)}{=} \phi\left(1 + \frac{\gamma}{2\alpha_1}\right).$$

Similarly,

$$\frac{1}{|W^C|} \sum_{v \in W^C} \phi\left(\frac{d_v}{d}\right) \geq \phi\left(\frac{1}{|W^C|} \sum_{v \in W^C} \frac{d_v}{d}\right) \stackrel{(8)}{=} \phi\left(1 - \frac{\gamma}{2\alpha_2}\right).$$

From the two inequalities above, we obtain

$$\frac{1}{n} \sum_{v \in V} \frac{d_v}{d} \log \frac{d_v}{d} \geq \alpha_1 \phi\left(1 + \frac{\gamma}{2\alpha_1}\right) + \alpha_2 \phi\left(1 - \frac{\gamma}{2\alpha_2}\right). \quad (9)$$

Now, set

$$y_1 = 1 + \frac{\gamma}{2\alpha_1} \quad \text{and} \quad y_2 = 1 - \frac{\gamma}{2\alpha_2}, \quad (10)$$

and observe that $\alpha_1 y_1 + \alpha_2 y_2 = 1$. By Taylor's theorem, there exist reals $\xi_1 \in (1, y_1)$ and $\xi_2 \in (y_2, 1)$ such that for $i = 1, 2$, we have

$$\phi(y_i) = \phi(1) + (y_i - 1)\phi'(1) + \frac{(y_i - 1)^2}{2}\phi''(\xi_i).$$

Since $\phi(1) = 0$, $\phi'(1) = 1$, and $\phi''(\xi_i) = 1/\xi_i$, we have

$$\begin{aligned} \alpha_1 \phi(y_1) + \alpha_2 \phi(y_2) &= \frac{\alpha_1}{2\xi_1}(y_1 - 1)^2 + \frac{\alpha_2}{2\xi_2}(y_2 - 1)^2 \\ &> \frac{\alpha_1}{2y_1}(y_1 - 1)^2 + \frac{\alpha_2}{2}(y_2 - 1)^2 \\ &\stackrel{(10)}{=} \frac{\gamma^2}{4(\gamma + 2\alpha_1)} + \frac{\gamma^2}{8\alpha_2} \\ &= \frac{\gamma^2(\gamma + 2)}{8(\gamma + 2\alpha_1)(1 - \alpha_1)}. \end{aligned} \quad (11)$$

The denominator is maximized with $\alpha_1 = 1/2 - \gamma/4$, there equaling $(\gamma + 2)^2$. Hence

$$\frac{1}{n} \sum_{v \in V} \frac{d_v}{d} \log \frac{d_v}{d} \stackrel{(9)}{\geq} \alpha_1 \phi(y_1) + \alpha_2 \phi(y_2) \stackrel{(11)}{\geq} \frac{\gamma^2}{\gamma + 2},$$

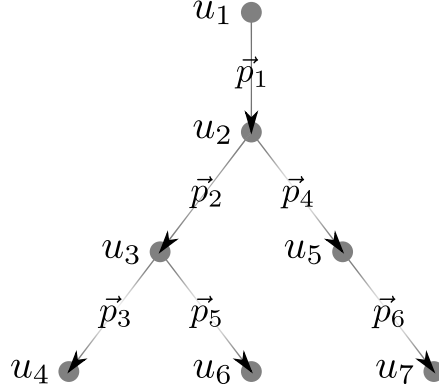


FIGURE 1. Tree labeling and orientation

proving (7). □

3. PROOF OF THEOREM 3

Let $T = (U, P)$ be a tree with $t \geq 3$ vertices, and let $G = (V, E)$ be a graph with $|V| = n$ vertices and average degree $d = d(G)$. We begin the proof by describing some notation.

We will consider the digraph $\vec{G} = (V, \vec{E})$ obtained from G by including, for each edge $\{v, v'\} \in E$, both arcs (v, v') and (v', v) into \vec{E} .

Moreover, we will label the vertices and orient the edges of T . For that we fix a leaf u_1 of T and let u_2 be the unique neighbor of u_1 in T . Furthermore, fix a labeling u_3, \dots, u_t of $U \setminus \{u_1, u_2\}$ such that, for every $i = 1, \dots, t$, the induced subgraph $T_i = T[\{u_1, \dots, u_i\}]$ is connected. Let $P = \{p_1, \dots, p_{t-1}\}$ be the corresponding labeling of the edges of T defined by the property that $E(T_i) = \{p_1, \dots, p_{i-1}\}$ for every $i = 2, \dots, t$. For every edge $p = \{u_i, u_j\}$ with $i < j$, we denote by \vec{p} the oriented pair (u_i, u_j) . Note that, from these definitions, it follows that, for every $j = 1, \dots, t-1$, we have $\vec{p}_j = (u_i, u_{j+1})$ for some $i < j$. Finally, we denote by $\vec{T} = (U, \vec{P})$ the oriented tree with $\vec{P} = \{\vec{p}: p \in P\}$ (see Figure 1).

Now, let $\Omega = \Omega_T(G)$ denote the family of all homomorphisms $\omega: U \rightarrow V$ (so that $|\Omega| = \hat{N}_T(G)$). We develop some notation for Ω . To begin, if $\vec{p} = (u, u') \in \vec{P}$, $\vec{e} = (v, v') \in \vec{E}$, and $\omega \in \Omega$, we write $\omega(\vec{p}) = \vec{e}$ if $\omega(u) = v$ and $\omega(u') = v'$. In this way, we may view the image of an element $\omega \in \Omega$ as a $(t-1)$ -tuple $\omega = (\omega(\vec{p}_1), \dots, \omega(\vec{p}_{t-1}))$. For $\vec{e}_1 \in \vec{E}$, define $\Omega_{\vec{e}_1} \subseteq \Omega$ to be the set of those homomorphisms $\omega \in \Omega$ for which $\omega(\vec{p}_1) = \vec{e}_1$. For $\omega \in \Omega_{\vec{e}_1}$, we write $\omega_- = (\omega(\vec{p}_2), \dots, \omega(\vec{p}_{t-1}))$. For an arc $(v, v') \in \vec{E}$, we shall sometimes abuse notation and write $(v, v') \in \omega_-$ to mean $(v, v') \in \{\omega(\vec{p}_2), \dots, \omega(\vec{p}_{t-1})\}$, and we denote by $m_{\omega_-}(v, v')$ the number of appearances of $(v, v') \in \{\omega(\vec{p}_2), \dots, \omega(\vec{p}_{t-1})\}$, i.e.,

$$m_{\omega_-}(v, v') = |\{j = 2, \dots, t-1: \omega(\vec{p}_j) = (v, v')\}|.$$

We consider the following procedure for generating a random element $\omega = \{\omega(\vec{p}_1), \dots, \omega(\vec{p}_{t-1})\} \in \Omega$. Select $\vec{e}_1 = \omega(\vec{p}_1) \in \vec{E}$ uniformly at random. For $2 \leq j < t$, suppose $\vec{e}_1 = \omega(\vec{p}_1), \dots, \vec{e}_{j-1} = \omega(\vec{p}_{j-1}) \in \vec{E}$ have been selected. Then, choose $\vec{e}_j = \omega(\vec{p}_j) = (v, v') \in \vec{E}$ by selecting v' uniformly at random from the set of neighbors of v in G . Moreover, a random element $\omega \in \Omega_{\vec{e}_1}$ is generated similarly, only that we condition on $\omega(\vec{p}_1) = \vec{e}_1$.

For fixed $\vec{e}_1 = (v_1, v_2) \in \vec{E}$ and fixed $\omega \in \Omega_{\vec{e}_1}$, we denote by $\mathbb{P}_{\vec{e}_1}(\omega)$ the probability of $\mathbb{P}(\omega = \omega)$, for a random homomorphism ω from the probability space $\Omega_{\vec{e}_1}$. It follows from the definition of $\Omega_{\vec{e}_1}$ that for every $\vec{e}_1 \in \vec{E}$ and every $\omega = (\vec{e}_1 = (v_1, v'_1), \dots, \vec{e}_{t-1} = (v_{t-1}, v'_{t-1})) \in \Omega$, we have

$$\mathbb{P}_{\vec{e}_1}(\omega) = \prod_{j=2}^{t-1} \frac{1}{d_{v_j}} = \prod_{(v, v') \in \omega_-} \left(\frac{1}{d_v} \right)^{m_{\omega_-}(v, v')}. \quad (12)$$

We now estimate the quantity $\hat{N}_T(G) = |\Omega| = \sum_{\vec{e}_1 \in \vec{E}} |\Omega_{\vec{e}_1}|$. To that end, observe that

$$\frac{|\Omega|}{nd} = \sum_{\vec{e}_1 \in \vec{E}} \frac{|\Omega_{\vec{e}_1}|}{nd} = \sum_{\vec{e}_1 \in \vec{E}} \sum_{\omega \in \Omega_{\vec{e}_1}} \frac{\mathbb{P}_{\vec{e}_1}(\omega)}{nd} \left(\frac{1}{\mathbb{P}_{\vec{e}_1}(\omega)} \right),$$

where $\sum_{\vec{e}_1 \in \vec{E}} \sum_{\omega \in \Omega_{\vec{e}_1}} (\mathbb{P}_{\vec{e}_1}(\omega)/(nd)) = 1$. Applying the Arithmetic-Geometric Mean Inequality¹ and using (12), we obtain

$$\begin{aligned} \frac{|\Omega|}{nd} &= \sum_{\vec{e}_1 \in \vec{E}} \sum_{\omega \in \Omega_{\vec{e}_1}} \frac{\mathbb{P}_{\vec{e}_1}(\omega)}{nd} \left(\frac{1}{\mathbb{P}_{\vec{e}_1}(\omega)} \right) \\ &\geq \prod_{\vec{e}_1 \in \vec{E}} \prod_{\omega \in \Omega_{\vec{e}_1}} \left(\frac{1}{\mathbb{P}_{\vec{e}_1}(\omega)} \right)^{\mathbb{P}_{\vec{e}_1}(\omega)/(nd)} \\ &\stackrel{(12)}{=} \prod_{\vec{e}_1 \in \vec{E}} \prod_{\omega \in \Omega_{\vec{e}_1}} \prod_{(v, v') \in \omega_-} d_v^{m_{\omega_-}(v, v') \mathbb{P}_{\vec{e}_1}(\omega)/(nd)}. \end{aligned} \quad (13)$$

For a fixed $\vec{e} = (v, v') \in \vec{E}$, we now collect the terms $d_v^{\mathbb{P}_{\vec{e}_1}(\omega)/(nd)}$ in the triple product above. To that end, observe that a factor $d_v^{\mathbb{P}_{\vec{e}_1}(\omega)/(nd)}$ appears for every $\vec{e}_1 \in \vec{E}$, $\omega \in \Omega_{\vec{e}_1}$ and $2 \leq s \leq t-1$ for which $\omega(\vec{p}_s) = \vec{e}$. Set, therefore, for $s = 2, \dots, t-1$ and $\vec{e} = (v, v') \in \vec{E}$

$$g_s(\vec{e}) = \frac{1}{nd} \sum_{\vec{e}_1 \in \vec{E}} \sum_{\omega \in \Omega_{\vec{e}_1}} \{ \mathbb{P}_{\vec{e}_1}(\omega) : \omega \in \Omega_{\vec{e}_1} \text{ and } \omega(\vec{p}_s) = \vec{e} \} \quad (14)$$

¹The AM-GM inequality states that for positive x_1, \dots, x_m and non-negative c_1, \dots, c_m , where $c_1 + \dots + c_m = 1$,

$$\sum_{i=1}^m c_i x_i \geq \prod_{i=1}^m x_i^{c_i}.$$

so that (13) is equivalently

$$|\Omega| \geq nd \prod_{\vec{e}=(v,v') \in \vec{E}} d_v^{\sum_{s=2}^{t-1} g_s(\vec{e})}. \quad (15)$$

In a moment, we prove that for a fixed $\vec{e} \in \vec{E}$ and $2 \leq s \leq t-1$,

$$g_s(\vec{e}) = \frac{1}{nd}. \quad (16)$$

Applying (16) to (15) yields

$$|\Omega| \geq nd \prod_{(v,v') \in \vec{E}} d_v^{\frac{t-2}{nd}} = nd \prod_{v \in V} d_v^{\frac{(t-2)d_v}{nd}},$$

as promised by Theorem 3. It remains to prove (16).

Proof of (16). Fix $\vec{e} \in \vec{E}$ and an integer s between 2 and $t-1$. We claim that $g_s(\vec{e})$ is the probability that \vec{e} will appear at the s -th step of a randomly generated $\omega \in \Omega$, i.e., $g_s(\vec{e}) = \mathbb{P}(\omega(\vec{p}_s) = \vec{e})$ (where this probability occurs in the space Ω). Indeed,

$$\begin{aligned} \mathbb{P}(\omega(\vec{p}_s) = \vec{e}) &= \sum_{\vec{e}_1 \in \vec{E}} \mathbb{P}(\omega(\vec{p}_s) = \vec{e} \mid \omega(\vec{p}_1) = \vec{e}_1) \mathbb{P}(\omega(\vec{p}_1) = \vec{e}_1) \\ &= \frac{1}{nd} \sum_{\vec{e}_1 \in \vec{E}} \mathbb{P}(\omega(\vec{p}_s) = \vec{e} \mid \omega(\vec{p}_1) = \vec{e}_1) \\ &= \frac{1}{nd} \sum_{\vec{e}_1 \in \vec{E}} \mathbb{P}_{\vec{e}_1}(\omega(\vec{p}_s) = \vec{e}) \\ &= \frac{1}{nd} \sum_{\vec{e}_1 \in \vec{E}} \sum \{ \mathbb{P}_{\vec{e}_1}(\omega) : \omega \in \Omega_{\vec{e}_1} \text{ and } \omega(\vec{p}_s) = \vec{e} \} \\ &= g_s(\vec{e}). \end{aligned} \quad (17)$$

We prove (16) by induction on s . We can easily extend the definition of $g_s(\vec{e})$ for the case $s=1$ by setting

$$g_1(\vec{e}) = \frac{1}{nd} \sum_{\vec{e}_1 \in \vec{E}} \sum \{ \mathbb{P}_{\vec{e}_1}(\omega) : \omega \in \Omega_{\vec{e}_1} \text{ and } \omega(\vec{p}_1) = \vec{e} \} \quad (18)$$

and, clearly, $g_1(\vec{e}) = 1/(nd)$, which establishes the base case.

Assume, for each $1 \leq r < s$ and each $\vec{e}' \in \vec{E}$, that $g_r(\vec{e}') = 1/(nd)$. Now, let $\vec{p}_s = (u_r, u_{s+1}) \in \vec{P}$, for some $1 \leq r \leq s$. Consider the arc $\vec{p}_{r-1} = (u_q, u_r) \in \vec{P}$, where $1 \leq q \leq r-1$. Observe that $\omega(\vec{p}_s) = \vec{e} = (v, v')$ only if $\omega(\vec{p}_{r-1}) = \vec{e}' = (v'', v)$ for some vertex v'' in the neighborhood $\Gamma_G(v)$

of v in the undirected graph G . Thus,

$$\begin{aligned} g_s(\vec{e}) &= \mathbb{P}(\boldsymbol{\omega}(\vec{p}_s) = \vec{e}) \\ &= \sum_{v'' \in \Gamma_G(v)} \mathbb{P}(\boldsymbol{\omega}(\vec{p}_s) = \vec{e} \mid \boldsymbol{\omega}(\vec{p}_{r-1}) = (v'', v)) \mathbb{P}(\boldsymbol{\omega}(\vec{p}_{r-1}) = (v'', v)). \end{aligned}$$

For every $v'' \in \Gamma_G(v)$, the induction hypothesis gives

$$\mathbb{P}(\boldsymbol{\omega}(\vec{p}_{r-1}) = (v'', v)) = \frac{1}{nd},$$

and by the definition of the probability space Ω , we have

$$\mathbb{P}(\boldsymbol{\omega}(\vec{p}_s) = \vec{e} \mid \boldsymbol{\omega}(\vec{p}_{r-1}) = (v'', v)) = \frac{1}{d_v}.$$

Thus,

$$g_s(\vec{e}) = |\Gamma_G(v)| \cdot \frac{1}{d_v} \cdot \frac{1}{nd} = \frac{1}{nd},$$

as desired. \square

4. PROOF OF THEOREM 5

The proof is similar to the proof of Theorem 3 and we only outline the differences here. We modify the probability space Ω to include only labeled copies of P_3 :

- (i) the first arc $\vec{e}_1 = (v_1, v_2)$ is chosen uniformly at random in \vec{E} ;
- (ii) the second arc $\vec{e}_2 = (v_2, v_3)$ is selected uniformly among all arcs $(v_2, v) \in \vec{E}$ with $v \neq v_1$;
- (iii) the third arc $\vec{e}_3 = (v_3, v_4)$ is selected uniformly among all arcs $(v_3, v) \in \vec{E}$ with $v \notin \{v_1, v_2\}$.

Note that the minimum degree condition of G is necessary to guarantee that the process above is feasible. Moreover, it is clear that $|\Omega| = N_{P_3}(G)$, since $N_{P_3}(G)$ denotes the number of labeled copies of P_3 in G .

In order to bound $|\Omega|$, we shall obtain an inequality along the lines of (13). Let $\omega = (\vec{e}_1 = (v_1, v_2), \vec{e}_2 = (v_2, v_3), \vec{e}_3 = (v_3, v_4)) \in \Omega$. Replacing (12) by the obvious bound

$$\mathbb{P}_{\vec{e}_1}(\omega) \geq \frac{1}{d_{v_2} - 2} \cdot \frac{1}{d_{v_3} - 2},$$

we obtain (using the arithmetic-geometric mean as before) the following analog of (13)

$$\frac{|\Omega|}{nd} \geq \prod_{\vec{e}_1 \in \vec{E}} \prod_{\omega \in \Omega_{\vec{e}_1}} \prod_{(v', v) \in \omega_-} (d_v - 2)^{\frac{\mathbb{P}_{\vec{e}_1}(\omega)}{nd}}, \quad (19)$$

where this time ω_- is the subpath (\vec{e}_1, \vec{e}_2) of the path $\omega = (\vec{e}_1, \vec{e}_2, \vec{e}_3)$. We remark that, in (19), there is a subtle difference in how we organized the products. In (19), the elements being multiplied are the degrees (minus two) of the *tail* vertex of each arc besides \vec{e}_3 , while in (13), they are the degrees

of the *head* vertex of each arc besides \vec{e}_1 . The reason for this subtle change, will become clear in a moment (see (20) below). Moreover, note that the $m_{\omega_-(v',v)}$ -term in the exponent of (13) does not appear in (19), since by definition of Ω no arc $\vec{e} \in \vec{E}$ can appear more than once in ω .

Let $g_s(\vec{e})$ for $s = 1, \dots, 3$, be defined as in (14) and (18). Note that the right hand side of (19) reduces to

$$\prod_{\vec{e}=(v',v) \in \vec{E}} (d_v - 2)^{g_1(\vec{e})+g_2(\vec{e})}.$$

The reason for changing the definition of ω_- and multiplying degrees of tails instead of heads above is that, as we will show below,

$$g_1(\vec{e}) = g_2(\vec{e}) = \frac{1}{nd} \quad (20)$$

while $g_3(\vec{e})$ may be different than $1/nd$ depending on the structure of G (triangles in G may force $g_3(\vec{e}) \neq 1/nd$). In fact, this is the reason the current method does not seem to apply to paths of length four or larger.

In order to verify (20), we first note that by definition of $g_1(\vec{e})$ in (18), we have $g_1(\vec{e}) = \frac{1}{nd}$. Moreover, (17) extends to this case here so that $g_2(\vec{e}) = \mathbb{P}(\omega(\vec{p}_2) = \vec{e})$ for every $\vec{e} \in \vec{E}$. The definition of Ω guarantees

$$\mathbb{P}(\omega(\vec{p}_2) = \vec{e} \mid \omega(\vec{p}_1) = (v'', v')) = \frac{1}{d_{v'} - 1}$$

for every $\vec{e} = (v', v) \in \vec{E}$ with $v \neq v''$. We therefore have

$$\begin{aligned} g_2(\vec{e}) &= \mathbb{P}(\omega(\vec{p}_2) = \vec{e}) \\ &= \sum_{v'' \in \Gamma_G(v') \setminus \{v\}} \mathbb{P}(\omega(\vec{p}_2) = \vec{e} \mid \omega(\vec{p}_1) = (v'', v')) \mathbb{P}(\omega(\vec{p}_1) = (v'', v')) \\ &= (d_{v'} - 1) \cdot \frac{1}{d_{v'} - 1} \cdot \frac{1}{nd} = \frac{1}{nd}, \end{aligned}$$

which proves (20).

Combining (19) and (20) we obtain

$$|\Omega| \geq nd \prod_{\vec{e}=(v',v) \in \vec{E}} (d_v - 2)^{\frac{2}{nd}} = nd \prod_{v \in V} (d_v - 2)^{\frac{2d_v}{nd}},$$

and Theorem 5 is proved. \square

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DEPARTMENT OF MATHEMATICS AND COMPUTER SCIENCE, EMORY UNIVERSITY, ATLANTA, GA 30322, USA

E-mail address: `ddellam@mathcs.emory.edu`

DEPARTMENT OF COMBINATORICS AND OPTIMIZATION, UNIVERSITY OF WATERLOO, WATERLOO, ON N2L 3G1, CANADA

E-mail address: `pehaxell@math.uwaterloo.ca`

DEPARTMENT OF DISCRETE MATHEMATICS, FACULTY OF MATHEMATICS AND COMPUTER SCIENCE, ADAM MICKIEWICZ UNIVERSITY, UMULTOWSKA 87, 61-614 POZNAŃ, POLAND AND DEPARTMENT OF MATHEMATICS AND COMPUTER SCIENCE, EMORY UNIVERSITY, ATLANTA, GA 30322, USA

E-mail address: `tomasz@amu.edu.pl`

DEPARTMENT OF MATHEMATICS, STATISTICS, AND COMPUTER SCIENCE, UNIVERSITY OF ILLINOIS, CHICAGO, IL 60607

E-mail address: `mubayi@math.uic.edu`

DEPARTMENT OF MATHEMATICS AND STATISTICS, UNIVERSITY OF SOUTH FLORIDA, TAMPA, FL 33620, USA

E-mail address: `bnagle@math.usf.edu`

INSTITUT FÜR INFORMATIK, HUMBOLDT-UNIVERSITÄT ZU BERLIN, UNTER DEN LINDEN 6, D-10099 BERLIN, GERMANY

E-mail address: `person@informatik.hu-berlin.de`

DEPARTMENT OF MATHEMATICS AND COMPUTER SCIENCE, EMORY UNIVERSITY, ATLANTA, GA 30322, USA

E-mail address: `rod1@mathcs.emory.edu`

FACHBEREICH MATHEMATIK, UNIVERSITÄT HAMBURG, BUNDESSTRASSE 55, D-20146 HAMBURG, GERMANY

E-mail address: `schacht@math.uni-hamburg.de`