



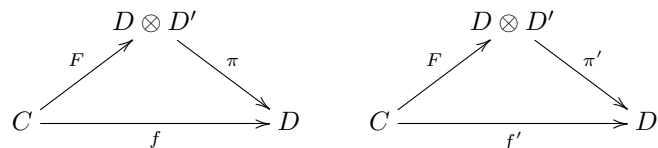
Sheet 5

In this sheet \mathbb{K} is a field, and all algebras, co-algebras and bialgebras are over \mathbb{K}

Problem 1. 1. Let H be a bialgebra and C be a sub-coalgebra of H , prove that the sub-algebra of H generated by C is a sub-bialgebra.

2. Show that a coalgebra C is cocommutative if and only if Δ is a coalgebra map.

3. Let C, D and D' be coalgebras such that C is cocommutative, and $f : C \rightarrow D$ and $f' : C \rightarrow D'$ two coalgebra maps. Define the canonical coalgebra maps $\pi : D \otimes D' \rightarrow D$ and $\pi' : D \otimes D' \rightarrow D$, and prove that there exists a unique map $F : C \rightarrow D \otimes D'$ such that the two following diagrams are commutatives.



Problem 2 (Rational modules, second part, see sheet 4). Let C be a coalgebra, and M and M' be two rational C^* -modules and L a C^* -module.

1. Prove that if N is a cyclic¹ sub-module of M then it is finite dimensional.

Solution. Let n be a generator of N . We write $\rho(n) = \sum_i n_i \otimes c_i$. Now if f is an element of C^* , $f \cdot n = \sum_i f(c_i)n_i$. Hence N is spanned as a vector space by the n_i and is therefore finite dimensional. \square

2. Prove that every finitely generated rational module is finite dimensional.

Solution. Same proof as before, with a finite set of generators instead of just one. \square

3. Prove that if a C^* -module N is a quotient of M , then it is rational.

Solution. Let us write $N = M/P$ and denote $\pi : M \rightarrow N$ the canonical projection. The module P is a submodule of M and hence is rational and is endowed with a comodule- C structure. There is a unique comodule- C structure on M/P making π a comodule- C morphism. From this comodule structure one can build an alternative (and rational) C^* -module structure. But then π is a C^* -module map for both C^* -module structure on N . This structure have to be equal since π is surjective. This prove that N is rational. \square

¹This a module generated by one element

4. Prove that $L^{\text{rat}} = \rho^{-1}(L \otimes C)$ is a rational sub-module. Prove that it contains all rational sub-modules of L .

Solution. Let $x \in L^{\text{rat}}$ and f in C^* . We want to show that $f \cdot x \in L^{\text{rat}}$, this means that $\rho(f \cdot x) \in L \otimes C$. Let us write $\rho(x) = \sum_i x_i \otimes c_i$. Let g be an element of C^* . We have:

$$\begin{aligned} g \cdot (f \cdot x) &= (gf) \cdot x \\ &= \sum_i x_i (gf)(c_i) \\ &= \sum_i x_i (g \otimes f)(\Delta(c_i)) \\ &= \sum_i \sum_{(c_i)} x_i g(c_{i(1)}) f(c_{i(2)}) \\ &= \sum_j x'_j g(c'_j). \end{aligned}$$

This means exactly that $\rho(f \cdot x) = \sum_j x'_j \otimes c'_j$. Hence, we have shown that L^{rat} is a C^* -module. It is trivial to show that $\rho(L^{\text{rat}} \subset L^{\text{rat}} \otimes C$ which means that L^{rat} is rational and that it contains any rational submodules of L . \square

5. Prove that $\phi : M \rightarrow M'$ is a C^* -module map iff f is a comodule- C map.

Solution. Suppose $\phi : M \rightarrow M'$ is a comodule- C morphism. Let f be an element of C^* and m an element of M . We write:

$$f \cdot m = \sum_{(m)} m_{(0)} f(m_{(1)}) \quad \text{and} \quad f \cdot \phi(m) = \sum_{(\phi(m))} \phi_{(0)} f(\phi(m)_{(1)}).$$

The map ϕ is suppose to be a map of comodule- C . This means:

$$\Delta(\phi(m)) \sum_{(\phi(m))} \phi(m)_{(0)} \otimes \phi(m)_{(1)} = \phi(m_{(0)}) \otimes m_{(1)} = (\phi \otimes \text{id}_C)(\Delta(m)).$$

Hence we have:

$$\begin{aligned} f \cdot \phi(m) &= \sum_{(\phi(m))} \phi(m)_{(0)} f(\phi(m)_{(1)}) \\ &= \sum_{(m)} \phi(m_{(0)}) f(m_{(1)}) \\ &= \phi \left(\sum_{(m)} m_{(0)} f(m_{(1)}) \right) \\ &= \phi(f \cdot m). \end{aligned}$$

Which means that ϕ is a morphism of C^* -modules.

Let us now suppose that ϕ is a C^* -module map. This implies that for any $m \in M$ and any $f \in C^*$, we have:

$$f \cdot \phi(m) = \sum_{(\phi(m))} \phi(m)_{(0)} f(\phi(m)_{(1)}) = \sum_{(m)} \phi(m_{(0)}) f(m_{(1)}) = \phi(f \cdot m).$$

This implies that $\Delta(\phi(m)) = \sum_{(m)} \phi(m_{(0)}) \otimes m_{(1)} = (\phi \otimes \text{id})(\Delta(m))$ and thus that ϕ is a comodule map. \square

Problem 3 (Fundamental theorem on Coalgebras). In this exercise C is a coalgebra and c an element of C .

1. Let $(C_i)_{i \in I}$ a collection of subcoalgebras of C . Prove that $\bigcap_{i \in I} C_i$ is a coalgebra. Let S be a subset of C , define the notion of *coalgebra generated by S* .

Solution. The first assertion is clear. For the second we simply consider the intersection of all sub-coalgebra of C containing S . \square

2. Recall how C is naturally endowed with a structure of C^* -module.

Solution. This is clear, C is even a rational. \square

3. Show that the sub- C^* -module N generated by c is finite dimensional (one should use the exercise about rational modules).

Solution. N is a submodule of a rational module, hence it is rational. Furthermore it is cyclic, so that it is finite dimensional. \square

4. (A little more difficult) Let $J = \{a \in C^* | a \cdot N = \{0\}\}$. Prove that J^\perp is a sub-coalgebra of C and that it is finite dimensional.

Solution. J is the kernel of $\pi : C^* \rightarrow \text{End}_{\mathbb{K}}(N)$ which is an algebra map. Hence we can deduce that J is a ideal of C^* with finite codimension. We claim that with this single hypothesis, we can deduce J^\perp is a finite dimensional sub-coalgebra of C : Indeed, one can check that if x is an element of J^\perp , then $\Delta(x)$ is in $J^\perp \otimes C \cap C \otimes J^\perp = J^\perp \otimes J^\perp$, so that J^\perp is a co-algebra.

Let us recall the definition of J^\perp and of $(J^\perp)^\perp$ and prove that $J \subset (J^\perp)^\perp$:

$$J^\perp = \{b \in C | f(b) = 0 \text{ for all } f \in J\} \quad \text{and} \quad (J^\perp)^\perp = \{g \in C^* | g(b) = 0 \text{ for all } b \in J^\perp\}.$$

Let f be an element of J and let b be an element of J^\perp . By definition of J^\perp , $f(b) = 0$. This is valid for every b , hence f is in $(J^\perp)^\perp$. The converse does not hold (J might not be a vector space for example). Let us prove that $\dim J^\perp = \text{codim} (J^\perp)^\perp$. Actually $V^* \simeq C^*/V^\perp$ holds for every sub-space V of C . If V is finite dimensional, this gives what we wanted. In the end, we have $\dim(J^\perp) = \text{codim} (J^\perp)^\perp \leq \text{codim} J < +\infty$. \square

5. Prove the following theorem:

Theorem 1. Every coalgebra is a sum³ of finite dimensional coalgebras.

Solution. Every coalgebra is the sum of all the subcoalgebra generated by one element which are finite dimensional, as we have just seen it. \square

²This was indeed completely obvious.

³Just a sum, not a direct sum

Problem 4. Let $(\mathcal{C}, \otimes, I, a, l, r)$ a (non-strict) monoidal category. In this problem we want to construct a strict monoidal category \mathcal{D} such that \mathcal{C} and \mathcal{D} are tensor equivalent.

1. We start with a (useful) example. Let \mathcal{V} be the monoidal category whose objects are non-negative integers and whose morphisms from m to n are matrices of size $n \times m$ with coefficient in a field \mathbb{K} , tensor products being given by the sum of integers. Prove that this category is tensor equivalent to \mathbb{K} -vect, the category of finite dimensional vector spaces over \mathbb{K} .
2. The objects of \mathcal{D} are finite sequences (the empty sequence is allowed) of objects of \mathcal{C} . We construct at the same time a (tensor) functor $F : \mathcal{D} \rightarrow \mathcal{C}$ even if \mathcal{D} is not completely defined. If $S = (V_1, V_2, \dots, V_l)$ is an object of \mathcal{D} , we set $F(S) = (\dots((V_1 \otimes V_2) \otimes V_3) \otimes \dots) \otimes V_l$ (what should be $F(\emptyset)$?). Define the hom-spaces of \mathcal{D} and the tensor product on objects of \mathcal{D} , denoted by \star .

Solution. The problem and its solution are derived from Quantum Groups from Christian Kassel. We define $F(\emptyset) = I$ and $\text{Hom}_{\mathcal{D}}(S, S') = \text{Hom}_{\mathcal{C}}(F(S), F(S'))$. The composition and the identity morphisms in \mathcal{D} are given by the composition and the identity morphisms in \mathcal{C} . The (strict) tensor product on objects of \mathcal{D} is given by the concatenation of sequences. The empty sequence being the unit. \square

3. Finish the definition of F and prove that it is fully faithful and essentially surjective (see the script or sheet 1, for the definitions). This proves that F is an equivalence of category which admit $G : \mathcal{D} \rightarrow \mathcal{C}, G(V) = (V)$ as an inverse.

Solution. The definition of F on the hom-spaces is completely trivial since it is really the identity map on each home-space. For this reason, the functor F is clearly fully faithful. It is as-well essentially surjective, since any object V of \mathcal{C} is equal (and hence isomorphic) to the image by F of the sequence (V) of length one with V as the only element of this sequence. Thanks to a theorem we proved earlier this gives that F is an equivalence of category. The proof of this theorem shows that we can indeed take G as prescribed to be an inverse. \square

4. For S and S' two objects of \mathcal{D} , let us define $\phi(S, S') : F(S) \otimes F(S') \rightarrow F(S \star S')$ inductively on the length of S' by:

$$\begin{aligned} \phi(\emptyset, S') &= l_{S'}, & \phi(S, \emptyset) &= r_S, & \phi(S, (V_1)) &= \text{id}_{F(S) \otimes V_1} \quad \text{and} \\ \phi(S, (V_1, \dots, V_{l+1})) &= (\phi(S, (V_1, \dots, V_l)) \otimes \text{id}_{V_{l+1}}) \circ a_{F(S), F((V_1, \dots, V_l)), V_{l+1}}^{-1}. \end{aligned}$$

Prove that if S, S' and S'' are objects of \mathcal{D} , we have the following equality:

$$\phi(S, S' \star S'') \circ (\text{id}_{F(S)} \otimes \phi(S', S'')) \circ a_{F(S), F(S'), F(S'')} = \phi(S, S' \star S'') \circ (\phi(S, S') \otimes \text{id}_{F(S'')}).$$

Solution. The maps ϕ should be thought as “re-parenthesisation”.

Note that this is exactly the “compatibility with the associativity” of definition 2.4.6 in the skript because the associators in \mathcal{D} are identity morphisms. This is done by induction on the length of S'' : If $S'' = \emptyset$, we have:

$$\begin{aligned} \phi(S, S') \circ (\text{id}_S \otimes \phi(S', \emptyset)) \circ a_{F(S), F(S'), I} &= \phi(S, S') \circ (\text{id}_S \otimes r_{F(S')}) \circ a_{F(S), F(S'), I} \\ &= \phi(S, S') \circ r_{F(S) \otimes F(S')} \\ &= r_{F(S) \otimes F(S')} \circ (\phi(S, S') \otimes \text{id}_I) \\ &= \phi(S \star S', \emptyset) \circ (\phi(S, S') \otimes \text{id}_I) \end{aligned}$$

Let know V be an object of the category \mathcal{C} . Let us suppose that the equality holds for the sequences S, S' and S'' .

$$\begin{aligned}
& \phi(S, S' \star S'' \star (V)) \circ (\text{id}_{F(S)} \otimes \phi(S', S'' \star (V))) \circ a_{F(S), F(S'), F(S'' \star (V))} \\
&= (\phi(S, S' \star S'') \otimes \text{id}_V) \circ a_{F(S), F(S' \star S''), V}^{-1} \circ (\text{id}_{F(S)} \otimes (\phi(S', S'') \otimes \text{id}_V)) \\
&\quad \circ (\text{id}_{F(S)} \otimes a_{F(S'), F(S''), V}^{-1}) \circ a_{F(S), F(S'), F(S'') \otimes V} \\
&= (\phi(S, S' \star S'') \otimes \text{id}_V) \circ ((\text{id}_{F(S)} \otimes \phi(S', S'')) \otimes \text{id}_V) \circ a_{F(S), F(S') \otimes F(S''), V}^{-1} \\
&\quad \circ (\text{id}_{F(S)} \otimes a_{F(S'), F(S''), V}^{-1}) \circ a_{F(S), F(S'), F(S'') \otimes V} \\
&= (\phi(S, S' \star S'') \otimes \text{id}_V) \circ ((\text{id}_{F(S)} \otimes \phi(S', S'')) \otimes \text{id}_V) \circ (a_{F(S), F(S'), F(S'')} \otimes \text{id}_V) \circ a_{F(S) \otimes F(S'), F(S''), V}^{-1} \\
&= (\phi(S \star S', S'') \otimes \text{id}_V) \circ ((\phi(S, S') \otimes \text{id}_{F(S'')}) \otimes \text{id}_V) \circ a_{F(S) \otimes F(S'), F(S''), V}^{-1} \\
&= (\phi(S \star S', S'') \otimes \text{id}_V) \circ a_{F(S \star S'), F(S''), V}^{-1} \circ (\phi(S, S') \otimes (\text{id}_{F(S'')} \otimes \text{id}_V)) \\
&= \phi(S \star S', S'' \star (V)) \circ (\phi(S, S') \otimes \text{id}_{F(S'' \star (V))})
\end{aligned}$$

□

5. Define the tensor product of two morphisms in \mathcal{D} and prove that with this structure \mathcal{C}^* is a strict monoidal category.

Solution. In the end we want the following diagrams to commutes:

$$\begin{array}{ccc}
F(S) \otimes F(S') & \xrightarrow{\phi(S, S')} & F(S \star S') \\
\downarrow F(f) \otimes F(g) & & \downarrow F(f \star g) \\
F(T) \otimes F(T') & \xrightarrow{\phi(T, T')} & F(T \star T')
\end{array}$$

This defines the tensor product on \mathcal{D} completely because F is trivial on the hom-space. One verify easily that \star is a functor, and it is strictly associative by construction. □

6. Prove that F and G are tensor functors. Conclude.

Solution. This is to be understood as “prove that F and G can be completed as tensor functors”. The triple (F, id_I, ϕ) is a tensor functor since the question 4 tells us that the required equalities (of the definition 2.4.6 of the script) hold (right and left unit constraint follow from the definition of $\phi(S, \emptyset)$ and $\phi(\emptyset, S)$). The triple $(G, \text{id}, \text{id})$ is as well a tensor functor (the id 's should be widely understood). Finally $FG = \text{id}_{\mathcal{C}}$, and the natural isomorphism $\theta : GF \rightarrow \text{id}_{\mathcal{D}}$ given by $\theta(S) = \text{id}_{F(S)}$ is a tensor natural transformation so that \mathcal{C} and \mathcal{D} are tensor equivalent. □