

# Forcing & the Continuum Hypothesis

## Lecture VII

30 October 2025

TODAY 2pm BEVIN ROOM @  
Cardwell College: FOUNDATIONS  
PhD Advice

We are in the middle of the proof of Gödel's 1938 theorem:



### THEOREM Lectures VI + VII

If  $M \models ZFC$ , then there is  $L \subseteq M$   
transitive containing all ordinals s.t.

1.  $L \models ZFC$
2.  $L \models GCH$
3.  $L$  is minimal in the sense  
that if  $N \subseteq M$  transitive with  
all ordinals and  $N \models ZFC$ ,  
then  $L \subseteq N$ .

We defined  $L$  by  
The definition of  $L$  is  
absolute for s.s. T;

this gives MINIMALITY (3.)

Now working on 1.

$$L_0 := \emptyset$$

$$L_{\alpha+1} := \mathcal{D}(L_\alpha)$$

$$L_\lambda := \bigcup_{\alpha < \lambda} L_\alpha$$

### Lecture VI

### Axioms of ZF

AXIOMS OF ZF	
<b>EXTENSIONAL AXIOM</b>	<b>EXTENSIONALITY</b> $\forall x, y (\forall u (u \in x \leftrightarrow u \in y) \rightarrow x = y)$
	<b>REGULARITY</b> $\forall x (\exists y (x \in y) \rightarrow \exists z (x \in z \wedge \forall u (u \in z \rightarrow u \notin x)))$
	<b>INFINITY</b> $\exists I \exists e (\forall v (v \in I \rightarrow v \in e) \wedge \forall x (x \in I \rightarrow \exists z (z \in I \wedge \forall u (u \in z \rightarrow u \in x))))$
<b>FUNCTIONAL AXIOM</b>	<b>PAIRING</b> $\forall x, y \exists z \forall u (u \in z \leftrightarrow (u = x \vee u = y))$
	<b>UNION</b> $\forall x \exists z \forall u (u \in z \leftrightarrow \exists v (x \in v \wedge u \in v))$
	<b>POWERSET</b> $\forall x \exists z \forall u (u \in z \leftrightarrow \forall v (v \subseteq x \rightarrow v \in z))$
	<b>SEPARATION</b> $\forall \varphi \forall x \exists z \forall u (u \in z \leftrightarrow u \in x \wedge \varphi(u, \vec{p}))$
	<b>REPLACEMENT</b> $\forall \varphi \forall x \forall y (\varphi(x, \vec{p}) \wedge \varphi(y, \vec{p}) \rightarrow x = y) \rightarrow \forall z \exists w (\forall x (x \in z \rightarrow \exists y (y \in w \wedge \varphi(y, \vec{p}))))$

# Lecture VI

## THE STRUCTURAL AXIOMS

Extensionality, Foundation, Infinity

### Lecture II:

The extent of absoluteness for transitive models

Lemma If  $M$  is a transitive set, then  $(M, \in) \models$  Extensionality + Foundation.

Since  $x = \omega$  is absolute and  $\omega$  satisfies the conditions of the Axiom of Infinity, we have that for all  $\alpha > \omega$ ,  $L_\alpha \models$  Infinity.

Now: The Functional Axioms  
Pairing, Union, Powerset, Separation, Replacement

### PAIRING & UNION

Pairing  $\forall x, y \exists p \forall z (z \in p \leftrightarrow z = x \vee z = y)$   
Union  $\forall x \exists u \forall z (z \in u \leftrightarrow \exists y \in x (z \in y))$

The operations  $\{x, y\}$  and  $\bigcup x$  are absolute operations for a s.s.t.

So, I only need to prove:

$\forall x, y \in L (\{x, y\} \in L \text{ and } \bigcup x \in L)$

So, we only need to show  $\{x, y\}, \bigcup x \in \mathcal{D}(L_\alpha)$  for some  $\alpha \rightarrow$  Lecture VII

### Lecture VI

Find  $\{x, y\}$  in  $L$ :

Assume  $x, y \in L$ , so take  $\alpha$  s.t.  $x, y \in L_\alpha$

Consider:

$\varphi(z, x, y) := z = x \vee z = y$

Then  $\mathcal{D}(\varphi, (x, y), L_\alpha) =$

$\{z \in L_\alpha; L_\alpha \models \varphi(z, x, y)\} =$

$\{z \in L_\alpha; L_\alpha \models z = x \vee z = y\} =$

$\{z \in L_\alpha; z = x \vee z = y\} = \{x, y\}$ .

So  $\{x, y\} \in \mathcal{D}(L_\alpha) \subseteq L$ .

Same proof for Union.

# Powerset

$$\forall x \exists p \forall z (z \in p \leftrightarrow z \subseteq x)$$

What is a candidate for the power set of  $x$  in  $L$ ?  
Call it  $p$ .

Clearly (by abs. of  $z \subseteq x$ ),  $p \subseteq \mathcal{P}(x)$ .

Also clearly  $p \subseteq L$ . ↑ POWER SET

So, a candidate could be  $p := \mathcal{P}(x) \cap L$ .

Obviously, if  $p \in L$ , then  $p$  satisfies the conditions of the power set axiom in  $L$ .

Is  $p \in L$ ?

Define  $\Omega := \{ \mathcal{P}_L(z); z \subseteq x, z \in L \}$

By Replacement, there is some  $\mathcal{D}$  s.t.  $\Omega \subseteq \mathcal{D}$ ,

so  $p = \mathcal{P}(x) \cap L \subseteq L \cap \mathcal{D}$ .

if  $\varphi(z, x) := z \subseteq x$ , then

$p = \mathcal{D}(\varphi, x, L \cap \mathcal{D})$ , so  $p \in L \cap \mathcal{D} \subseteq L$ .

Remark. We have no good bound on  $\mathcal{D}$  at the moment since it was given by the Replacement Axiom. We shall need to improve on that later.

# Separation $\varphi$

$$\forall x \forall p \exists s \forall z (z \in s \leftrightarrow z \in x \wedge \varphi(z, p))$$

Which set do we need:

$$\{z \in x; L \models \varphi(z, p)\}$$

no case

Attempt

Assume  $x \in L_\alpha$ .

Take formula  $\psi(z, x, p) := z \in x \wedge \varphi(z, p)$

$$D(\psi, (x, p), L_\alpha) =$$

$$\{z \in L_\alpha; L_\alpha \models z \in x \wedge \varphi(z, p)\}$$

$$= \{z \in x; L_\alpha \models \varphi(z, p)\}$$

Problem: Only works if  $\varphi$  is absolute between  $L_\alpha$  and  $L$ .

LRT to the rescue:

Find  $\beta > \alpha$  s.t.  $\varphi$  is abs. between  $L_\beta$  and  $L$   
 and define  $D(\psi, (x, p), L_\beta) = \{z \in x; L_\beta \models \varphi(z, p)\}$   
 $= \{z \in x; L \models \varphi(z, p)\}$

Replacement  
on ES#2:

Same proof ideas  
as power set  
& separation

COMBINED.

Lövy  
Reflection  
Theorem  
(from Lecture  
IV)

## MAIN THEOREM FOR TODAY

For every TC ZFC finite, there is a transitive model M.T.

The main tool for proving MT is:

### Lövy Reflection Theorem

If  $\alpha \mapsto Z_\alpha$  is a hierarchy and  $\varphi$  is any formula, then  $\forall \beta \exists \gamma > \beta$  s.t.  $\varphi$  is absolute between  $Z_\beta$  and  $Z_\gamma$ .

Def:  $\alpha \mapsto Z_\alpha$  s.t.  
 (a)  $Z_\alpha$  is a transitive set  
 (b)  $\text{Ord} \cap Z_\alpha = \alpha$   
 (c)  $\alpha < \beta \implies Z_\alpha \in Z_\beta$   
 (d)  $\lambda$  limit  $\implies Z_\lambda = \bigcup_{\alpha < \lambda} Z_\alpha$

is called a hierarchy  
 We write  $Z$  for the proper class that is the union of the  $Z_\alpha$ .

### Proof of MT from LRT

Consider the hierarchy

von Neumann

$$V_0 := \emptyset$$

$$V_{\alpha+1} := \mathcal{P}(V_\alpha)$$

This is a hierarchy (ES#1)

$$V_\lambda := \bigcup_{\alpha < \lambda} V_\alpha$$

## Remarks on the Axiom of Choice

So far, AC was not used.

We therefore get: in any model  $M$  of ZF,

there is an inner model  $L \models ZF$ .

(plus MINIMALITY)

However, even if AC is not true in  $M$ , we can provide a wellorder of  $L$ :

Fix some wellorder  $<_\omega$  on  $L_\omega$  of order type  $\omega$ .

Assume that  $<_\alpha$  is a wellorder of  $L_\alpha$ ,  
define (lexicographically) a wellorder of

$\text{Fin} \times L_\alpha$  [par abus de langage write  $<_\alpha$  for this]

and then write for  $x \in L_{\alpha+1}$

$$w(x) := \min_{<_\alpha} \{ (\varphi, p) ; x = D(\varphi, p, L_\alpha) \}$$

Make this into an end-extension of  $<_\alpha$

by

$$x <_{\alpha+1} y \iff \begin{aligned} &x, y \in L_\alpha \wedge x <_\alpha y \\ &\text{OR } x \in L_\alpha \wedge y \in L_{\alpha+1} \setminus L_\alpha \\ &\text{OR } x, y \notin L_\alpha \wedge w(x) < w(y). \end{aligned}$$

Because the orders end-extend,

$< := \bigcup_{\alpha \in \text{Ord}} <_\alpha$  is a wellorder of  $L$ .

However, this was a recursive definition, so it's absolute and so

COROLLARY  $L \models <$  is a global wellorder of  $L$ .  
 $\text{Con}(ZF) \Rightarrow \text{Con}(ZFC)$ .

Now we have established 1. + 3. (ZFC and Minimality) and can proceed to 2. (GCH).

## A closer look at Minimality

If  $T$  is s.s., then for each transitive  $X \models T$ , we have  $\bigcup_{\alpha \in X} L_\alpha \subseteq X$ .

## A converse

### AXIOM OF CONSTRUCTIBILITY

$\forall x \exists \alpha$  ( $\alpha$  is an ordinal  $\wedge x \in L_\alpha$ )

often written as  $V=L$ .

Prop If  $T$  is s.s.,  $X$  transitive s.t.  $X \models T \wedge V=L$ , then  $\bigcup_{\alpha \in X} L_\alpha = X$ .

Proof.  $\subseteq$  is minimality as above;  $\supseteq$  is  $V=L$ .

### THEOREM GÖDEL'S CONDENSATION LEMMA

There is a sentence  $\sigma$  (the condensation sentence) s.t. for all tns  $X$

if  $X \models \sigma$ , then there is  $\alpha$  s.t.  $X = L_\alpha$ .

Proof. Take the s.s.  $T$  (finite) as in Proposition and write  $\sigma := \neg \exists \lambda$  ( $\lambda$  is a limit ordinal  $\wedge$  there is no largest ordinal  $< \lambda$  s.t.  $L_\lambda \models T$ ).

If  $X$  tns,  $X \cap \text{Ord} = \lambda$  for some  $\lambda$ .

So if  $X \models \sigma \implies \lambda$  is a limit ordinal.

By Prop  $X = \bigcup_{\alpha < \lambda} L_\alpha = L_\lambda$  (since  $\lambda$  limit). q.e.d.