

# III

## Third Lecture of FORCING & THE CONTINUUM HYPOTHESIS

1 February 2025

### Transitive models

$M$  transitive  $\Rightarrow (M, \in) \models \text{Extensionality} + \text{Foundation}$

$\Delta_0$ ,  $\Sigma_1$ ,  $\Pi_1$  formulas

Theorem If  $M$  transitive,  $T \subseteq ZFC$ ,  $M \models T$ , then

$\varphi \left\{ \begin{array}{l} \Delta_0^T \\ \Sigma_1^T \\ \Pi_1^T \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} \varphi \text{ absolute for } M \\ \varphi \text{ upwards absolute for } M \\ \varphi \text{ downwards absolute for } M \end{array} \right.$

Operations absolute for  $M$ :  
requires both absoluteness  
of definition and closure

### EXAMPLES What is $\Delta_0$ ?

- |  |                         |
|--|-------------------------|
| 1. $x \in y$   | 10. $z = x \setminus y$ |
| 2. $x = y$   | 11. $z = x \cup \{x\}$  |
| 3. $x \subseteq y \iff \forall w \in x (w \in y)$  | 12. $z$ is transitive   |
| 4. $z = \{x\} \iff x \in z \wedge \forall w \in z (w = x)$   | 13. $z = \bigcup x$     |
| 5. $z = \{x, y\}$  |                         |
| 6. $z = (x, y) = \{\{x, \emptyset\}, \{x, y\}\}$   |                         |
| 7. $z = \emptyset \iff \forall w \in z (w \neq w)$   |                         |
| 8. $z = x \cup y \iff x \subseteq z \wedge y \subseteq z$<br>$\wedge \forall w \in z (w \in x \vee w \in y)$ |                         |
| 9. $z = x \cap y$  |                         |

### More examples

14.  $z$  is an ordered pair  
 $\exists s \in z \exists t \in z \exists x \in s \exists y \in t$   
 $((\forall w \in s (w = x) \wedge \forall v \in t (v = y \vee v = \emptyset)) \wedge \forall w \in z (w = s \vee w = t))$
15.  $z = a \times b$
16.  $z$  is a relation
17.  $z = \text{dom}(x)$
18.  $z = \text{ran}(x)$
19.  $z$  is a function
20.  $z$  is injective
21.  $z$  is surjective
22.  $z$  is bijective

# ORDINALS

" $x$  is an ordinal" :  $\Leftrightarrow x$  is transitive  $\wedge$   
 $(x, \in)$  is a wellorder

We know: being wellfounded is not expressible  
in FOL

[ES#1]

Because all trs model satisfy Foundation, we have that  
if  $M$  is trs

$M \models x$  is transitive  $\wedge (x, \in)$  is linearly ordered  
characterizes ordinals.

clearly  $\Delta_0$

So: being an ordinal is absolute for transitive  
models.

Thus  $M \cap \text{Ord} = \{x \in M; M \models x \text{ is an ordinal}\}$

This is transitive, thus there is  $\alpha \in \text{Ord}$  s.t.

$$\alpha = M \cap \text{Ord}.$$

Also absolute:

" $x$  is successor ordinal"  $\exists y \in x \forall w \in x$   
" $x$  is limit ordinal"  $(w \in y \vee w = y)$

" $x$  is nonzero limit ordinal"

$x = \omega$   $[\forall w \in x$   $w$  is a successor or  $0$   
 $\wedge x$  is limit]

ES#1 :  $x = \omega + 1$   
 $x = \omega \cdot 3$   
 $x = \omega^2 \dots$

# CARDINALS

"x is a cardinal" :  $\iff x$  is an ordinal  $\wedge$   
 $\forall f \forall y \in x \quad \underbrace{f: y \rightarrow x}_{f \text{ is not a surj.}} \implies$

?? not bounded ??  
 bdd

Observe: this is  $\Pi_1$  and therefore downwards abs.

Remarks.

1. We may not want this to be absolute. If it does, we couldn't change cardinal behaviour.
2. We can't obviously bound  $\forall f$  since the natural bound would be  $\{h; h: y \rightarrow x\}$   
 $p(y \times x)$

These, however, are not (yet?!) on our list of absolute concepts.

3. Note that neither 1. nor 2. is an argument since there could be an equivalent formula that is  $\Delta_0$ .

## NON-ABSOLUTENESS

Assume that  $M \models \text{ZFC}$  is transitive & countable.

$$M \cap \text{Ord} = \alpha < \omega_1.$$

However,  $M \models \text{ZFC} \Rightarrow M \models$  there are uncountable cardinals.

Let  $\beta < \alpha$  be s.t.  $M \models \beta$  is the least uncountable cardinal.

But  $\beta$  is a countable ordinal, so not a cardinal.

Consequence All cardinal in  $M$  except for  $\aleph_0$  are going to be fake  
 $\Rightarrow$  "x is a cardinal" can't be absolute.

Note: This also shows that " $x = \aleph(y)$ " cannot be absolute.

Take  $y$  s.t.  $M \models y = \aleph(\omega)$ .

Then  $y \subseteq \aleph(\omega)$ , but is countable since  $y \subseteq M$ .

Thus  $y \neq \aleph(\omega)$ .

Therefore " $x = \aleph(y)$ " is not absolute.

# General proof strategy

Instead of substructures, we will restrict our attention to transitive substructures:

in particular, to  $M$  transitive

s.t.  $M \models ZFC$

$[\forall x x \in M \Rightarrow x \subseteq M$

equivalently

$x \in M \wedge \forall y (x \in y \Rightarrow y \in M)]$

Cohen's proof becomes:

If  $M$  is a countable transitive set s.t.  $M \models ZFC$ ,  
then there is cttb trs set  $N \supseteq M$  s.t.

$N \models ZFC + \neg CH$ .

Q: Is this really solving the original problem?  
i.e.,  $\text{Con}(ZFC) \Rightarrow \text{Con}(ZFC + \neg CH)$ .

It's not obvious that  $\text{Con}(ZFC) \Rightarrow$  there is a ctm  
**COUNTABLE TRANSITIVE MODEL**  
of  $ZFC$ !

A. That's not only not obvious, but false.  
Let's prove that  $\text{Con}(ZFC) \not\Rightarrow$  there is a ctm of  $ZFC$ .

Why? Note that  $\text{Con}(ZFC)$ , or  $\text{Con}(T)$  for any  $T$ ,  
is  $\Delta_0$ . So, it's absolute for trs models.

So if  $M$  is a ctm of  $ZFC$ , then  $\text{Con}(ZFC)$   
is true,

so by absoluteness  $M \models \text{Con}(ZFC)$ .

So  $M \models ZFC + \text{Con}(ZFC)$ . Contradicts Gödel's  
Incompleteness Theorem.

We can get a proof of

$$\text{Con}(\text{ZFC}) \Rightarrow \text{Con}(\text{ZFC} + \neg \text{CH})$$

via a trick (ES#1).

### Cohen Lemma

For every finite  $T \subseteq \text{ZFC}$  there is finite  $T^* \subseteq \text{ZFC}$

s.t. if  $M$  true of  $T^*$ , there there is  
 $N \supseteq M$  s.t.  $N$  true of  $T + \neg \text{CH}$ .

This reduces the problem to:

find true of  $T^*$  for sufficiently  
large finite  $T^* \subseteq \text{ZFC}$ .

For this:

Lévy Reflection Theorem

Def.

We call an assignment

$$\alpha \mapsto Z_\alpha$$

a **HIERARCHY**, if

- (i)  $Z_\alpha$  is a trans set,
- (ii)  $\text{Ord} \cap Z_\alpha = \alpha$ ,
- (iii)  $\alpha < \beta \implies Z_\alpha \subseteq Z_\beta$ , and
- (iv) a limit  $\implies Z_\lambda = \bigcup_{\alpha < \lambda} Z_\alpha$ .

If  $\{Z_\alpha; \alpha \in \text{Ord}\}$  is a hierarchy, we can define

$$Z := \bigcup_{\alpha \in \text{Ord}} Z_\alpha. \quad \text{This is a proper class with } \text{Ord} \subseteq Z.$$

and  $\rho_Z(x) := \min \{ \alpha; x \in Z_\alpha \}$   
a notion of **Z-rank**.

Paradigmatic example: von Neumann hierarchy  
 $\forall \alpha; V_\alpha$  and  $V$  is the entire universe.

### Lévy Reflection Theorem

If  $Z$  is a hierarchy,  $\varphi$  is a formula.  
Then there unboundedly many  $Z_\alpha$   
s.t.  $\varphi$  is absolute between  $Z_\alpha$  and  $Z$ .

**Proposition 2.3.5 (Tarski–Vaught Test)** Suppose that  $M$  is a substructure of  $N$ . Then,  $M$  is an elementary substructure if and only if, for any formula  $\phi(v, \bar{w})$  and  $\bar{a} \in M$ , if there is  $b \in N$  such that  $N \models \phi(b, \bar{a})$ , then there is  $c \in M$  such that  $N \models \phi(c, \bar{a})$ .

TVT

**TVT $_{\Phi}$**  : Let  $M \subseteq N$  and  $\Phi$  be a collection of formulas closed under subformulas. Then TFAE

(i) all formulas in  $\Phi$  are absolute between  $M$  and  $N$

(ii) for all  $\varphi \in \Phi$ , the TV-condition holds:  
 if  $\varphi = \exists x \psi$ , then f.a.  $\vec{y} \in M$   
 if there is  $a \in N$  s.t.  $N \models \psi(a, \vec{y})$ ,  
 then there is  $b \in M$  s.t.  $N \models \psi(b, \vec{y})$ .

Robert Lawson Vaught



Vaught in 1974

**Born** April 4, 1926  
Alhambra, California

**Died** April 2, 2002 (aged 75)  
Berkeley, California

**Nationality** American

**Alma mater** University of California, Berkeley

Alfred Tarski



Tarski in 1968

**Born** Alfred Teitelbaum  
January 14, 1901  
Warsaw, Congress Poland

**Died** October 26, 1983 (aged 82)  
Berkeley, California, US

**Nationality** Polish, American

**Education** University of Warsaw (Ph.D., 1924)