



LARGE CARDINALS

THIRD LECTURE

30 January 2023

RECAP

$I(\kappa) : \iff \kappa \text{ is inaccessible}$

IC : $\iff \exists \kappa I(\kappa)$

1. $I(\kappa) \implies \kappa = \aleph_\kappa$
2. $I(\kappa) \implies V_\kappa \models \text{ZFC}$
3. $I(\kappa) \implies \text{Cons}(\text{ZFC} + \neg \text{IC})$
[If κ is least inaccessible, then $V_\kappa \models \text{ZFC} + \neg \text{IC}$.]

Together:

1. $I(\kappa)$ implies κ is large
2. IC is not provable
3. The largeness of κ features in the proof of 2.

} IC is
a
LARGE
CARDINAL
AXIOM

The consistency strength hierarchy.

Fix a base theory B

In our context, we can always assume $B = ZFC$.

[Proof-theorists would be some weak arithmetic.]

We consider theories T s.t.

$B \subseteq T$, T is computable enumerable and deductively closed
[if $T \vdash \varphi \Rightarrow \varphi \in T$].

On theories like this, we define

$T \leq_{\text{Cons}} S : \Leftrightarrow B \vdash \text{Cons}(S) \Rightarrow \text{Cons}(T)$

$T \equiv_{\text{Cons}} S : \Leftrightarrow T \leq_{\text{Cons}} S \ \& \ S \leq_{\text{Cons}} T$

$T <_{\text{Cons}} S : \Leftrightarrow T \leq_{\text{Cons}} S \ \& \ S \not\leq_{\text{Cons}} T$

EQUICONSISTENCY

If T is c.e., so is $T + \varphi$.

Notation

If T is a theory & φ is a sentence.

$T + \varphi$ is the deductive closure of $T \cup \{\varphi\}$.

Properties of $\leq_{\text{Cons}}$

1. By deductive closure, there is exactly one inconsistent theory, write

$\perp$

for this.

This theory is the maximal element in $\leq_{\text{Cons}}$ and we get that

$$T <_{\text{Cons}} \perp \iff T \text{ is consistent.}$$

2.



David Hilbert
1862-1943

The HILBERTIAN
DREAM
("Finitism"):

There is a base theory $\mathcal{B}$
s.t. φ is true $\iff$
 $\mathcal{B} \vdash \varphi$.

THIS IS NOT TRUE

The Hilbertian dream was shattered
by Gödel's Incompleteness Theorem.

If this existed, then of course all
consistent theories are equiconsistent.

Thus: The interest in $\leq_{\text{Cons}}$ derives from
the INCOMPLETENESS PHENOMENON
which shows that $\leq_{\text{Cons}}$ is undecidable.



3. Con is stronger than just
 "proves more theorems".

Example CH.

Proofs by Gödel & Cohen:

$$\begin{array}{ll} 1938 & \text{Con}(\text{ZFC}) \Rightarrow \text{Con}(\text{ZFC} + \text{CH}) \\ 1962 & \text{Con}(\text{ZFC}) \Rightarrow \text{Con}(\text{ZFC} + \neg \text{CH}) \end{array}$$

Thus $\text{ZFC} \equiv_{\text{Con}} \text{ZFC} + \text{CH} \equiv_{\text{Con}} \text{ZFC} + \neg \text{CH}$

But of course $\text{ZFC} + \text{CH}$ proves more theorems than ZFC .

4. Gödel's 2nd Incompleteness Theorem implies

If $T \neq \perp$, then $T <_{\text{Con}} T + \text{Con}(T)$.

5. Careful There are consistent theories T
 s.t. $T + \text{Con}(T) = \perp$.

Example G2 implies $\text{ZFC} \vdash \text{Con}(\text{ZFC})$
 [Assuming that ZFC is consistent] or in other words $T := \text{ZFC} + \neg \text{Con}(\text{ZFC})$ is consistent.

Since $T \supseteq \text{ZFC}$, $\text{Con}(T) \Rightarrow \text{Con}(\text{ZFC})$,

so $T + \text{Con}(T) \vdash \text{Con}(\text{ZFC})$

Since $T + \text{Con}(T) \vdash \neg \text{Con}(\text{ZFC})$.

$\equiv T$,

So $T + \text{Con}(T) = \perp$.

6. There is a stronger notion of consistency, called ω -consistency

[essentially adding an infinitary proof rule that corresponds to the naive idea of semantics of " $\forall n \in \mathbb{N}$ " and demanding that no contradiction can be derived even with their additional rule]

If T is ω -consistent, then $T + \text{Cons}(T)$ is consistent.

7. Luckily, our proof of $\text{ZFC} + \text{IC} \vdash \text{Cons}(\text{ZFC})$ gives much more:

$$\text{ZFC}_0 := \text{ZFC}$$

$$\text{ZFC}_{i+1} := \text{ZFC}_i + \text{Cons}(\text{ZFC}_i)$$

[We could even add

$$\text{ZFC}_\omega := \text{"for all } i, \text{ZFC}_i \text{"}$$

and continue into the transfinite.]

Want that

$$\text{ZFC}_0 <_{\text{Cons}} \text{ZFC}_1 <_{\text{Cons}} \text{ZFC}_2 <_{\text{Cons}} \dots$$

We proved $\text{ZFC} + \text{IC} \vdash \text{Cons}(\text{ZFC})$
by a very concrete model, viz. V_κ ,
of $\sum \text{ZFC}$.

Since $w \subseteq V_k$ and therefore

$\exists n \in N$ and $\forall n \in N$
are interpreted the same way in V_k & V ,
we get that all formulas whose quantifiers
are restricted to N will get the same
truth value in V_k and V . $V \models \phi \Leftrightarrow V \models "V_k \models \phi"$

(ARITHMETICAL FORMULAS)

In other words, arithmetical formulas
are absolute between V_k & V .

IMPORTANT For any T , $\text{Cons}(T)$ is
an arithmetical formula.

<u>BOOTSTRAP</u> :	Suppose $V \models IC$	1
$\Rightarrow$	$V \models "V_k \models ZFC"$	2
$\Rightarrow$	$V \models \text{Cons}(ZFC)$	3
Absoluteness	$V \models "V_k \models \text{Cons}(ZFC)"$	4
$2+4$	$V \models "V_k \models ZFC_1"$	5
$\Rightarrow$	$V \models \text{Cons}(ZFC_1)$	6
Absoluteness	$V \models "V_k \models \text{Cons}(ZFC_1)"$	7
$5+7$	$V \models "V_k \models ZFC_2"$	8

AND SO ON.

To summarize:

$ZFC + IC$ proves

$$ZFC_0 <_{\text{Cons}} ZFC_1 <_{\text{Cons}} \dots <_{\text{Cons}} ZFC_i <_{\text{Cons}} \dots <_{\text{Cons}} ZFC + IC$$

Natural question

Is this about inaccessibility or rather about $V_\kappa \models ZFC$?

Definition

An ordinal α is called WORLDLY if $V_\alpha \models ZFC$.

Are "inaccessible" and "worldly" the same thing?

Answer will be (Lectures IV & V):

NO!

In particular, we'll show that if $IC(\kappa) \Rightarrow$ there are many worldly cardinals below κ .

The difference between INACCESSIBLE and WORLDLY cardinals comes down to the difference between

$V_\kappa \models \text{Replacement}$
and

V_κ satisfies SOR
mentioned in Lecture II.