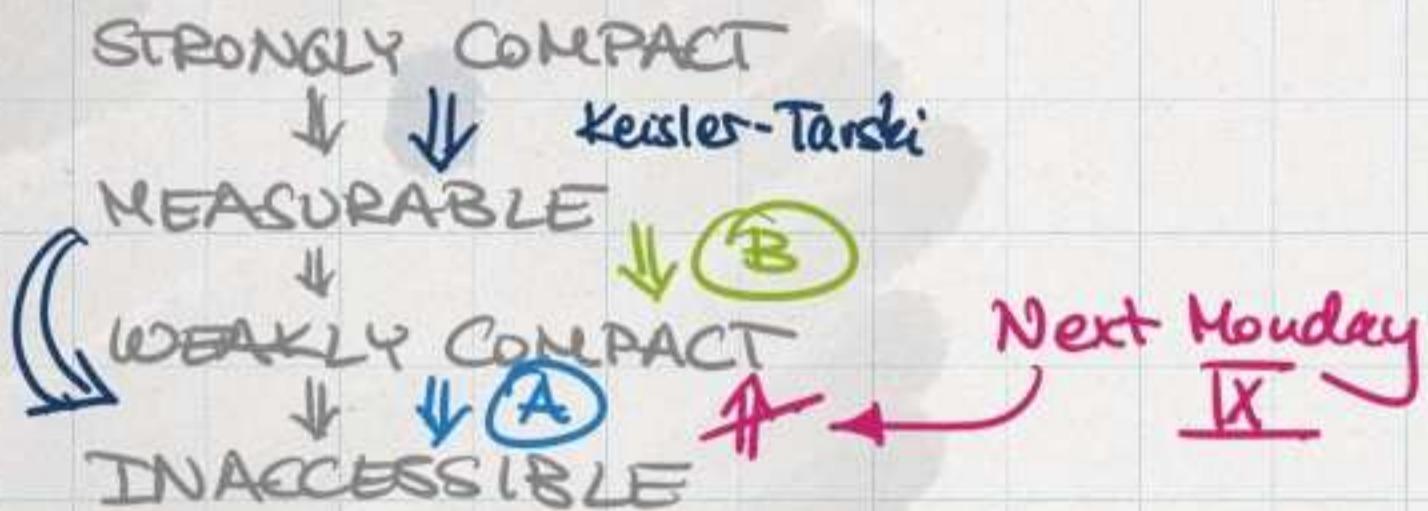


# Large Cardinals

## Eighth Lecture

16 February  
2022



Theorem A weakly compact  $\Rightarrow$  inaccessible

Proof. We already know "regular".  
So need to show strong limit.

Towards a contradiction, assume  $2^\lambda \geq \kappa$ .

for some  $\lambda < \kappa$ .

Like language with constant symbols  
 $c_\alpha, c_\alpha^0, c_\alpha^1$  for  $\alpha < \lambda$ .

[This is OK for weak compactness.]

$$\Phi^* \vdash \left[ \begin{array}{l} \bigwedge_{\alpha < \lambda} c_\alpha^0 \neq c_\alpha^1 \\ \bigwedge_{\alpha < \lambda} c_\alpha = c_\alpha^0 \vee c_\alpha = c_\alpha^1 \end{array} \right]$$

-

$$q_f := \bigvee_{\alpha < \lambda} c_\alpha \neq c_\alpha^{f(\alpha)}$$

for  $f: \lambda \rightarrow 2$

If  $\mathcal{M} \models \Phi$ , we can define a function

$$\begin{aligned} g: \lambda &\longrightarrow 2 \\ \text{by } g(\alpha) &:= \begin{cases} 0 & \text{if } \mathcal{M} \models c_\alpha = c_\alpha^0 \\ 1 & \text{if } \mathcal{M} \models c_\alpha = c_\alpha^1. \end{cases} \end{aligned}$$

$\Phi^*$  is inconsistent since the additional formula says that no such function can exist.

To conclude our contradiction, we only need to show that  $\Phi^*$  is  $\kappa$ -satisfiable.

$$\Phi_0 \subseteq \Phi^* \text{ s.t. } |\Phi_0| < \kappa \leq 2^\lambda.$$

Therefore, there is some  $f: \lambda \rightarrow 2$  s.t.

$q_f \notin \Phi_0$ . But now interpret  $c_\alpha$  by  $f$ , and that will satisfy  $\Phi_0$ . q.e.d.

## The method of ultraproducts

If  $\mathcal{U}$  is an ultrafilter on  $\kappa$  and let  $(M_\alpha ; \alpha < \kappa)$  be a sequence of first-order structures.

Write  $\vec{M}$  for this sequence.

Consider choice functions for  $\vec{M}$ :

$$f : \kappa \longrightarrow \bigcup_{\alpha < \kappa} M_\alpha$$

$$\text{s.t. } f(\alpha) \in M_\alpha$$

Say  $f \sim_U g \iff$

$$\{\alpha ; f(\alpha) = g(\alpha)\} \in \mathcal{U}$$

This is an equivalence relation on the set of choice functions.

$\mathcal{U} / \sim_U (\vec{M}, \mathcal{U})$  is the quotient of the set of choice functions modulo

$$\sim_U.$$

We write  $[f]$  for  $\{g ; f \sim_U g\}$ .

**ULTRA-  
PRODUCT**

Check that the obvious definitions

[Example.] If  $R$  is a binary  
relation symbol

$[f]R[g] \iff$

$$\{ \alpha < \kappa ; M_\alpha \models f(\alpha) R g(\alpha) \} \in U.$$

are welldefined and give a structure  
of the right type.

Theorem (Łoś's Theorem)

If  $\varphi$  is a formula in  $n$  variables,  
then

$$V \models (\vec{m}, U) \models \varphi([f_1], \dots, [f_n])$$

$\iff$

$$\{ \alpha < \kappa ; M_\alpha \models \varphi(f_1(\alpha), \dots, f_n(\alpha)) \} \in U.$$

[Essentially: induction on complexity of  $\varphi$ .]

Observation 1 If  $\varphi$  is a sentence, Łoś gives us

$$UH(\vec{M}, U) \models \varphi \iff \{\alpha < \kappa; M_\alpha \models \varphi\} \in U.$$

So if all  $M_\alpha \models \varphi$ , then so is  $UH(\vec{M}, U)$ .

Observation 2 If  $\vec{m}$  is constant, i.e., if there is an  $M$  s.t.  $M_\alpha = \vec{m}$  for all  $\alpha$ , then we call  $UH(\vec{M}, U)$  an ultrapower and write  $UH(M, U)$ .

In this case, we can define the following map

$$\begin{array}{ccc} m & \xrightarrow{\quad} & [c_m] \\ \uparrow & & \uparrow \\ M & & c_m: \alpha \rightarrow M \end{array}$$

constant function

This is an embedding from  $M$  into  $UH(M, U)$ .

### Observation 3

This map

$$j_U: m \mapsto [c_m]$$

called the ultrapower embedding is  
an elementary embedding (by Łoś):

[Need to show:

$$M \models \varphi(m_1, \dots, m_n) \iff U \Vdash (M, U) \models$$

$$m_1, \dots, m_n \in M$$

$$\varphi(j_U(m_1), \dots, j_U(m_n))$$

$$\iff U \Vdash (M, U) \models \varphi([c_{m_1}], \dots, [c_{m_n}])$$

$$\xleftarrow{\text{Łoś}} \left\{ \alpha < \kappa; M \models \varphi(c_{m_1}(\alpha), \dots, c_{m_n}(\alpha)) \right\} \in U$$

$$\iff \left\{ \alpha < \kappa; \underbrace{M \models \varphi(m_1, \dots, m_n)} \right\} \in U$$

This is independent of  $\alpha$ ,  
so either true or false.

If true:  $\{ \dots \} = \kappa \in U$

If false:  $\{ \dots \} = \emptyset \notin U$ .

$$\iff M \models \varphi(m_1, \dots, m_n)$$

We can think of  $M$  as an elementary  
substructure of  $(U, U)$  by identifying  
 $M$  with  $\{j_U(m); m \in M\}$ .

Theorem (B) measurable  $\Rightarrow$  w.c.

Proof. Suppose  $\kappa$  is measurable. Then we know that  $\kappa$  is inaccessible.

[Lemma.] If  $L$  is an L $\kappa\kappa$ -language with at most  $\kappa$  many non-logical symbols, then  $|L| \leq \kappa$ .

Proof. Clearly, at most  $\kappa$  many atomic formulas.

Inductively, we only need to show that the construction steps for L $\kappa\kappa$ -languages preserve "card.  $\leq \kappa$ ", i.e., if  $X$  is a set of formulas with  $|X| \leq \kappa$ , then the closure of  $X$  under the L $\kappa\kappa$ -constructions has card.  $\leq \kappa$ .

In this case, closing under  $\bigwedge_{\alpha < \lambda} \varphi_\alpha$  will give cardinality  $\kappa^{<\kappa} = \bigcup_{\lambda < \kappa} \kappa^\lambda$ . Since  $\kappa$  is inaccessible,  $\kappa^{<\kappa} = \kappa$ .

Similarly,  $\exists^{\lambda} \bar{x} \varphi$  will give rise to  $\kappa^{<\kappa}$  many formulas, so  $\leq \kappa$ .  $\square$

See also  
Jack  
Thm S2D  
p. 57

If  $\Phi$  is any set of formulas in an  $\mathcal{L}_\kappa$ -language with  $\leq \kappa$  nonlogical symbols, we can write it as.

$$\Phi = \{\varphi_\alpha; \alpha < \kappa\}$$

by our lemma.

Need to show: if  $\Phi$  is  $\kappa$ -satisfiable, then it's satisfiable.

Define  $\bar{\Phi}_\lambda := \{\varphi_\alpha; \alpha < \lambda\}$ . By assumption,  $\bar{\Phi}_\lambda$  is satisfiable. Let

$$M_\lambda \models \bar{\Phi}_\lambda.$$

$$\text{Define } M := \text{Ult}(\vec{M}, U)$$

Claim  $M \models \Phi$ .

Let  $\varphi_\alpha$  be arbitrary.

$$\{\lambda < \kappa; M_\lambda \models \varphi_\alpha\}$$

$$\supseteq \{\lambda < \kappa; \alpha < \lambda\} \in U.$$

$$\text{So } \{\lambda < \kappa; M_\lambda \models \varphi_\alpha\} \in U$$

$$\xleftarrow{\text{L.o.s}} \text{Ult}(\vec{M}, U) \models \varphi_\alpha.$$

REMINDER

If  $|A| < \kappa$ , then

$A \notin U$  by  
 $\kappa$ -completeness +  
non-triviality.

q.e.d.

a nontrivial  
 $\kappa$ -complete  
uf. on  $\kappa$   
[exists by  
measurability]

Next goal

inaccessible  $\not\Rightarrow$  w.c.

Technique : REFLECTION and  
the Kiesler Extension Property.



If  $K$  is a weakly compact,  
 $V_K$  is a transitive set

$$X \not\cong V_K$$

s.t.

$$(V_K, \in) \preccurlyeq (X, \in).$$

[Compare the situation of worldly cardinals.]