

# INFINITE GAMES

## Lecture VI

3 February 2021

FROM  
LECTURE  
V

### Theorem (Gale-Stewart)

All closed sets  $A \subseteq M^\omega$  are quasidetermined.

We are in the middle of the proof:

$l$   $A = [T]$ ;  $l(s) := \underline{II}$  if  $s \notin T$   
PARTIAL LABELLING  $l: M^{<\omega} \rightarrow \{\underline{I}, \underline{II}\}$

Given  $l$ ,  $l^+ \supseteq l$  defined by

if  $l(s)$  is even and for all  $m \in M$   $l(sm) = \underline{II}$ , then

$$l^+(s) := \underline{II}$$

$l^+$   $\otimes$  if  $l(s)$  is odd and for some  $m \in M$   $l(sm) = \underline{II}$ , then

$$l^+(s) := \underline{II}.$$

RECURSIVE  
DEFINITION

$$l_0 := l; \quad l_{\alpha+1} := l_\alpha^+; \quad l_\lambda := \bigcup_{\delta < \lambda} l_\delta$$

FIXED PT  $l_\alpha$  By Replacement, there is a fixed pt  $l_{\alpha+1} = l_\alpha$ .

$$\hat{l}: M^{<\omega} \rightarrow \{\underline{I}, \underline{II}\}; \quad \hat{l}(s) := \begin{cases} \underline{II} & \text{if } l_\alpha(s) = \underline{II} \\ \underline{I} & \text{o/w.} \end{cases}$$

Claim

if  $\hat{l}(\emptyset) = \underline{I}$ , then player  $\underline{I}$  has a winning quasistrategy.

if  $\hat{l}(\emptyset) = \underline{II}$ , then player  $\underline{II}$  has a w. q. s.

FROM  
LECTURE  
V

This claim proves the G-S Theorem.

Subclaim 1 If  $\hat{\ell}(\emptyset) = \underline{I}/\underline{\Pi}$ , define

$$Q_{\underline{I}} := \{s \in \omega^{<\omega}; \hat{\ell}(s) = \underline{I}\}$$

$$Q_{\underline{\Pi}} := \{s \in \omega^{<\omega}; \hat{\ell}(s) = \underline{\Pi}\}.$$

Then  $Q_{\underline{I}}/Q_{\underline{\Pi}}$  is a  $\underline{I}/\underline{\Pi}$ -game tree.

[Need to show:

① if  $\hat{\ell}(s) = \underline{I}$  &  $lh(s)$  is even then there is  $m$  s.t.  $\hat{\ell}(sm) = \underline{I}$

② if  $\hat{\ell}(s) = \underline{I}$  &  $lh(s)$  is odd then for all  $m$   $\hat{\ell}(sm) = \underline{I}$

③ if  $\hat{\ell}(s) = \underline{\Pi}$  &  $lh(s)$  is even then for all  $m$   $\hat{\ell}(sm) = \underline{\Pi}$

④ if  $\hat{\ell}(s) = \underline{\Pi}$  &  $lh(s)$  is odd then there is  $m$  s.t.  $\hat{\ell}(sm) = \underline{\Pi}$

We had  $l_\alpha = l_{\alpha+1} = (l_\alpha)^+$ , so ③ & ④

follow immediately from \*

Similarly, if

$lh(s)$  is even  $\hat{\ell}(s) = \underline{I} \implies s \notin \text{dom}(l_\alpha)$

$\implies$  there is  $m$   $sm \notin \text{dom}(l_\alpha)$

$\implies \hat{\ell}(sm) = \underline{I}$ .

And ② is similar]

Subclaim 2 If  $\hat{l}(\emptyset) = \underline{I}$ , then  $Q_I$  is a w.q.s. for  $\underline{I}$ .

[Need to show  $[Q_I] \subseteq A$ . So fix some  $x \in [Q_I]$ . So  $\forall n$

$$\begin{aligned} x \upharpoonright n &\in Q_I \\ \Rightarrow \hat{l}(x \upharpoonright n) &= \underline{I} \\ \Rightarrow \hat{l}(x \upharpoonright n) &\neq \underline{II} \\ \Rightarrow l(x \upharpoonright n) &\neq \underline{II} \\ \Rightarrow x \upharpoonright n &\in T \end{aligned}$$

$l(s) = \underline{I}$   
 $\iff$   
 $s \in \underline{I}$

So  $x \in A$ .]

Remark What we did so far, didn't really need trees / closed sets. Let

$l: M^{<\omega} \rightarrow \{\underline{I}\}$  be any partial labelling. Do the transfinite recursion  $l_{\alpha+1} := (l_\alpha)^\dagger$ ,  $l_0 := l, \dots$

Find fixed pt  $l_\alpha$ ; define  $\hat{l}$ ; define  $Q_I$ . Then  $Q_I$  is a q.s. that avoids all  $s \in \text{dom}(l)$ .

For player  $\text{II}$ , it's not always the case that  $Q_{\text{II}}$  is winning

[cf. Example (12) on ES #1]

since you could stay on labels  $\text{II}$  without ever leaving the tree.

So, we need to work a bit harder.

Remember the age function:

$$\text{age}: Q_{\text{II}} \rightarrow \alpha + 1$$

$$\text{age}(s) := \text{the least } \beta \text{ s.t. } s \in \text{down}(l_\beta)$$

This means:

$$\text{age}(s) = 0 \iff s \in \text{down}(l_0) = \text{down}(l)$$

$$\iff s \notin T.$$

Subclaim 3 If  $\hat{\ell}(s) = \underline{\Pi}$  and  $\ell(s)$  is even then  $\text{age}(s) = 0$  or for all  $m$   $\text{age}(sm) < \text{age}(s)$ .  
 If  $\hat{\ell}(s) = \underline{\Pi}$  and  $\ell(s)$  is odd then  $\text{age}(s) = 0$  or there is  $m$   $\text{age}(sm) < \text{age}(s)$ .

[Follows directly from  $\ast$ .]

Define  $\hat{Q}_{\underline{\Pi}} \subseteq Q_{\underline{\Pi}}$  by

$\emptyset \in \hat{Q}_{\underline{\Pi}}$  and

$sm \in \hat{Q}_{\underline{\Pi}} \iff \text{age}(sm) = 0$   
 or

$sm \in Q_{\underline{\Pi}}$  and

$\text{age}(sm) < \text{age}(s)$

INFORMALLY

Playing by  $\hat{Q}_{\underline{\Pi}}$  means:

"play into positions labelled  $\underline{\Pi}$  reducing the age if you can".

If  $\hat{L}(\emptyset) = \Pi$ , then by Subclaim 3,  
 $Q_{\Pi}$  is a q.s.

Subclaim 4 If  $\hat{L}(\emptyset) = \Pi$ , then  $\hat{Q}_{\Pi}$   
 is winning for  $\Pi$ .

$$[\hat{Q}_{\Pi}] \cap A = \emptyset$$

[Suppose  $x \in [\hat{Q}_{\Pi}]$ . Consider

$$a_n := \text{age}(x \upharpoonright n)$$

This is a <sup>strictly</sup> decreasing seq. of ordinals  
 until it hits 0 by construction  
 of  $\hat{Q}_{\Pi}$ . Since no infinite <sup>strictly</sup> dec.

seq. of ordinals exist, we find  $k$   
 s.t.  $a_k = \text{age}(x \upharpoonright k) = 0$ .

$$\implies x \upharpoonright k \notin T$$

$$\implies x \notin A$$

qed (Claim)  
 qed (AST)

Remark If  $M$  is wellorderable (e.g.,  $M = \mathbb{N}$ ), then GST says:

Every closed subset of  $M^\omega$  is determined.

Also: Every complement of a closed set is determined.

[Follows directly from the proof.]

---

## End of Lecture IV

MOTTO / HOPE : If a set is NICE or SIMPLE, then it is determined.

If NICE = CLOSED, then this is just GST.

Next goal : Define a topology on  $M^\omega$ .  
Let's focus on the case  $M = \mathbb{N}$  or even  $M = 2 = \{0, 1\}$ .

Definition  $x, y \in \omega^\omega$

$$d(x, y) := \begin{cases} 0 & \text{if } x = y \\ 2^{-n} & \text{if } x \upharpoonright n = y \upharpoonright n \\ & \text{and } x(n) \neq y(n) \end{cases}$$

This is a metric on  $\omega^\omega$ .

What are the open balls for this metric

$$B_\varepsilon(x) = \{y \in \omega^\omega; d(x,y) < \varepsilon\}$$

$$\varepsilon = 2^{-n}$$

$$B_{2^{-n}}(x) = \{y \in \omega^\omega; d(x,y) < 2^{-n}\}$$

$$= \{y; y \upharpoonright_{n+1} = x \upharpoonright_{n+1}\}$$

OPEN BALLS ARE DETERMINED  
BY FINITE SEQ.

If  $s \in \omega^{<\omega}$ , write

$$[s] := \{x \in \omega^\omega; s \subseteq x\}$$

$$\text{So } B_{2^{-n}}(x) = [x \upharpoonright_{n+1}].$$

Thus the topology of the metric space  
is generated by the basic open sets  
 $\{[s]_{\omega^\omega}; s \in \omega^{<\omega}\}$ .

This space is called **BAIRE** space;  
if we restrict to  $2^\omega$ , then it  
is called **CANTOR** space.



If you think of

$$\omega^\omega \text{ as } \prod_{i \in \omega} X_i \quad \text{with } X_i = \omega$$

$$\text{and } 2^\omega \text{ as } \prod_{i \in \omega} Y_i \quad \text{with } Y_i = 2 = \{0, 1\}$$

then Baire space is just the product topology on  $\prod_{i \in \omega} X_i$  with the discrete

topology on  $\omega$ ;

and Cantor space is the product topology on  $\prod_{i \in \omega} Y_i$  with the discrete

topology on  $2$ .

Tychonoff implies that Cantor space is compact, but Baire space is not.

Easy to see the latter:

$$\omega^\omega = \bigcup_{n \in \omega} \left[ \langle n \rangle \right] \\ = \{x; x(0) = n\}$$

Next time:

Show that

$$A = [T] \subseteq \omega^\omega$$



$A$  is closed in  
Baire space.