

# SET THEORY 2021-12-13

- THE LEO IS THE ULTIMATE ARBITER
- SUMMARY OF WHAT WE HAVE ACHIEVED THUS FAR.
  - TO DEFEAT A POTENTIAL PROOF OF A STATEMENT  $\psi$ , E.G., CH, TCH, SUSLIN'S HYPOTHESIS - TAKE FINITELY MANY INSTANCES OF) AXIOMS SAY  $\varphi_1, \dots, \varphi_n$  AND BUILD A COUNTABLE STRUCTURE  $N$  SUCH THAT  $\varphi_1^N \wedge \dots \wedge \varphi_n^N \wedge (\neg\psi)^N$  HOLDS.
  - GIVEN FINITELY MANY AXIOMS  $\varphi_1, \dots, \varphi_n$  WE CAN FIND A COUNTABLE TRANSITIVE  $M$  SUCH THAT  $\varphi_1^M \wedge \dots \wedge \varphi_n^M$  HOLDS.
  - FOR CONVENIENCE WE ASSUME  $M$  SATISFIES ALL OF ZFC. [AFTERWARDS YOU CAN INSPECT THE PROOF AND SEE WHICH INSTANCES OF THE AXIOMS YOU USED (AND PRETEND YOU ASSUMED THAT  $M$  IS SATISFIED ONLY THOSE).
  - WE CAN EXTEND  $M$  USING A SINGLE SET  $G$  TO OBTAIN  $N = M[G]$  SUCH THAT  $(\neg\psi)^N$  HOLDS AND  $N$  ALSO SATISFIES ALL OF ZFC.

- HOW :- TAKE A PARTIAL ORDER  $IP$  IN  $M$
- $G$  IS TO BE A FILTER ON  $IP$ 
    - $p, q \in G \rightarrow \exists r \in G (r \leq p \wedge r \leq q)$
    - $p \in G \wedge p \leq q \rightarrow q \in G$
  - $G \cap D \neq \emptyset$  WHENEVER  $D \in M$  IS DENSE IN  $IP$   
 DENSE:  $(\forall p \in IP) (\exists q \in D) (q \leq p)$
  - SUCH  $G$  EXIST BECAUSE  $M$  IS COUNTABLE FOR EVERY  $p \in IP$  THERE IS SUCH A  $G$ .
  - SUCH  $G$  ARE CALLED  $M$ -GENERIC FILTERS ON  $IP$  [SOME AUTHORS: " $IP$ -GENERIC IN  $M$ "]

- USEFUL NOTIONS:
- $p \perp q$  (INCOMPATIBLE) NO  $r$  WITH  $r \leq p \wedge r \leq q$
  - $p \parallel q$  (COMPATIBLE) SOME  $r$  WITH  $r \leq p \wedge r \leq q$
  - ANTICHAIN  $A \subseteq IP$ :  
 $p, q \in A \wedge p \neq q \rightarrow p \perp q$
  - ZORN: MAXIMAL ANTICHAINS EXIST.

USEFUL:  $G$  IS AN  $M$ -GENERIC FILTER ON  $\mathbb{P}$   
IFF  $-G$  IS A FILTER

-  $G \cap A \neq \emptyset$  FOR EVERY MAXIMAL  
ANTICHAIN  $A$  IN  $\mathbb{P}$  THAT IS IN  $M$   
(NOTE  $|G \cap A| = 1$  IN THAT CASE)

FORCING: - BUILDING  $M[G]$  OUT OF  $M \cup \{G\}$   
- CLOSE RELATION BETWEEN  
TRUTH IN  $M[G]$  AND TRUTH IN  $M$ .

-  $\mathbb{P}$ -NAMES:

$\tau$  IS A  $\mathbb{P}$ -NAME

IFF  $\tau$  IS A RELATION [SET OF ORDERED  
PAIRS]

SUCH THAT

IF  $\langle \pi, p \rangle \in \tau$  THEN - IT IS A  $\mathbb{P}$ -NAME  
-  $p \in \mathbb{P}$ .

$M^{\mathbb{P}}$  IS THE COLLECTION OF  $\mathbb{P}$ -NAMES  
THAT ARE IN  $M$

- FOR US IT IS A SET

- FOR  $M$  IT IS A PROPER CLASS

-  $VAL(\tau, G) = \{ VAL(\pi, G) : (\exists p \in G)(\langle \pi, p \rangle \in \tau) \}$   
ABBREVIATION  $\tau_G$  (OTHER NOTATION  $\tau[G]$ ).

- WE DEFINE  $\Vdash$  (FROM THE OUTSIDE):

$p \Vdash \varphi(\tau_1, \dots, \tau_n)$  IFF  $\varphi(\tau_{1G}, \dots, \tau_{nG})$   <sup>$M[G]$</sup>   
FOR EVERY  $M$ -GENERIC  
FILTER  $G$  ON  $\mathbb{P}$  WITH  $p \in G$

[NOTE: IN ALL NON-TRIVIAL CASES WE HAVE  $G \neq \emptyset$   
IF NON-TRIVIAL:  $(\forall p \in \mathbb{P})(\exists q, r \in \mathbb{P})(q, r \leq p \wedge q \perp r)$   
FOR IN THAT CASE  $\mathbb{P} \setminus G$  IS DENSE]

- WE ALSO DEFINE  $\Vdash^*$  (AND RELATIVISED TO  $M$ )

$p \Vdash^* \varphi(\tau_1, \dots, \tau_n)$

BY RECURSION ON COMPLEXITY OF  $\varphi$ .

FIRST  $\tau_1 = \tau_2$

THEN  $\tau_1 \in \tau_2$

THEN  $\neg \varphi$  AND  $\varphi \wedge \gamma$

THEN  $(\exists \sigma)(\varphi(\sigma, \tau_1, \dots, \tau_n)) \dots (\exists \sigma \in M^{\mathbb{P}}) \dots$

### MAIN TECHNICAL RESULT

- ① IF  $p \in G$  AND  $(p \Vdash^* \varphi(z_1, \dots, z_n))^M$   
THEN  $(\varphi(\text{VAL}(z_1, G), \dots, \text{VAL}(z_n, G)))^{M[G]}$
- ② IF  $(\varphi(\text{VAL}(z_1, G), \dots, \text{VAL}(z_n, G)))^{M[G]}$   
THEN  $(\exists p \in G) (p \Vdash^* \varphi(z_1, \dots, z_n))^M$ .

BETTER VERSION OF LAST WEEK'S NOTES:

JUST FOR  $\tau_1 = \tau_2$ :  $p \Vdash^* \tau_1 = \tau_2$  MEANS

- (i) FOR ALL  $\langle \pi_1, s_1 \rangle \in \tau_1$  THE  $D(\pi_1, s_1, p)$   
 $\cdot D \{ q \leq p : q \leq s_1 \rightarrow (\exists \langle \pi_2, s_2 \rangle \in \tau_2) (q \leq s_2 \wedge q \Vdash^* \pi_1 = \pi_2) \}$
- (ii) FOR ALL  $\langle \pi_2, s_2 \rangle \in \tau_2$   
 $\{ q \leq p : q \leq s_2 \rightarrow (\exists \langle \pi_1, s_1 \rangle \in \tau_1) (q \leq s_1 \wedge q \Vdash^* \pi_1 = \pi_2) \}$   
 ARE DENSE BELOW  $p$ .

THEN (i) GIVES:  $\tau_{1,G} \subseteq \tau_{2,G}$  AND (ii) GIVES  $\tau_{2,G} \subseteq \tau_{1,G}$

- ① SO TAKE  $\alpha \in \tau_{1,G}$ : THERE IS  $\langle \pi_1, s_1 \rangle \in \tau_1$   
 WITH  $s_1 \in G$  AND  $\alpha = \text{VAL}(\pi_1, G) = \pi_{1,G}$   
 THERE IS  $\tau \in G$  WITH  $\tau \leq p$  AND  $\tau \leq s_1$   
 THEN  $\tau \Vdash^* \tau_1 = \tau_2$  AND SO THERE IS  $q \in G \cap D(\pi_1, s_1, \tau)$   
 BUT  $q \leq \tau \leq s_1$  SO  $(\exists \langle \pi_2, s_2 \rangle \in \tau_2) (q \leq s_2 \wedge q \Vdash^* \pi_1 = \pi_2)$   
 WE GET  $s_2 \in G$ , SO  $\pi_{2,G} \in \tau_{2,G}$   
 AND  $q \Vdash^* \pi_1 = \pi_2$   
 BY INDUCTION  $\pi_{1,G} = \pi_{2,G}$  SO  $\alpha \in \tau_{2,G}$ .

- ② ASSUME  $\tau_{1,G} = \tau_{2,G}$   
 LET  $D$  BE THE UNION OF  $\{ \tau : \tau \Vdash^* \tau_1 = \tau_2 \}$ ,  
 $D_1 = \{ \tau : (\exists \langle \pi_1, s_1 \rangle \in \tau_1) (\tau \leq s_1 \wedge (\forall \langle \pi_2, s_2 \rangle \in \tau_2) (\forall q \in P) (q \leq s_2 \wedge q \Vdash^* \pi_1 = \pi_2) \rightarrow q \perp \tau) \}$ , AND  
 $D_2 = \{ \tau : (\exists \langle \pi_2, s_2 \rangle \in \tau_2) (\tau \leq s_2 \wedge (\forall \langle \pi_1, s_1 \rangle \in \tau_1) (\forall q \in P) (q \leq s_1 \wedge q \Vdash^* \pi_1 = \pi_2) \rightarrow q \perp \tau) \}$ .

WE WANT  $G$  TO INTERSECT  $\{ \tau : \tau \Vdash^* \tau_1 = \tau_2 \}$   
 IF  $\tau \in G \cap D_1$ , THEN WE GET  $\langle \pi_1, s_1 \rangle \in \tau_1$  WITH  $\tau \leq s_1$  AND  $\tau \in G$   
 AND SO  $\pi_{1,G} \in \tau_{1,G}$  AND ALSO  
 IF  $\alpha \in \tau_{2,G}$  THEN  $\alpha = \pi_{2,G}$  FOR SOME  $\langle \pi_2, s_2 \rangle \in \tau_2$   
 WITH  $s_2 \in G$

IF  $\pi_{1,G} = \pi_{2,G}$  THEN BY INDUCTION THERE  
 IS  $p \in G$  SUCH THAT  $p \Vdash^* \pi_1 = \pi_2$   
 BUT TAKE  $q \in G$  WITH  $q \leq p, \tau_1, s_2$   
 THEN  $q \leq s_2$  AND  $q \perp \tau$  SO  $q \Vdash^* \pi_1 = \pi_2$   
 BUT ALSO  $q \leq p$  SO  $q \Vdash^* \pi_1 = \pi_2$   
CONTRADICTION

SIMILARLY  $D_a \cup D_c \cup D_u = \emptyset$

SO WE MUST SHOW  $D_a \cup D_c \cup D_u$  IS DENSE

LET  $p \in P$  AND ASSUME  $p \notin D_a$  SO  $p \Vdash \dot{z}_1 = \dot{z}_2$

SAY (C) FAILS; WE FIND  $q \leq p$  THAT IS IN  $D_c$

TAKE  $\langle \pi_1, s_1 \rangle \in \dot{z}_1$  AND USE "NOT DENSE BELOW  $p$ " TO FIND  $q \leq p$  SUCH THAT NO  $q \leq r$  WORK AGAIN (C)

SO  $(\forall q \leq r) (q \leq s_1 \wedge (\forall \langle \pi_2, s_2 \rangle \in \dot{z}_2) (\neg (q \leq s_2 \wedge q \Vdash \pi_1 = \pi_2)))$

IN PARTICULAR  $r \leq s_2$

NOW IF  $\langle \pi_2, s_2 \rangle \in \dot{z}_2$  AND  $q \leq s_2$  AND  $q \Vdash \pi_1 = \pi_2$

THEN  $q \perp r$

FOR IF  $s \leq q, r$  THEN BOTH ON  $s \Vdash \pi_1 = \pi_2$

$(s \leq s_2 \wedge s \Vdash \pi_1 = \pi_2)$  BECAUSE  $s \leq q$  AND

$\neg (s \leq s_2 \wedge s \Vdash \pi_1 = \pi_2)$  BECAUSE  $s \leq r$

SO  $r \in D_c$

### WHAT WE USE IN PRACTICE

① FOR ALL  $p \in P$

$p \Vdash \varphi(\dot{z}_1, \dots, \dot{z}_n)$  IFF  $(p \Vdash \varphi(\dot{z}_1, \dots, \dot{z}_n))^\eta$

② FOR ALL  $M$ -GENERIC FILTERS  $G$  ON  $P$

$\varphi(\dot{z}_{1,G}, \dots, \dot{z}_{n,G})^{M[G]}$  IFF  $(\exists p \in G) (p \Vdash \varphi(\dot{z}_1, \dots, \dot{z}_n))$

① DEFINABILITY OF FORCING ; ② TRUTH LEMMA

NOW EXPLAIN WHY ZFC HOLDS IN  $M[G]$   
(LAST PAGE OF LAST WEEK'S NOTES)

### BACK TO $\neg CH$

COHEN STARTED FROM A MODEL  $M$  THAT SATISFIED GCH.

GÖDEL'S CONSTRUCTIBLE UNIVERSE [1938-40]

WITH COHEN WE TOOK  $P = FN(\omega_2^M \times \omega, 2)$

FROM AN  $M$ -GENERIC FILTER  $G$  WE OBTAINED AN INJECTIVE MAP  $f: \omega_2^M \rightarrow \mathcal{P}^{M[G]}(\omega)$ .

- FROM  $G$  TO  $UG: \omega_2^M \times \omega \rightarrow 2$

- FROM  $UG$  TO  $f(\alpha) = \{n: UG(\alpha, n) = 1\}$

-  $D_{\alpha, m} = \{p: \langle \alpha, m \rangle \in \text{DOM } p\}$  IS DENSE

-  $E_{\alpha, \beta} = \{p: (\exists m) (\langle \alpha, m \rangle, \langle \beta, m \rangle \in \text{DOM } p \wedge p(\alpha, m) \neq p(\beta, m))\}$  IS DENSE IF  $\alpha \neq \beta$ .

GENERAL LEMMA:

LET  $X, Y \in \mathcal{M}$ ; LET  $g: X \rightarrow Y$  BE A MAP IN  $\mathcal{M}[G]$   
THERE IS A MAP  $F: X \rightarrow [Y]^{\leq \omega}$  IN  $\mathcal{M}$   
SUCH THAT  $(\forall x \in X)(g(x) \in F(x))$ .

"EVERY MAP IN  $\mathcal{M}[G]$  IS CAPTURED BY A  
COUNTABLE PIPE FROM  $\mathcal{M}$ ".

PROOF

TAKE A NAME  $\gamma$  IN  $\mathcal{M}$  SUCH THAT  $g = \text{VAL}(\gamma, G)$   
AND  $p \in G$  SUCH THAT  
 $p \Vdash (\gamma : \check{X} \rightarrow \check{Y})$

DEFINE, IN  $\mathcal{M}$ , THE RELATION  $R_\gamma$  BY

$\langle q, \langle x, y \rangle \rangle \in R_\gamma$  IFF  $q \in P$  AND  
 $q \leq p \rightarrow q \Vdash \gamma(\check{x}) = \check{y}$   
 $q \perp p \rightarrow \check{y} = y_0$

WHERE  $y_0$  IS A FIXED ELEMENT OF  $Y$ .

IF  $H$  IS  $\mathcal{M}$ -GENERIC THEN  $R_\gamma[H]$  IS A  
MAP FROM  $X$  TO  $Y$ :

- IF  $p \notin H$  THEN  $R_\gamma[H] = X \times \{y_0\}$
- IF  $p \in H$  THEN  $R_\gamma[H] = \dots$  A MAP.
- $R_\gamma[H] = \text{VAL}(\gamma, H)$   
 $\langle x, y \rangle \in R_\gamma[H] \Leftrightarrow (\exists q \in H)(\langle q, \langle x, y \rangle \rangle \in R_\gamma)$   
 $\Leftrightarrow (\exists q \in H)(q \Vdash \gamma(\check{x}) = \check{y})$   
 $\Leftrightarrow \gamma_G(x) = y.$

DEFINE, IN  $\mathcal{M}$ ,  $F: X \rightarrow \mathcal{P}(Y)$  BY

$F(x) = \{y : (\exists q \in P)(\langle q, \langle x, y \rangle \rangle \in R_\gamma)\}$

CLEARLY  $g(x) \in F(x)$  FOR ALL  $x$  (REALLY!)

- FOR ALL  $x$  THE SET  $F(x)$  IS COUNTABLE  
APPLY AC:  $y \mapsto q_y$  FROM  $F(x)$  TO  $P$   
SUCH THAT  $\langle q_y, \langle x, y \rangle \rangle \in R_\gamma$   
OR  $q_y \Vdash \gamma(\check{x}) = \check{y}$

IF  $y \neq z$  THEN  $q_y \perp q_z$   
(IF  $\kappa \leq q_y, q_z$  THEN  
 $\kappa \Vdash \check{y} = \gamma(\check{x}) = \check{z}$ )

SO  $\{q_y : y \in F(x)\}$  IS AN ANTICHAIN

AND WE HAVE SEEN THAT  
(ANTICHAINS IN  $\mathcal{P}(W_2^M \times W_2)$ )  
ARE COUNTABLE

### COROLLARY

FOR ALL  $\alpha \in ON \wedge M$  WE HAVE  
 $(\alpha \text{ IS A CARDINAL})^M \iff (\alpha \text{ IS A CARDINAL})^{MEG}$

### PROOF

$\Leftarrow$  CLEAR IF  $\beta < \alpha$  AND  $f \in \mathcal{P}^M$  IS A SURJECTION FROM  $\beta$  ONTO  $\alpha$  THEN  $f$  IS ALSO A MEMBER OF  $MEG$  THAT IS A SURJECTION.

IN GENERAL " $\alpha$  IS A CARDINAL" IS DOWNWARD ABSOLUTE.

$\Rightarrow$  LET  $\beta < \alpha$  AND LET  $f \in \mathcal{P}^{MEG}$  BE A MAP FROM  $\beta$  TO  $\alpha$ .

LET  $F : \beta \rightarrow \{\alpha\}^{\omega}$  BE AS IN THE LEMMA THEN  $|\cup \{F(\beta) : \beta < \beta\}^M| \leq |\beta|^M \cdot \aleph_0 < \alpha$ .

WHY: WLOG  $\alpha \geq \aleph_0$ . BECAUSE THE NATURAL NUMBERS AND  $\omega$  ARE ABSOLUTE.

### EXERCISE

PROVE:  $CF^M \alpha = CF^{MEG} \alpha$  FOR ALL  $\alpha$ .

LET'S CALCULATE  $2^{\aleph_0}$  IN  $MEG$ .

WE KNOW  $2^{\aleph_0} \geq \omega_2^M$

ASSUME  $\forall \alpha \in G, \text{val}(\alpha, G) \leq \omega$

LET  $p \in G$  BE SUCH THAT  $p \Vdash \dot{\mathcal{C}} = \dot{\omega}$

MAKE A RELATION AGAIN (OR ANOTHER NAME)

$$\mathcal{C}' = \left\{ \langle \check{\eta}, q \rangle : \left( q \Vdash p \rightarrow q \Vdash \check{\eta} \in \dot{\mathcal{C}} \right) \wedge \left( q \perp p \rightarrow \check{\eta} = \emptyset \right) \right\}$$

THEN  $\emptyset \Vdash \mathcal{C}' \leq \dot{\omega}$

$p \Vdash \mathcal{C}' = \dot{\mathcal{C}}$  AND  $\text{val}(\mathcal{C}', G) = \text{val}(\dot{\mathcal{C}}, G)$

SO ALL SUBSETS OF  $\omega$  COME FROM

RELATIVELY NICE NAMES:  $S$

SUBSETS OF  $\{\check{\eta} : \check{\eta} \in \omega\} \times \mathbb{P}$ .

PROPERTIES OF SUCH A SUBSET/NAME  $\mathcal{C}'$ :

IF  $\langle \check{\eta}, q \rangle \in \mathcal{C}'$  AND  $r \leq q$  THEN  $\langle \check{\eta}, r \rangle \in \mathcal{C}'$   
BECAUSE IF  $q \Vdash \check{\eta} \in \dot{\mathcal{C}}$  THEN  $r \Vdash \check{\eta} \in \dot{\mathcal{C}}$

MAKE  $\mathcal{C}'$  EVEN NICER: TAKE, FOR EVERY  $\check{\eta}$ , A MAXIMAL ANTICHAIN  $A_{\check{\eta}}$  IN  $\{q : \langle \check{\eta}, q \rangle \in \mathcal{C}'\}$

AND LET  $Z'' = \bigcup_{\text{non}} \{x\} \times A_m$ .

NOW WE HAVE

- $Z''$  IS A COUNTABLE SET
- $\text{VAL}(Z'', G) = \text{VAL}(Z', G) = \text{VAL}(Z, G)$ .

SO HOW MANY SUCH NAMES ARE THERE?

WELL, AT MOST

$$| [ \omega \times \text{FN}(\omega_2^{\aleph_0} \times \omega, 2) ]^{\aleph_0} |^{\aleph_0}$$

BUT  $( | \text{FN}(\omega_2^{\aleph_0} \times \omega, 2) | = \aleph_2^{\aleph_0} )^{\aleph_0}$

SO  $| [ \omega \times \text{FN}(\omega_2^{\aleph_0} \times \omega, 2) ]^{\aleph_0} |^{\aleph_0} \leq (\aleph_2^{\aleph_0})^{\aleph_0} = \aleph_2^{\aleph_0}$

BY THE GCH IN  $M$

WE CONCLUDE  $(2^{\aleph_0} = \aleph_2)^{M[G]}$ .

EXERCISE  $(2^{\aleph_1} = \aleph_2)^{M[G]}$

SO  $2^{\aleph_0} = 2^{\aleph_1}$  IS CONSISTENT

SOLOVAY [1965]:

$2^{\aleph_0}$  CAN BE ANYTHING IT OUGHT TO BE

LET  $\kappa$  BE AN INFINITE CARDINAL OF  $M$  SUCH THAT  $\text{CF} \kappa > \aleph_0$  IN  $M$ .

THEN THERE IS AN EXCELLENT EXTENSION  $N$  OF  $M$  IN WHICH  $2^{\aleph_0} = \kappa$ .

[THIS MODEL  $N$  SATISFIES GCH]

EXCELLENT: -  $\text{ON} \cap N = \text{ON} \cap M$

-  $\text{CF}^N \alpha = \text{CF}^M \alpha$  FOR ALL  $\alpha$ .

HOW: TAKE  $G$  TO BE  $\aleph_0$ -GENERIC ON  $\text{FN}(\kappa \times \omega, 2)$

EVERYTHING GOES THROUGH AS IN THE CASE  $\kappa = \omega_2^M$ .

YOU GET  $2^{\aleph_0} = \kappa$  FOR ALL  $\lambda < \kappa$ .

EXERCISE

• IF  $\kappa = \aleph_m^M$  FOR SOME  $m \in \omega \setminus \{0, 1\}$

THEN -  $2^{\aleph_m^M} = \aleph_m^M$  FOR  $m \in M$

-  $2^{\aleph_m^M} = \aleph_{m+1}^M$  FOR  $m \geq m$ .

• IF  $\kappa = \aleph_{\omega+1}^M$  THEN  $2^{\aleph_m^M} = \aleph_{\omega+1}^M$  FOR ALL  $m \in \omega$  AND  $2^{\aleph_\omega} = \aleph_{\omega+1}^M$  AS WELL

• WHAT HAPPENS TO  $2^{\aleph_0}$  IF WE USE  $\text{FN}(\aleph_\omega \times \omega, 2)$ ?

SO WE HAVE THE FOLLOWING CONSISTENCIES

- GÖDEL:  $AC + GCH$
- COHEN:  $2^{\aleph_0} = 2^{\aleph_1} = \aleph_2^{\aleph_1}$   
[CALCULATE  $2^{\aleph_\alpha}$  FOR  $\alpha \geq 2$ ]
- COHEN/SOLOVAY  $2^{\aleph_0} = \aleph_\alpha^{\aleph_1}$  AS LONG AS  $CF \alpha > \aleph_0$

WHAT ELSE?

- LET  $\kappa$  BE REGULAR AND  $\lambda$  SUCH THAT  $CF \lambda > \kappa$ . THEN WE CAN ACHIEVE
  - $2^\kappa = \lambda$
  - $2^{\aleph_\alpha} = \aleph_{\alpha+1}^\lambda$  IF  $\aleph_\alpha < \kappa$ . [SOLOVAY, EXCELLEN] ]

FOR DEFINITENESS LET  $\kappa = \aleph_1$ .

HOMEWORK 12 (42)

LET  $P = FN(\lambda \times \omega_1, 2, \aleph_1)$

THE SET OF FUNCTIONS  $p$  WITH  $DOM p \in [\lambda \times \omega_1]^{\leq \aleph_0}$  AND  $RNG p \subseteq 2$ .

ORDERED AS BEFORE:  $p \leq q$  IF  $p \supseteq q$ .

HOMEWORK: ALL ANTICHAINS HAVE CARDINALITY AT MOST  $\aleph_1$  [SOLOVAY STARTED FROM GCH]

SO EVERY MAP  $g: X \rightarrow Y$  WITH  $X, Y \in \Pi$  IS CAPTURED BY A PIPE  $F: X \rightarrow [Y]^{\leq \aleph_1}$

SO, IF  $CF \alpha > \aleph_1$ , IN  $\Pi$  THEN  $CF \alpha > \aleph_1$ , IN  $\Pi[G]$ .

WHAT IF  $CF \alpha \leq \aleph_1$ ? CAN WE GET  $CF \alpha = \aleph_0$ ?

NEW PHENOMENON:

IF  $g: \omega \rightarrow X$  WITH  $X \in \Pi$  AND  $g \in \Pi[G]$

THEN  $g \in \Pi$ .

LET  $\gamma$  BE A NAME FOR  $g$ , SO  $g = VAL(\gamma, G)$

LET  $p \in G$  BE SUCH THAT  $p \Vdash \gamma: \check{\omega} \rightarrow \check{X}$

LET  $q \leq p$ .  $q \Vdash p$

THEN  $q \Vdash (\exists z)(z \in \check{X} \wedge \gamma(\check{\delta}) = z)$

SO THERE ARE  $\tau \in q$  AND  $\sigma$  SUCH THAT

$\tau \Vdash (\sigma \in \check{X} \wedge \gamma(\check{\delta}) = \sigma)$

BUT THEN THERE ARE  $s \leq \tau$  AND  $\langle \check{x}, \mathbb{1} \rangle \in \check{X}$

SUCH THAT  $s \Vdash (\check{x} = \sigma \wedge \gamma(\check{\delta}) = \sigma)$

OR  $s \Vdash (\gamma(\check{\delta}) = \check{x})$



RECURSIVELY :

TAKE  $q_0 \in q$  AND  $x_0 \in X$

SUCH THAT  $q_0 \Vdash \gamma(\check{0}) = \check{x}_0$

AND  $q_{m+1} \in q_m$  AND  $x_{m+1} \in X$

SUCH THAT  $q_{m+1} \Vdash \gamma(\check{m+1}) = \check{x}_{m+1}$

EACH  $q_m$  IS A COUNTABLE FUNCTION

AND SO  $\dot{q} = \bigcup_{m \in \omega} q_m$  IS A COUNTABLE FUNCTION TOO

WE HAVE  $\dot{q} \in \text{Fn}(\lambda \times \omega_1, 2, S_1)$

AND FOR ALL  $n$   $\dot{q} \Vdash \gamma(\check{n}) = \check{x}_n$

DEFINE  $\dot{f} = \{ \langle m, x_m \rangle : m \in \omega \}$  NOTE  $\dot{f} \in M$

THEN  $\dot{q} \Vdash (\gamma = \dot{f})$

CONCLUSION  $D_\gamma = \{ \dot{q} \in P : (\exists \dot{f} \in X^\omega) (\dot{q} \Vdash \gamma = \dot{f}) \}$

IS DENSE BELOW  $p$ .

TAKE  $\dot{q} \in D_\gamma \cap G$  WITH ITS  $\dot{f}$

THEN  $\dot{q} = \dot{f} \in M$

So:  $P^\omega(\omega) = P^{\text{REG}}(\omega)$  WE STILL HAVE  $2^{S_0} = S_1$

- IF  $\text{CF}^\omega \alpha = S_1$  THEN  $\text{CF}^{\text{REG}} \alpha = S_1$

NO NEW FUNCTIONS FROM  $\omega$  TO  $\alpha$ .

TO SEE  $2^{S_1} = \lambda$  DO THE SAME COMPUTATIONS

AS FOR THE CASE OF  $\text{Fn}(\omega_2 \times \omega, 2)$

-  $2^{S_1} \geq \lambda$  BECAUSE  $G$  GIVES AN INJECTION FROM  $\lambda$  INTO  $P^{\text{REG}}(\omega_1)$

-  $2^{S_1} \leq \lambda$  USE ANTICHAINS AGAIN [AND GCH]

-  $|\text{Fn}(\lambda \times \omega_1, 2, S_1)| \leq \lambda^{S_0} = \lambda$

-  $|\text{Fn}(\lambda \times \omega_1, 2, S_1)^{S_1}| \leq \lambda^{S_1} = \lambda$ .