

# Set Theory

## FIFTH LECTURE

11 October 2021

### Ordinal Arithmetic

We shall now define addition, multiplication and exponentiation of ordinal numbers, using Transfinite Recursion.

**Definition 2.18 (Addition).** For all ordinal numbers  $\alpha$

- (i)  $\alpha + 0 = \alpha$ ,
- (ii)  $\alpha + (\beta + 1) = (\alpha + \beta) + 1$ , for all  $\beta$ ,
- (iii)  $\alpha + \beta = \lim_{\xi \rightarrow \beta} (\alpha + \xi)$  for all limit  $\beta > 0$ .

**Definition 2.19 (Multiplication).** For all ordinal numbers  $\alpha$

- (i)  $\alpha \cdot 0 = 0$ ,
- (ii)  $\alpha \cdot (\beta + 1) = \alpha \cdot \beta + \alpha$  for all  $\beta$ ,
- (iii)  $\alpha \cdot \beta = \lim_{\xi \rightarrow \beta} \alpha \cdot \xi$  for all limit  $\beta > 0$ .

**Definition 2.20 (Exponentiation).** For all ordinal numbers  $\alpha$

- (i)  $\alpha^0 = 1$ ,
- (ii)  $\alpha^{\beta+1} = \alpha^\beta \cdot \alpha$  for all  $\beta$ ,
- (iii)  $\alpha^\beta = \lim_{\xi \rightarrow \beta} \alpha^\xi$  for all limit  $\beta > 0$ .

$$\begin{aligned}\alpha + \beta &= \bigcup_{\xi < \beta} \alpha + \xi \\ &= \bigcup \{ \alpha + \xi, \xi < \beta \} \\ &= \sup_{\xi < \beta} \alpha + \xi\end{aligned}$$

DEFINITION HAVE HIGH DEGREE OF LEFT/RIGHT ASYMMETRY.

We saw:

$$1 + \omega = \omega$$

$$\neq \omega + 1$$

$$2 \cdot \omega = \omega$$

$$\neq \omega \cdot 2$$

**Theorem 2.26 (Cantor's Normal Form Theorem).** Every ordinal  $\alpha > 0$  can be represented uniquely in the form

$$\alpha = \omega^{\beta_1} \cdot \underline{k_1} + \dots + \omega^{\beta_n} \cdot \underline{k_n},$$

where  $n \geq 1$ ,  $\alpha \geq \underline{\beta_1} > \dots > \underline{\beta_n}$ , and  $k_1, \dots, k_n$  are nonzero natural numbers.

↑  
GROUP INTERACTION #3

# SWALLOWING

① If  $\alpha < \beta$ , then  $\omega^\alpha + \omega^\beta = \omega^\beta$ .

[By subtraction (GI#3), we find  $\delta > 0$

s.t.  $\beta = \alpha + \delta$

$$\omega^\alpha + \omega^\beta = \omega^\alpha + \omega^{\alpha + \delta}$$

$$= \omega^\alpha + \omega^\alpha \cdot \omega^\delta$$

$$= \omega^\alpha (1 + \omega^\delta)$$

$$\underbrace{(\omega^\delta)}_{\omega^\delta}$$

$$= \omega^\alpha \cdot \omega^\delta = \omega^{\alpha + \delta}$$

$$= \omega^\beta$$

② Consequence.

If  $\gamma < \omega^\beta$ , then  $\gamma + \omega^\beta = \omega^\beta$ .

Consider CNF of  $\gamma$ :

Since  $\gamma < \omega^\beta$

$\rightarrow$

$\gamma_1 < \beta$ .

$$\gamma = \omega^{\gamma_1} + \dots + \omega^{\gamma_n} \quad \text{with } \gamma_1 \geq \gamma_2 \geq \dots \geq \gamma_n$$

$$\leq \omega^{\gamma_1} \cdot n < \omega^{\gamma_1} \cdot \omega = \omega^{\gamma_1 + 1}$$

$$\gamma + \omega^\beta \leq \omega^{\gamma_1} \cdot n + \omega^\beta = \omega^\beta$$

Apply ①  $n$  times

$$\leq \omega^\beta$$

③ Remarks about distributivity.

$$\alpha(\beta + \gamma) = \alpha\beta + \alpha\gamma \quad \text{holds,}$$

$$(\alpha + \beta)\gamma \neq \alpha\gamma + \beta\gamma. \quad \text{in general}$$

$$\begin{aligned} \underline{(\omega + 1) \cdot 2} &= (\omega + 1) + (\omega + 1) \\ &= \omega + \underbrace{(1 + \omega)}_{= \omega} + 1 \end{aligned}$$

$$= \underline{\omega \cdot 2 + 1}$$

$$\underline{2 \cdot (\omega + 1)} = 2 \cdot \omega + 2 \cdot 1$$

$$= \underline{\omega + 2}$$

$$\begin{aligned} \text{④ } (\omega^\alpha + \gamma) \cdot 2 &= \omega^\alpha + \underbrace{\gamma + \omega^\alpha}_{\text{Apply ②}} + \gamma \\ &\text{where } \gamma < \omega^\alpha \end{aligned}$$

$$= \omega^\alpha \cdot 2 + \gamma$$

$$\text{In general, } (\omega^\alpha + \gamma) \cdot u = \boxed{\omega^\alpha \cdot u + \gamma}$$

$$\text{Therefore } (\omega^\alpha + \gamma) \cdot \omega = \bigcup_{n \in \mathbb{N}} \underline{(\omega^\alpha + \gamma) \cdot n}$$

$$= \omega^\alpha \cdot \omega = \omega^{\alpha+1}$$

## EXAMPLE

$$\underbrace{(\omega^3 + \omega^2 \cdot 2 + 1)}_{\alpha} (\omega^2 + \omega \cdot 3 + 2)$$

$$= \alpha \omega^2 + \alpha \omega \cdot 3 + \alpha \cdot 2$$

$$\alpha \cdot 2 = (\omega^3 + \omega^2 \cdot 2 + 1) 2 = \omega^3 + \cancel{\omega^2 \cdot 2} + \cancel{1} + \omega^3 + \omega^2 \cdot 2 + 1$$
$$= \omega^3 \cdot 2 + \omega^2 \cdot 2 + 1$$

$$\alpha \cdot \omega = (\omega^3 + \omega^2 \cdot 2 + 1) \cdot \omega = \omega^4 \quad \text{by } \textcircled{4}$$

$$\alpha \cdot \omega^2 = \omega^5 \quad \text{by } 4$$

$$= \omega^5 + \omega^4 \cdot \underline{3} + \omega^3 \cdot \underline{2} + \omega^2 \cdot \underline{2} + \underline{1}$$

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$$(\underline{x^3} + \underline{2x^2} + \underline{1})(\underline{x^2} + \underline{3x} + \underline{2})$$

$$= \underline{x^5} + \underline{5x^4} + \underline{8x^3} + \underline{5x^2} + \underline{3x} + \underline{2}$$

# NORMAL OPERATIONS & FIXED PTS

Remember "definable operations": a functional formula  $\Phi$ .

We call such a definable operation an ordinal operation if for all  $\alpha$  ordinal there is a (unique)  $\beta$  ordinal s.t.  $\Phi(\alpha, \beta)$ .

In Jech's book, he uses symbols like  $F(\alpha) = \beta$ .

Definition An ordinal operation  $F$  is called

NORMAL if  $[\Rightarrow \text{f.o. } \alpha \quad F(\alpha) \geq \alpha]$

$$\textcircled{1} \quad \alpha < \beta \longrightarrow F(\alpha) < F(\beta)$$

$\textcircled{2}$  [CONTINUITY]

if  $\lambda$  is a limit ordinal, then

$$F(\lambda) = \bigcup \{ F(\alpha); \alpha < \lambda \}$$

that's exactly what we had in ordinal arithmetic.

Examples are the operations

$$\begin{array}{l} \beta \longmapsto \alpha + \beta \\ \beta \longmapsto \alpha \cdot \beta \\ \beta \longmapsto \alpha / \beta \end{array}$$

All of them are normal.

Def. If  $F$  is an ordinal operation,  $\alpha$  is a fixed pt of  $F$  if  $F(\alpha) = \alpha$ .

Thm If  $F$  is a normal ordinal operation and  $\alpha$  is an ordinal, then there is some  $\gamma > \alpha$  s.t.  $\gamma$  is a fixed pt. for  $F$ .

Proof. Define the following sequence by recursion:

$$\gamma_0 := \alpha + 1.$$

$$\gamma_{u+1} := F(\gamma_u).$$

Needs Acc.  
Thm w/o  
fixed range.

Case 1. One of the  $\gamma_u$  is a fixed pt.  
Done.

Case 2. Not. That means  $\gamma_u$  is a strictly increasing seq. of ordinals. Therefore

$\gamma := \bigcup \{ \gamma_u; u \in \mathbb{N} \}$  is a limit ordinal.

$$F(\gamma) = \bigcup \{ F(\alpha); \alpha < \gamma \}$$

$$\stackrel{\text{by continuity}}{=} \bigcup \{ F(\gamma_u); u \in \mathbb{N} \} \quad (\text{q.e.d.})$$

$$= \bigcup \{ \gamma_{u+1}; u \in \mathbb{N} \} = \gamma$$

## Illustration

Take function  $\beta \mapsto \omega + \beta$ .

This is normal, therefore has fixed pts.

What is the smallest fixed pt?

$$\gamma_0 = 0$$

$$\gamma_1 = \omega + 0 = \omega \cdot 1$$

$$\gamma_2 = \omega + (\omega + 0) = \omega \cdot 2$$

$$\gamma_3 = \omega + \gamma_2 = \omega \cdot 3$$

$$\gamma_n = \omega \cdot n$$

$$\bigcup \{ \gamma_n; n \in \mathbb{N} \} = \omega \cdot \omega = \omega^2.$$

[This is in line with (1) and (2) above since  $\omega^2$  is a fixed pt for all operators

$$\beta \mapsto \gamma + \beta \quad \text{where } \gamma < \omega^2.]$$

Remark. These fixed pts are closely related to the questions on gamma numbers and delta numbers on QI#3.

Back to our original motivation:

- We wanted ordinals as canonical representatives from the isomorphism classes of wellorders.
- We have proved that if  $\alpha, \beta$  ordinals and  $(\alpha, \epsilon) \cong (\beta, \epsilon)$ , then  $\alpha = \beta$ .
- We still need to show that each isomorphism class contains an ordinal.

Theorem If  $(X, <)$  is a wellorder,  
(ZF) there is a unique ordinal  $\alpha$  s.t.

$$(X, <) \cong (\alpha, \epsilon).$$

## REPRESENTATION THEOREM FOR WELLORDERS

Proof. We define by recursion a function  
on  $X$ :

$$F(x) = \text{ran}(F \upharpoonright <[x])$$



Claim 1  $F$  is order-preserving between  $<$  and  $\in$   
if  $x < y$ , then  $F(x) \in F(y)$

[By recursion equations.]

Define  $A := \text{ran}(F)$ .

Claim 2  $A$  is an ordinal.

$A$  is well ordered by  $\in$  [follows from Claim 1]

$A$  is transitive:

$$\forall v \in w \in A$$

$\stackrel{''}{=} F(x)$  by rec. eq.  $v = F(y)$  for some  $y < x$

$$\Rightarrow v \in A.$$

So  $(X, <) \cong (A, \in)$ .

q.e.d.

The theorem that there is no set of all ordinals means that INTUITIVELY there must be lots of them.

But: are there even uncountable ordinals?

RECURSION EQUATION

$$F(x) = \text{ran}(F \upharpoonright <[x])$$

Theorem

## HARTOGS'S THEOREM

If  $X$  is any set, then there is an ordinal  $\alpha$  s.t. there is no injection from  $\alpha$  into  $X$ .

[Consequence: if  $X := \mathbb{N}$ , then  $\alpha$  is an uncountable ordinal.]

Proof

Idea: Consider  $A_X$ , the collection of all ordinals  $\alpha$  s.t. there is an inj. from  $\alpha$  to  $X$ .

Show that  $A_X$  is a set. Show that  $A_X$  is transitive. So  $A_X$  is an

ordinal:  $A_X \notin A_X$ .

Consider

$W_X := \{ (A, \mathcal{R}) ; A \subseteq X \text{ and } (A, \mathcal{R}) \text{ is a wellorder} \}$

This is a set by Separation and the usual operations  $\mathcal{P}, \times, \dots$

By Representation Thm, each  $(A, \mathcal{R}) \in W_X$  corresponds to a unique ordinal  $\alpha$  s.t.

$$(A, \mathcal{R}) \cong (\alpha, \in).$$

This is a functional formula, so by Replacement

$$\hat{W}_X := \{ \alpha ; \exists (A, R) \in W_X \text{ s.t. } (\alpha, \epsilon) \cong (A, R) \}$$

is a set.

Claim  $\alpha \in A_X \iff \alpha \in \hat{W}_X$

" $\Leftarrow$ " The isomorphism  $(\alpha, \epsilon) \cong (A, R)$  is an injection from  $\alpha$  to  $X$ .

" $\Rightarrow$ ": Suppose  $f: \alpha \rightarrow X$  is an injection.

$A := \text{ran}(f)$ , so  $f: \alpha \rightarrow A$  is a bijection.

Define  $R \subseteq A \times A$  by

$$f(\beta) R f(\gamma) \iff \beta \in \gamma$$

Then  $f$  is an isom. between

"Coding"  $(\alpha, \epsilon)$  and  $(A, R)$

$\implies \alpha \in \hat{W}_X$  . q.e.d.

Claim  $A_X$  is transitive.

[ $\alpha \in \beta \in A_X \xrightarrow{\text{T.S.}} \alpha \in A_X$ . If  $\alpha \in \beta \implies \alpha \in \beta \implies$  if  $\beta \in A_X$ , then  $\alpha \in A_X$ .]

## Remark.

This is a constructive proof. It provides a concrete example  $A_X$  of an ordinal that doesn't inject into  $X$  and it is the least example:

if  $\gamma$  is any ordinal:

$$\textcircled{1} \quad \gamma \in A_X \implies \gamma \text{ injects into } X$$

$$\textcircled{2} \quad \gamma = A_X$$

$$\textcircled{3} \quad A_X \in \gamma \implies A_X \subseteq \gamma \\ \implies \gamma \text{ does not inject into } X.$$

Concretely: The set of all countable ordinals is the least uncountable ordinal.

Definition This ordinal is called the Hartogs-  
Aleph of  $X$ , in symbols,

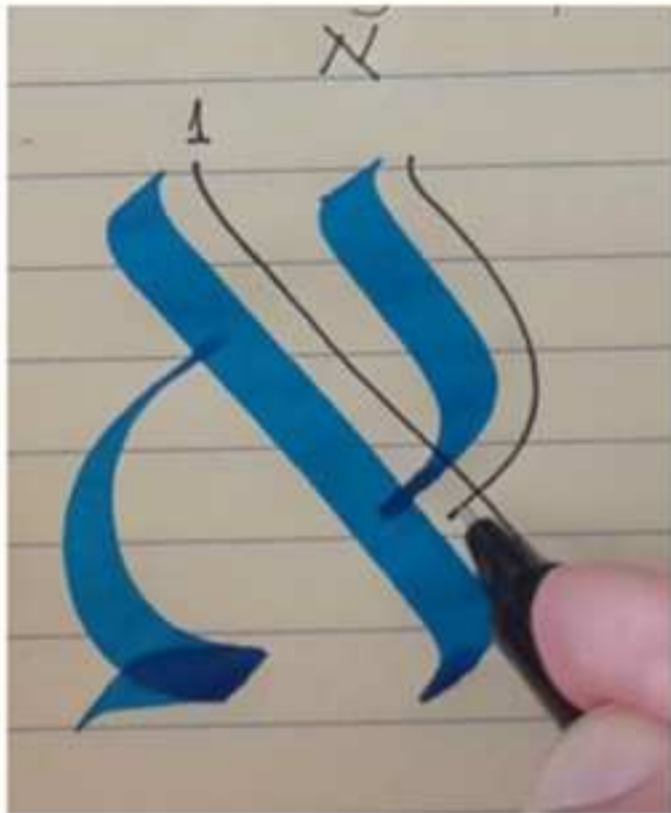
$$\aleph'(X) := \{ \alpha ; \alpha \text{ injects into } X \}$$

The first uncountable ordinal

$$\omega_1 := \aleph_1 := \aleph(\mathbb{N})$$

After that

$$\omega_2 := \aleph_2 := \aleph(\aleph_1)$$



The last missing ZF axiom:

FOUNDATION

[Jech: REGULARITY]

$$\text{FST} \subseteq Z \subseteq \text{ZF}_0$$

FST + I<sub>f</sub>

Z + Regl.


There is no intrinsic argument for Foundation, but it makes the theory of models of set theory much nicer.

Another point: so far, we have not yet excluded the possibility of a set containing itself as an element.

ANSC Axiom of no self-containment

$$\forall x (x \notin x)$$

Note that


$$\models \text{ANSC},$$

but we have  $\exists x \exists y (x \in y \wedge y \in x)$

# The Cumulative Hierarchy of Sets

# VON NEUMANN HIERARCHY.

We define, by transfinite induction,

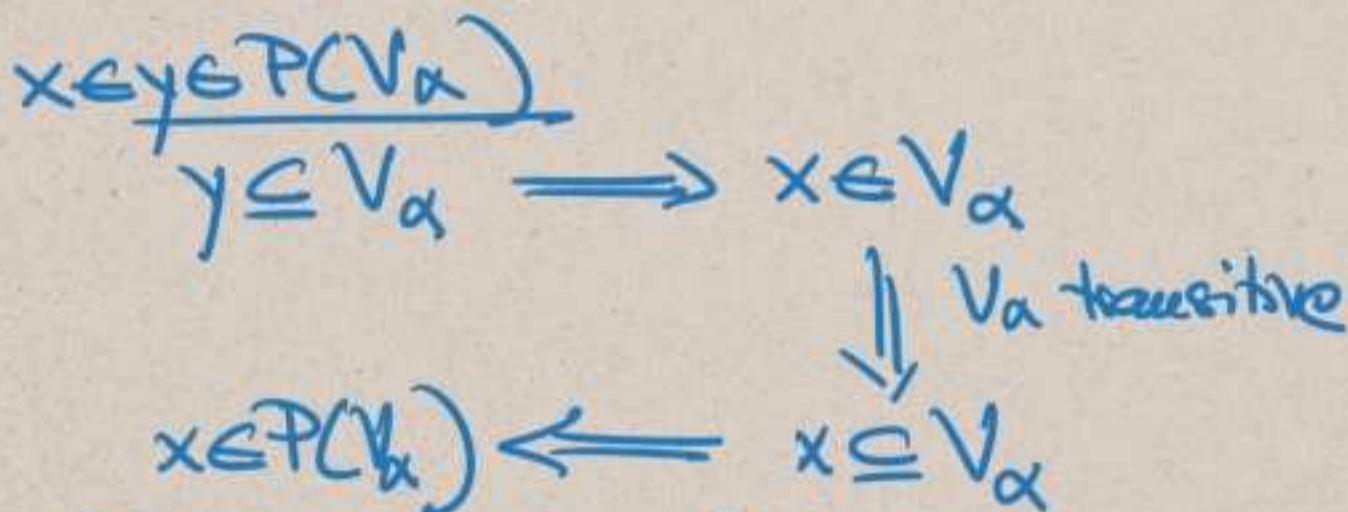
$$\begin{aligned} V_0 &= \emptyset, & V_{\alpha+1} &= P(V_\alpha), \\ V_\alpha &= \bigcup_{\beta < \alpha} V_\beta & \text{if } \alpha \text{ is a limit ordinal.} \end{aligned}$$

The sets  $V_\alpha$  have the following properties (by induction):

- (i) Each  $V_\alpha$  is transitive.
- (ii) If  $\alpha < \beta$ , then  $V_\alpha \subset V_\beta$ .
- (iii)  $\alpha \subset V_\alpha$ .

Proof (i):  $V_0$  is transitive. & it is limit  
 If  $V_\alpha$  is transitive for  $\alpha < \lambda$ , then  
 $V_\lambda$  is transitive.

To Show: If  $V_\alpha$  is transitive, then so is  $P(V_\alpha)$ .



Proof (ii) Transfinite ind. on  $\beta$ .

If  $\beta = 0$ : nothing to show.

If  $\beta$  limit, then the claim follows from definition.

The successor case reduces to  $\beta = \gamma + 1$  and  $V_\gamma \subseteq V_{\gamma+1}$ . Suppose  $x \in V_\gamma \implies x \subseteq V_\gamma$  (trans.)  $\implies x \in V_{\gamma+1}$ .

Strengthening of (iii):

$$O_\alpha := \{ \beta \in V_\alpha; \beta \text{ is an ordinal} \}$$

Claim  $O_\alpha = \alpha$ .

Proof by induction  $\alpha = 0$ . Clear.

If  $\beta$  is a limit,

$$O_\beta = \bigcup \{ O_\alpha; \alpha < \beta \}$$

$$\stackrel{IH}{=} \bigcup \{ \alpha; \alpha < \beta \}$$

$$= \beta$$

Suppose  $O_\alpha = \alpha$ . Show that  $O_{\alpha+1} = \alpha+1$ .

① Show that  $\alpha \in O_{\alpha+1}$ .

We have by IH,  $\alpha \subseteq V_\alpha$ , so  $\alpha \in P(V_\alpha) = V_{\alpha+1}$ .

② Show that  $\alpha+1 \notin O_{\alpha+1}$ .

Otherwise  $\alpha+1 \in V_{\alpha+1}$

$$\Rightarrow \alpha+1 \subseteq V_\alpha$$

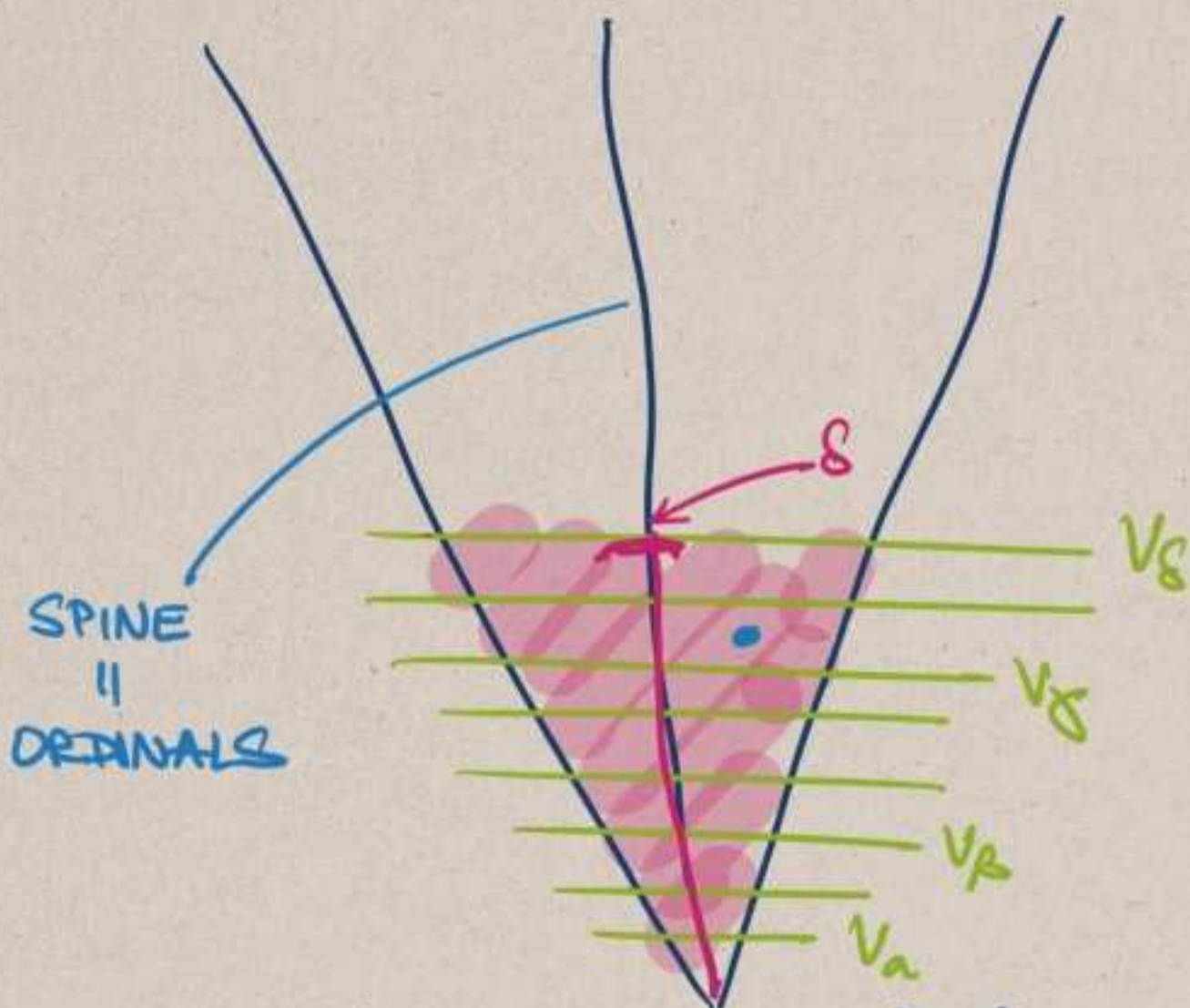
$$\Rightarrow \alpha \in \alpha+1 \subseteq V_\alpha$$

$$\Rightarrow \alpha \in V_\alpha \text{ in contradiction to } O_\alpha = \alpha.$$

$$O_\alpha = \alpha.$$

q.e.d.





This stratified picture of the cumulative hierarchy allows us to define a rank function (MIRIMANOFF rank)

if  $x$  is in the cumulative hierarchy (i.e.,  $\exists \alpha \ x \in V_\alpha$ )

$$\rho(x) := \alpha \iff x \in V_{\alpha+1} \setminus V_\alpha$$

This gives us the opportunity to prove properties in the cumulative hierarchy by induction on the Mirimanoff rank.

# AXIOM OF FOUNDATION / REGULARITY.

$$\forall x (x \neq \emptyset \longrightarrow \exists m (m \in x \wedge m \cap x = \emptyset))$$

FST +

Foundation implies ANSC:

if  $x \in x$ , form  $y := \{x\}$ .

if  $m$  is any element of  $y$

$[m = x]$ ,

then  $m \cap y \ni x$ .

$\implies m \cap y \neq \emptyset$ .

It also excludes:



and the existence of an infinite descending

chain:  $f$  with  $\text{dom}(f) = \mathbb{N}$

s.t.  $f(n+1) \in f(n)$ .

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$$ZF = ZF_0 + \text{Foundation}.$$

Theorem ( $\epsilon$ -induction)

(ZF) If  $\varphi$  is any formula s.t.

for all  $x$   
if  $\forall y (y \in x \rightarrow \varphi(y))$   
 $\rightarrow \varphi(x)$

then  $\varphi$  holds for all  $x$ .

Proof. If  $x$  is any set, let us define by recursion

$$f(0) := x$$

$$f(n+1) := \bigcup f(n)$$

By Rec. Thm. w/o range we get a function  $f$  with  $\text{dom}(f) = \mathbb{N}$  and

can form  $\bigcup_{n \in \mathbb{N}} f(n) =: TC(x)$

TRANSITIVE CLOSURE  
OF  $x$

Then  $TC(x)$  is transitive and  $x \subseteq TC(x)$ .

[In  $ZF_0$ , every set is contained in a transitive set.]

Suppose  $\varepsilon$ -induction is false and  
derive a contradiction.

(\*) for any  $x$  if  $\forall y (y \in x \rightarrow \varphi(y)) \rightarrow \varphi(x)$

But there is some  $x_0$  s.t.  $\varphi(x_0)$  fails.

Case 1. For all  $y \in x_0$ ,  $\varphi(y)$  holds.

This contradicts (\*).

Case 2. For some  $y \in x_0$ ,  $\varphi(y)$  doesn't hold.

$$\{y \in x_0; \neg \varphi(y)\} \neq \emptyset.$$

Let  $T := TC(x_0)$ .

$$W := \{y \in T; \neg \varphi(y)\} \neq \emptyset.$$

By Foundation, find  $m$  minimal in  $W$ ,  
i.e.,  $m \in W$  and  $\underline{m \cap W = \emptyset}$ .

$$\{z \in m; z \in W\} = \emptyset$$

$$\{z \in m; z \notin W\} = m$$

$\Rightarrow$  for all  $z \in m$  we  
have  $\varphi(z)$

(\*)  $\Rightarrow \varphi(m)$ . Contradiction.

q.e.d.

## Corollary (ZF)

$$\forall x \exists \alpha \quad x \in V_\alpha.$$

Proof. Suppose  $x$  is an arbitrary set.  
And suppose for all  $y \in x$  there  
is  $\alpha$  s.t.  $y \in V_\alpha$ .

By Replacement, form

$$\{g(y)^{+1}; y \in x\}$$

$$\text{and } \gamma := \bigcup \{g(y)^{+1}; y \in x\}$$

is an ordinal.

Then if  $y \in x$ ,  $y \in V_\gamma$ .

Thus  $x \subseteq V_\gamma \implies x \in V_{\gamma+1}$ .  
q.e.d.

Thus, Foundation implies that the von Neumann hierarchy is the set-theoretic universe.

# Consistency proofs with the von Neumann hierarchy

Most axioms are  $\boxed{\forall \exists}$

So showing that  $V_\alpha \models \Phi$  is about analysing what the axiom does to the rank of  $V_\alpha$ :

Ex.  $x, y$   $\rho(x) = \alpha$   $\rho(y) = \beta$  w.l.o.g.  $\alpha \leq \beta$

So, if  $\rho(\{x, y\}) = \beta + 1$ .

$\alpha, \beta \mapsto \max(\alpha, \beta) + 1$ ,  
then  $V_\gamma \models \text{Pair}$ .

Similarly for

Union:  $\rho(x) = \alpha$   
then  $\rho(Ux) \leq \alpha$ .

Power set:  $\rho(x) = \alpha$   
then  $\rho(P(x)) = \alpha + 1$

Separation

If  $\rho(x) = \alpha$ , then  $\rho(\{y \in x; \varphi(y)\}) = \alpha$ .

Together

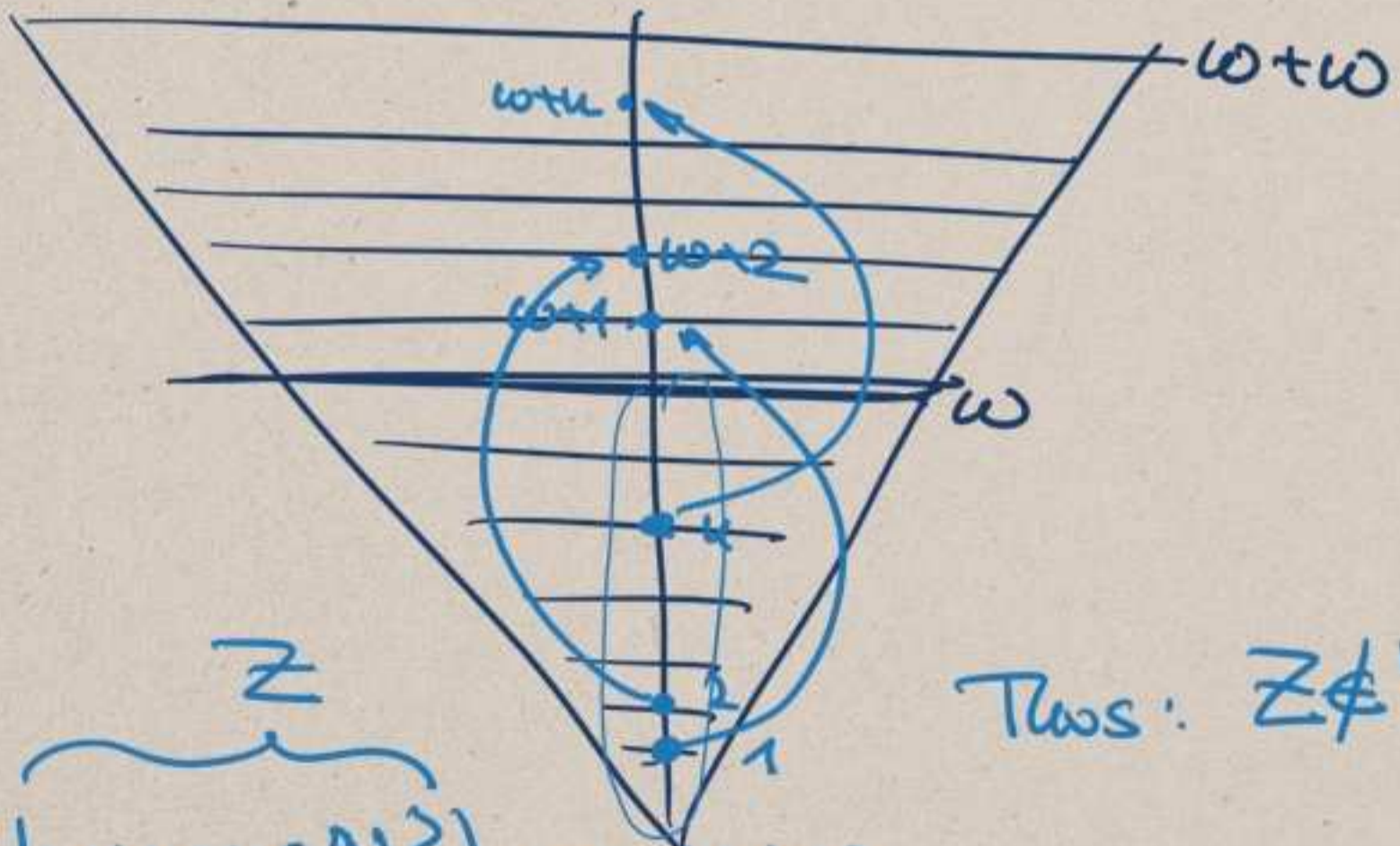
if  $\gamma$  is limit, then  $V_\gamma \neq \text{FST}$ .

if  $\omega \in V_\gamma$ , then  $V_\gamma \neq \text{lcf}$ .

$\Leftrightarrow$   
 $\gamma \geq \omega + 1$

Together

if  $\gamma > \omega$  is a limit, then  $V_\gamma \neq \mathbb{Z}$ .



Thus:  $\mathbb{Z} \notin V_{\omega+\omega}$ .

$$f(\{\omega+u; u \in \mathbb{N}\}) = \omega+\omega$$