

## Exercise 21

Prove the Knaster-Tarski Fixed Point Theorem, the Banach decomposition Theorem and derive the cantor-Schröder-Bernstein Theorem.

**Knaster-Tarski Fixed Point Theorem (KT-FPT):** Let  $X$  be a set and  $F : P(X) \rightarrow P(X)$  a  $\subseteq$ -monotone function. Then  $F$  has a fixed point.

**Banach Decomposition theorem (BDT):** Let  $X$  and  $Y$  be sets and  $f : X \rightarrow Y$  and  $g : Y \rightarrow X$  be arbitrary functions. Then there are disjoint decompositions  $X = X_1 \cup X_2$  and  $Y = Y_1 \cup Y_2$  such that  $f[X_1] = Y_1$  and  $f[Y_2] = X_2$ .

**Cantor-Schröder-Bernstein Theorem (CSBT):** Let  $X$  and  $Y$  be sets. If there is an injection from  $X$  to  $Y$  and one from  $Y$  to  $X$ , then there exists a bijection between  $X$  and  $Y$ .

We start by proving KI-FPT. Let  $X$  be a set and  $F : P(X) \rightarrow P(X)$  a  $\subseteq$ -monotone function. Define the following set:

$$\mathcal{A} := \{A \subseteq X \mid A \subseteq F(A)\}$$

Consider the following set:  $B = \bigcup \mathcal{A}$ . We will show  $B = F(B)$  by showing two inclusions:

- $B \subseteq F(B)$   
Let  $b \in B$ . Then there is  $A \in \mathcal{A}$  such that  $b \in A$ . As  $A \subseteq F(A)$ , we see  $b \in F(A)$ . By monotonicity of  $F$  and  $A \subseteq B$ , we see  $F(A) \subseteq F(B)$ . So  $b \in F(B)$ . Thus,  $B \subseteq F(B)$ .
- $F(B) \subseteq B$   
As  $B \subseteq F(B)$ ,  $F(B) \subseteq F(F(B))$  by monotonicity. Thus,  $F(B) \in \mathcal{A}$  and so,  $F(B) \subseteq B$ .

We conclude that  $B$  is a fixed point of  $F$ .

Now, we prove BDT. Let  $X$  and  $Y$  be sets and  $f : X \rightarrow Y$  and  $g : Y \rightarrow X$  functions. Consider the following function:

$$F : P(X) \rightarrow P(X) : S \mapsto X \setminus g[Y \setminus f[S]]$$

We show this function is monotone: Let  $A$  and  $B$  be  $X$ -subsets such that  $A \subseteq B$ , then:

$$\begin{aligned} A \subseteq B &\rightarrow f[A] \subseteq f[B] \\ &\rightarrow Y \setminus f[B] \subseteq Y \setminus f[A] \\ &\rightarrow g[Y \setminus f[B]] \subseteq g[Y \setminus f[A]] \\ &\rightarrow X \setminus g[Y \setminus f[A]] \subseteq X \setminus g[Y \setminus f[B]] \\ &\rightarrow F(A) \subseteq F(B) \end{aligned}$$

So indeed,  $F$  is monotone. By KT-FPT, this function has a fixed point  $C$ :

$$C = F(C) = X \setminus g[Y \setminus f[C]]$$

Then, also:  $X = C \cup g[Y \setminus f[C]]$  and  $Y = f[C] \cup (Y \setminus f[C])$ . So, we choose  $X_1 = C$ ,  $X_2 = g[Y \setminus f[C]]$ ,  $Y_1 = f[C]$  and  $Y_2 = Y \setminus f[C]$  (note these are disjoint). This proves BDT.

Finally, we will derive CSBT. Let  $X$  and  $Y$  be sets and  $f : X \rightarrow Y$  and  $g : Y \rightarrow X$  be injective functions. By BDT, there are disjoint compositions  $X = X_1 \cup X_2$  and  $Y = Y_1 \cup Y_2$  such that  $f[X_1] = Y_1$  and  $g[Y_2] = X_2$ . As  $f$  and  $g$  are injections, so are  $f \upharpoonright X_1$  and  $g \upharpoonright Y_2$ . Note

that these functions are reversible when domain and range are restricted:  $f$  acts as a bijection between  $X_1$  and  $Y_1$  and  $g$  acts as a bijection between  $Y_2$  and  $X_2$ .

Now, consider the following function:

$$F : X \rightarrow Y : x \mapsto \begin{cases} f(x) & \text{if } x \in X_1 \\ g^{-1}(x) & \text{if } x \in X_2 \end{cases}$$

This function is one-to-one, as  $f$  and  $g^{-1}$  are. If  $y \in Y$ , then either  $y \in Y_1$  or  $Y_2$ . In the former case, we know there is  $x \in X_1$  such that  $f(x) = y$ , as  $f[X_1] = Y_1$ . In the latter case, we know there is a unique  $x$  such that  $g(y) = x$ . Thus,  $g^{-1}(x) = y$ . Thus,  $F$  is onto. We conclude  $F$  is a bijection. This proves CSBT.

22)  $X$  set of pairwise disj. non-empty sets.  $\mathcal{C}$  is a choice set  
set for  $X$  iff  $\forall x \in X, |x \cap \mathcal{C}| = 1$

(ACS): For all  $X$  as above, there exists a choice set

[To show: In ZF, (AC)  $\Leftrightarrow$  (ACS)]

Proof. " $\Rightarrow$ " Let  $X$  be a set of pairw. disj. non-empty sets. By assumption,  $\emptyset \notin X$ . By (AC), there exists a choice function  $f$  s.t.  $\forall x \in X, f(x) \in x$ . We then consider as our candidate choice set  $\mathcal{C} = f(X)$ .

For all  $x \in X$  we have  $f(x) \cap x \neq \emptyset$ , since  $f(x) \in x$ .

Suppose by contradiction that  $f(X) \cap x \neq \{f(x)\}$ , i.e.

there exists an  $z \in f(X) \cap x$ ,  $z \neq f(x)$ . Since  $z \in f(X)$

$z = f(y)$  for some  $y \in X$ . But then we know that

$f(y) \in y$  since  $f$  is a choice function. By (1) and (2)

we get  $f(y) \in x \cap y$ , which is a contradiction since

by assumption  $\forall x, y \in X, x \cap y = \emptyset$

" $\Leftarrow$ " Let  $S$  be a family of nonempty sets. Consider the set  $X = \{ \{x\} \times x, x \in S \}$ . Since  $x \neq \emptyset \forall x \in S, \emptyset \notin X$ .

Also, the elements of  $X$  are pairwise disjoint:

let  $\{x\} \times x, \{y\} \times y \in X$ , then  $\{x\} \times x \cap \{y\} \times y = \emptyset$

since  $x \neq y \forall x, y \in S$  (by extensionality). We then

apply (ACS) to  $X$ :  $\exists$  a choice set  $\mathcal{C}$  s.t.  $\forall \{x\} \times x \in X$

$|\mathcal{C} \cap \{x\} \times x| = 1$ , furthermore  $\mathcal{C} \cap \{x\} \times x = (x, x')$  where

$x' \in x$ . Then define the map  $f$  s.t.  $\forall x \in S$

$f(x) = x'$ , where  $x' = \mathcal{C} \cap \{x\} \times x$ . By construction,

$f$  is functional and it is a choice function for  $S$ .

2.3 ~~2.3~~ Prove Zorn's lemma in ZFC.

If  $(P, <)$  is a partially ordered set such that each chain in non-empty  $P$  has an upper bound, then  $P$  has a maximal elt.

Let  $(P, <)$  be a non-empty partially ordered set s.t. each chain  $C$  has an upper bound:

$$u \in C \text{ s.t. } c \leq u \quad \forall c \in C.$$

Now, define by transfinite recursion a sequence as follows:

Let  $c$  be the choice function.

$$\text{Def } a_0 := c(P) \in P \quad (\text{application AC})$$

$$\text{and } P_{a_0} := \{x \in P \mid a_0 < x\}$$

If  $P_{a_0} = \emptyset$ ,  $a_0$  is maximal and we are done.

If  $P_{a_0} \neq \emptyset$ , def  $a_1 := c(P_{a_0}) \in P_{a_0}$  and so:

$$a_0 := c(P) \in P \quad (\text{Application AC})$$

$$a_{n+1} := c(P_{a_n}) \in P_{a_n} \quad "$$

$$\text{and } a_\lambda := c(\bigcap_{\gamma < \lambda} P_\gamma) \in \bigcap_{\gamma < \lambda} P_\gamma \quad \text{for } \lambda \text{ limit ord. and } \gamma < \lambda.$$

If  $P$  has no maximal elt, this definition would define a function from the ordinals to  $P$  that is injective.

$$f: \begin{cases} 0 \mapsto c(P) \\ n+1 \mapsto c(P_{a_n}) \\ \lambda \mapsto c(\bigcap_{\gamma < \lambda} P_\gamma) \end{cases}$$

This is in contradiction with Hartogs's theorem, which says that for each set  $P$  there exists an ordinal that does not inject into  $P$ , so  $P$  must have a maximal elt.  $\square$