



Core Logic

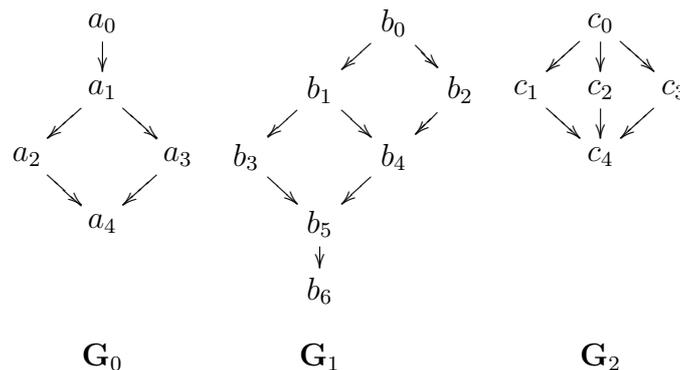
2007/2008; 1st Semester
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Homework Set # 11

Deadline: November 28th, 2007

Exercise 37 (9 points).

Consider the following three directed graphs G_0 , G_1 and G_2 . We say that a vertex x is **below** a vertex y if there is a path from y to x . We say that x is a **bottom element** in a graph if it is below all other elements in the graph. Our three graphs have bottom elements: a_4 , b_6 , and c_4 are the bottom elements of G_0 , G_1 , and G_2 , respectively.



In such a graph, if x and y are vertices, we define the **greatest lower bound of x and y** as follows. If $x = y$ or x is below y , we say that x is the greatest lower bound of x and y . Otherwise (i.e., if x and y are distinct vertices and neither is below the other), we call z the greatest lower bound of x and y if

- z is below x ,
- z is below y , and
- if any w is below both x and y , then either $w = z$ or w is below z .

Notice that a greatest lower bound has to be unique if it exists. Also notice that in the three given graphs, each pair of vertices has a greatest lower bound (you do not have to prove this, but please check for yourself by trying three or four examples).

We can now define a binary operation \wedge on the vertices by letting $x \wedge y$ be the greatest lower bound of x and y . Using this, we define a (possibly partial) unary operation $-$ on the vertices as follows: $-x = y$ if and only

- $x \wedge y$ is the bottom element, and
- if for any w , $x \wedge w$ is the bottom element, then either $w = y$ or w is below y .

(In other words: y is the greatest vertex such that $x \wedge y$ is the bottom element if this exists uniquely.)

- (1) For each of G_0 , G_1 , and G_2 , find out whether $-$ is a total function and give an argument (1½ points each).

- (2) With the given operation $-$, does G_0 satisfy the formula $--x = x$? (Give an argument; 2 points)
- (3) With the given operation $-$, does G_1 satisfy the formula $-- -x = -x$? (Give an argument; 2½ points)

Exercise 38 (7 points).

- (1) As mentioned in the lecture: Find wellorders W and W^* such that $W \oplus W^*$ is not isomorphic to $W^* \oplus W$ and explain why (2 points).
- (2) Similarly, find wellorders W and W^* such that $W \otimes W^*$ is not isomorphic to $W^* \otimes W$ and explain why (2 points).
- (3) In the first two tasks, you can choose one wellorder to be finite. Why can't both wellorders be finite in such an example (1 point)?
- (4) Consider $L := \langle \mathbb{Q}, \leq \rangle$ to be the rational numbers with the usual ordering. Find out whether $L \oplus L$ is isomorphic to L and give an argument (2 points).

Hint. The Cantor Isomorphism Theorem (sometimes called “back-and-forth theorem”) for countable linear orders may help. If you use it, you don't have to prove it, but please state it clearly with a proper reference to the literature and make sure that you apply it properly.

Exercise 39 (6 points).

We are modelling Achilles and the turtle as a transfinite process on the real line \mathbb{R} . Please give arguments for all answers.

- (1) Achilles' position at time t is given by A_t , the turtle's position is given by T_t . We start with $A_0 := 0$ and $T_0 := 1$. For every index i , we define $A_{i+1} := A_i + |T_i - A_i|$, $T_{i+1} := T_i + \frac{1}{2} \cdot |T_i - A_i|$, and

$$\begin{aligned} T_\infty &:= \lim_{i \in \mathbb{N}} T_i, \\ A_\infty &:= \lim_{i \in \mathbb{N}} A_i, \\ T_{\infty+\infty} &:= \lim_{i \in \mathbb{N}} T_{\infty+i}, \text{ and} \\ A_{\infty+\infty} &:= \lim_{i \in \mathbb{N}} A_{\infty+i}. \end{aligned}$$

Determine the least index i such that $A_i = T_i$ (1 point). Where is Achilles at time $\infty + \infty$ (1 point)?

- (2) Now the positions are given by A_t^* and T_t^* defined as follows. For each index $i \in \{0, 1, 2, \dots, \infty, \infty + 1, \infty + 2, \infty + 3, \dots\}$, we define the *value* $v(i)$ as follows:

$$v(i) := n \text{ if } i = n \text{ or } i = \infty + n.$$

We start with $A_0^* := 0$ and $T_0^* := 1$. For every index i , we define $A_{i+1}^* := A_i^* + \frac{1}{2^{v(i)}}$, $T_{i+1}^* := T_i^* + \frac{1}{2^{v(i)+1}}$, and

$$\begin{aligned} T_\infty^* &:= \lim_{i \in \mathbb{N}} T_i^*, \\ A_\infty^* &:= \lim_{i \in \mathbb{N}} A_i^*, \\ T_{\infty+\infty}^* &:= \lim_{i \in \mathbb{N}} T_{\infty+i}^*, \text{ and} \\ A_{\infty+\infty}^* &:= \lim_{i \in \mathbb{N}} A_{\infty+i}^*. \end{aligned}$$

Compute $A_{\infty+5}^*$, $T_{\infty+12}^*$, $A_{\infty+\infty}^*$ and $T_{\infty+\infty}^*$ (1 point each).