

ON THE INTERSECTION OF INFINITE MATROIDS

Elad Aigner-Horev* Johannes Carmesin Jan-Oliver Fröhlich

University of Hamburg
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Abstract

We show that the *infinite matroid intersection conjecture* of Nash-Williams implies the infinite Menger theorem proved recently by Aharoni and Berger.

We prove that this conjecture is true whenever one matroid is nearly finitary and the second is the dual of a nearly finitary matroid, where the nearly finitary matroids form a superclass of the finitary matroids.

In particular, this proves the infinite matroid intersection conjecture for finite-cycle matroids of 2-connected, locally finite graphs with only a finite number of vertex-disjoint rays.

1 Introduction

The infinite Menger theorem¹ was conjectured by Erdős in the 1960s and proved recently by Aharoni and Berger [1]. It states that for any two sets of vertices S and T in a connected graph, there is a set of vertex-disjoint S - T -paths whose maximality is witnessed by an S - T -separator picking exactly one vertex from each of these paths.

The complexity of the only known proof of this theorem and the fact that the finite Menger theorem has a short matroidal proof, make it natural to ask whether a matroidal proof of the infinite Menger theorem exists. In this paper, we propose a way to approach this issue by proving that a certain conjecture of Nash-Williams regarding infinite matroids implies the infinite Menger theorem.

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¹see Theorem 3.1 below.

Recently, Bruhn, Diestel, Kriesell, Pendavingh and Wollan [5] found axioms for infinite matroids in terms of independent sets, bases, circuits, closure and (relative) rank. These axioms allow for duality of infinite matroids as known from finite matroid theory, which settled an old problem of Rado. With these new axioms it is possible now to look which theorems of finite matroid theory have infinite analogues.

Here, we shall look at the *matroid intersection theorem*, which is a classical result in finite matroid theory [9]. It asserts that *the maximum size of a common independent set of two matroids M_1 and M_2 on a common ground set E is given by*

$$\min_{X \subseteq E} \text{rk}_{M_1}(X) + \text{rk}_{M_2}(E \setminus X), \quad (1)$$

where rk_{M_i} denotes the rank function of the matroid M_i .

In this paper, we consider the following conjecture of Nash-Williams, which first appeared in [2] and serves as an infinite analogue to the finite matroid intersection theorem².

Conjecture 1.1. [The infinite matroid intersection conjecture]
Any two matroids M_1 and M_2 on a common ground set E have a common independent set I admitting a partition $I = J_1 \cup J_2$ such that $\text{cl}_{M_1}(J_1) \cup \text{cl}_{M_2}(J_2) = E$.

Here, $\text{cl}_M(X)$ denotes the *closure* of a set X in a matroid M ; it consists of X and the elements spanned by X in M (see [9]). Originally, Nash-Williams's Conjecture just concerned *finitary matroids*, those all of whose circuits are finite.

1.1 Our results

Aharoni and Ziv [2] proved that Conjecture 1.1 implies the infinite analogues of König's and Hall's theorems. We strengthen this by showing that this conjecture implies the celebrated *infinite Menger theorem* (Theorem 3.1 below), which is known to imply the infinite analogues of König's and Hall's theorems [7].

Theorem 1.2. *The infinite matroid intersection conjecture for finitary matroids implies the infinite Menger theorem.*

²An alternative notion of infinite matroid intersection was recently proposed by Christian [6].

In finite matroid theory, an exceptionally short proof of the matroid intersection theorem employing the well-known *finite matroid union theorem* [9, 10] is known. The latter theorem asserts that for two finite matroids $M_1 = (E_1, \mathcal{I}_1)$ and $M_2 = (E_2, \mathcal{I}_2)$ the set system

$$\mathcal{I}(M_1 \vee M_2) = \{I_1 \cup I_2 \mid I_1 \in \mathcal{I}_1, I_2 \in \mathcal{I}_2\} \quad (2)$$

forms the set of independent sets of their *union matroid* $M_1 \vee M_2$.

In a previous paper [3, Proposition 1.1], we showed that for infinite matroids M_1 and M_2 , the set system $\mathcal{I}(M_1 \vee M_2)$ is not necessarily a matroid. This then raises the question of whether the traditional connection between (infinite) matroid union and intersection still holds. In this paper, we prove the following.

Theorem 1.3. *If M_1 and M_2 are matroids on a common ground set E and $M_1 \vee M_2^*$ is a matroid, then Conjecture 1.1 holds for M_1 and M_2 .*

Throughout, M^* denotes the dual of a matroid M .

In [3] we show that the ‘largest’ class of matroids for which one can have a union theorem is essentially a certain superclass of the finitary matroids called the *nearly finitary* matroids (to be defined next). This, together with Theorem 1.3, enables us to make additional progress on Conjecture 1.1, as set out below.

Nearly finitary matroids are defined as follows [3]. For any matroid M , taking as circuits only the finite circuits of M defines a (finitary) matroid with the same ground set as M . This matroid is called the *finitarization* of M and denoted by M^{fin} .

It is not hard to show that every basis B of M extends to a basis B^{fin} of M^{fin} , and conversely every basis B^{fin} of M^{fin} contains a basis B of M . Whether or not $B^{\text{fin}} \setminus B$ is finite will in general depend on the choices for B and B^{fin} , but given a choice for one of the two, it will no longer depend on the choice for the second one.

We call a matroid M *nearly finitary* if every base of its finitarization contains a base of M such that their difference is finite.

Next, let us look at some examples of nearly finitary matroids. There are three natural extensions to the notion of a finite graphic matroid in an infinite context [5]; each with ground set $E(G)$. The most studied one is the *finite-cycle matroid*, denoted $M_{FC}(G)$, whose circuits are the finite cycles of G . This is a finitary matroid, and hence is also nearly finitary.

The second extension is the *algebraic-cycle matroid*, denoted $M_A(G)$, whose circuits are the finite cycles and double rays of G [5, 4]³.

³ $M_A(G)$ is not necessarily a matroid for any G ; see [8].

Proposition 1.4. $M_A(G)$ is a nearly finitary matroid if and only if G has only a finite number of vertex-disjoint rays.

The third extension is the *topological-cycle matroid*, denoted $M_C(G)$ ⁴, whose circuits are the topological cycles of G (Thus $M_C^{\text{fin}}(G) = M_{FC}(G)$ for any finitely separable graph G ; see Section 5.2 or [4] for definitions).

Proposition 1.5. Suppose that G is 2-connected and locally finite. Then, $M_C(G)$ is a nearly finitary matroid if and only if G has only a finite number of vertex-disjoint rays.

Having introduced nearly finitary matroids, we now state the result of [3].

Theorem 1.6. [Nearly finitary union theorem [3]]
If M_1 and M_2 are nearly finitary matroids, then $M_1 \vee M_2$ is a nearly finitary matroid.

The following is a consequence of Theorem 1.6 and Theorem 1.3.

Corollary 1.7. Conjecture 1.1 holds for M_1 and M_2 whenever M_1 is nearly finitary and M_2 is the dual of a nearly finitary matroid.

Aharoni and Ziv [2] proved that the infinite matroid intersection conjecture is true whenever one matroid is finitary and the other is a countable direct sum of finite-rank matroids. Note that Corollary 1.7 does not imply this result of [2] nor is it implied by it.

Proposition 1.5 and Corollary 1.7 can be used to prove the following.

Corollary 1.8. Suppose that G and H are 2-connected, locally finite graphs with only a finite number of vertex-disjoint rays. Then their finite-cycle matroids $M_{FC}(G)$ and $M_{FC}(H)$ satisfy the intersection conjecture.

Similar results are true for the *algebraic-cycle matroid*, the *topological-cycle matroid*, and their duals.

This paper is organized as follows. Additional notation, terminology, and basic lemmas are given in Section 2. In Section 3 we prove Theorem 1.2. In Section 4 we prove Theorem 1.3, and in Section 5 we prove Propositions 1.4 and 1.5 and Corollary 1.8.

⁴ $M_C(G)$ is a matroid for any G ; see [4].

2 Preliminaries

Notation and terminology for graphs are that of [7], and for matroids that of [9, 5].

Throughout, G always denotes a graph where $V(G)$ and $E(G)$ denote its vertex and edge sets, respectively. We write M to denote a matroid and write $E(M)$, $\mathcal{I}(M)$, $\mathcal{B}(M)$, and $\mathcal{C}(M)$ to denote its ground set, independent sets, bases, and circuits, respectively.

We review the definition of a matroid as this is given in [5]. A set system \mathcal{I} is the set of independent sets of a matroid if it satisfies the following *independence axioms*:

- (I1) $\emptyset \in \mathcal{I}$.
- (I2) $[\mathcal{I}] = \mathcal{I}$, that is, \mathcal{I} is closed under taking subsets.
- (I3) Whenever $I, I' \in \mathcal{I}$ with I' maximal and I not maximal, there exists an $x \in I' \setminus I$ such that $I + x \in \mathcal{I}$.
- (IM) Whenever $I \subseteq X \subseteq E$ and $I \in \mathcal{I}$, the set $\{I' \in \mathcal{I} \mid I \subseteq I' \subseteq X\}$ has a maximal element.

The following is a well-known fact for finite matroids (see, e.g., [9]), which generalizes easily to infinite matroids.

Lemma 2.1. [5, Lemma 3.11]

Let M be a matroid. Then, $|C \cap C^| \neq 1$, whenever $C \in \mathcal{C}(M)$ and $C^* \in \mathcal{C}(M^*)$.*

We end this section with the definition of *exchange chains*. For a set $X \subseteq E(M)$, an X -*circuit* is a circuit containing X . For sets $I_1 \in \mathcal{I}(M_1)$ and $I_2 \in \mathcal{I}(M_2)$, and elements $x \in I_1 \cup I_2$ and $y \in E(M_1) \cup E(M_2)$ (possibly in $I_1 \cup I_2$), a tuple $Y = (y_0 = y, \dots, y_n = x)$ is called an *even (I_1, I_2, y, x) -exchange chain* (or *even (I_1, I_2, y, x) -chain*) of *length n* if the following terms are satisfied.

- (X1) For an even i , there exists a $\{y_i, y_{i+1}\}$ -circuit $C_i \subseteq I_1 + y_i$ of M_1 .
- (X2) For an odd i , there exists a $\{y_i, y_{i+1}\}$ -circuit $C_i \subseteq I_2 + y_i$ of M_2 .

If $n \geq 1$, then (X1) and (X2) imply that $y_0 \notin I_1$ and that, starting with $y_1 \in I_1 \setminus I_2$, the elements y_i alternate between $I_1 \setminus I_2$ and $I_2 \setminus I_1$; the single exception being y_n which might lie in $I_1 \cap I_2$.

By an *odd exchange chain* (or *odd chain*) we mean an even chain with the words ‘even’ and ‘odd’ interchanged in the definition. Consequently, we say *exchange chain* (or *chain*) to refer to either of these notions. Furthermore, a subchain of a chain is also a chain; that is, given an (I_1, I_2, y_0, y_n) -chain (y_0, \dots, y_n) , the tuple (y_k, \dots, y_l) is an (I_1, I_2, y_k, y_l) -chain for $0 \leq k \leq l \leq n$.

Lemma 2.2. [3, Lemma 4.4]

If there exists an (I_1, I_2, y, x) -chain, then $(I + y) - x \in \mathcal{I}(M_1 \vee M_2)$ where $I := I_1 \cup I_2$. Moreover, if $x \in I_1 \cap I_2$, then $I + y \in \mathcal{I}(M_1 \vee M_2)$.

3 From infinite matroid intersection to the infinite Menger theorem

In this section, we prove Theorem 1.2; asserting that the infinite matroid intersection conjecture implies the infinite Menger theorem.

Given a graph G and $S, T \subseteq V(G)$, a set $X \subseteq V(G)$ is called an S - T *separator* if $G - X$ contains no S - T path. The infinite Menger theorem reads as follows.

Theorem 3.1 (Aharoni and Berger [1]). *Let G be a connected graph. Then for any $S, T \subseteq V(G)$ there is a set \mathcal{L} of vertex-disjoint S - T paths and an S - T separator $X \subseteq \bigcup_{P \in \mathcal{L}} V(P)$ satisfying $|X \cap V(P)| = 1$ for each $P \in \mathcal{L}$.*

Infinite matroid union cannot be used in order to obtain the infinite Menger Theorem directly via Theorem 1.3 and Theorem 1.2. Indeed, in [3, Proposition 1.1] we construct a finitary matroid M and a co-finitary matroid N such that their union is not a matroid. Consequently, one cannot apply Theorem 1.3 to the finitary matroids M and N^* in order to obtain Conjecture 1.1 for them. However, it is easy to see that Conjecture 1.1 is true for these particular M and N^* .

Next, we prove Theorem 1.2.

Proof of Theorem 1.2. Let G be a connected graph and let $S, T \subseteq V(G)$ be as in Theorem 3.1. We may assume that $G[S]$ and $G[T]$ are both connected. Indeed, an S - T separator with $G[S]$ and $G[T]$ connected gives rise to an S - T separator when these are not necessarily connected. Abbreviate $E(S) := E(G[S])$ and $E(T) := E(G[T])$, let M be the finite-cycle matroid $M_F(G)$, and put $M_S := M/E(S) - E(T)$ and $M_T := M/E(T) - E(S)$; all three matroids are clearly finitary.

Assuming that the infinite matroid intersection conjecture holds for M_S and M_T , there exists a set $I \in \mathcal{I}(M_S) \cap \mathcal{I}(M_T)$ which admits a partition $I = J_S \cup J_T$ satisfying

$$\text{cl}_{M_S}(J_S) \cup \text{cl}_{M_T}(J_T) = E,$$

where $E = E(M_S) = E(M_T)$. We regard I as a subset of $E(G)$.

For the components of $G[I]$ we observe two useful properties. As I is disjoint from $E(S)$ and $E(T)$, the edges of a cycle in a component of $G[I]$ form a circuit in both, M_S and M_T , contradicting the independence of I in either. Consequently,

$$\text{the components of } G[I] \text{ are trees.} \quad (3)$$

Next, an S -path⁵ or a T -path in a component of $G[I]$ gives rise to a circuit of M_S or M_T in I , respectively. Hence,

$$|V(C) \cap S| \leq 1 \text{ and } |V(C) \cap T| \leq 1 \text{ for each component } C \text{ of } G[I]. \quad (4)$$

Let \mathcal{C} denote the components of $G[I]$ meeting both of S and T . Then by (3) and (4) each member of \mathcal{C} contains a unique S - T path and we denote the set of all these paths by \mathcal{L} . Clearly, the paths in \mathcal{L} are vertex-disjoint.

In what follows, we find a set X comprised of one vertex from each $P \in \mathcal{L}$ to serve as the required S - T separator. To that end, we show that one may alter the partition $I = J_S \cup J_T$ to yield a partition

$$I = K_S \cup K_T \text{ satisfying } \text{cl}_{M_S}(K_S) \cup \text{cl}_{M_T}(K_T) = E \text{ and (Y.1-4),} \quad (5)$$

where (Y.1-4) are as follows.

(Y.1) Each component C of $G[I]$ contains a vertex of $S \cup T$.

(Y.2) Each component C of $G[I]$ meeting S but not T satisfies $E(C) \subseteq K_S$.

(Y.3) Each component C of $G[I]$ meeting T but not S satisfies $E(C) \subseteq K_T$.

(Y.4) Each component C of $G[I]$ meeting both, S and T , contains at most one vertex which at the same time

- (a) lies in S or is incident with an edge of K_S , and
- (b) lies in T or is incident with an edge of K_T .

⁵A non-trivial path meeting $G[S]$ exactly in its end vertices.

Postponing the proof of (5), we first show how to deduce the existence of the required S - T separator from (5). Define a pair of sets of vertices (V_S, V_T) of $V(G)$ by letting V_S consist of those vertices contained in S or incident with an edge of K_S and defining V_T in a similar manner. Then $V_S \cap V_T$ may serve as the required S - T separator. To see this, we verify below that (V_S, V_T) satisfies all of the terms (Z.1-4) stated next.

(Z.1) $V_S \cup V_T = V(G)$;

(Z.2) for every edge e of G either $e \subseteq V_S$ or $e \subseteq V_T$;

(Z.3) every vertex in $V_S \cap V_T$ lies on a path from \mathcal{L} ; and

(Z.4) every member of \mathcal{L} meets $V_S \cap V_T$ at most once.

To see (Z.1), suppose v is a vertex not in $S \cup T$. As G is connected, such a vertex is incident with some edge $e \notin E(T) \cup E(S)$. The edge e is spanned by K_T or K_S ; say K_T . Thus, $K_T + e$ contains a circle containing e or $G[K_T + e]$ has a T -path containing e . In either case v is incident with an edge of K_T and thus in V_T , as desired.

To see (Z.2), let $e \in cl_{M_T}(K_T) \setminus K_T$; so that $K_T + e$ has a circle containing e or $G[K_T + e]$ has T -path containing e ; in either case both end vertices of e are in V_T , as desired. The treatment of the case $e \in cl_{M_S}(K_S)$ is similar.

To see (Z.3), let $v \in V_S \cap V_T$; such is in S or is incident with an edge of K_S , and in T or is incident with an edge in K_T . Let C be the component of $G[I]$ containing v . By (Y.1-4), $C \in \mathcal{C}$, i.e. it meets both, S and T and therefore contains an S - T path $P \in \mathcal{L}$. Recall that every edge of C is either in K_S or K_T and consider the last vertex w of a maximal initial segment of P in $C - K_T$. Then w satisfies (Y.4a), as well as (Y.4b), implying $v = w$; so that v lies on a path from \mathcal{L} .

To see (Z.4), we restate (Y.4) in terms of V_S and V_T : each component of \mathcal{C} contains at most one vertex of $V_S \cap V_T$. This clearly also holds for the path from \mathcal{L} which is contained in C .

It remains to prove (5). To this end, we show that any component C of $G[I]$ contains a vertex of $S \cup T$. Suppose not. Let e be the first edge on a $V(C)$ - S path Q which exists by the connectedness of G . Then $e \notin I$ but without loss of generality we may assume that $e \in cl_{M_S}(J_S)$. So in $G[I] + e$ there must be a cycle or an S -path. The latter implies that C contains a vertex of S and the former means that Q was not internally disjoint to $V(C)$, yielding contradictions in both cases.

We define the sets K_S and K_T as follows. Let C be a component of $G[I]$.

1. If C meets S but not T , then include its edges into K_S .
2. If C meets T but not S , then include its edges into K_T .
3. Otherwise (C meets both of S and T) there is a path P from \mathcal{L} in C . Denote by v_C the last vertex of a maximal initial segment of P in $C - J_T$. As C is a tree, each component C' of $C - v_C$ is a tree and there is a unique edge e between v_C and C' . For every such component C' , include the edges of $C' + e$ in K_S if $e \in J_S$ and in K_T otherwise, i.e. if $e \in J_T$.

Note that, by choice of v_C , either v_C is the last vertex of P or the next edge of P belongs to J_T . This ensures that K_S and K_T satisfy (Y.4). Moreover, they form a partition of I which satisfies (Y.1-3) by construction. It remains to show that $\text{cl}_{M_S}(K_S) \cup \text{cl}_{M_T}(K_T) = E$.

As $K_S \cup K_T = I$, it suffices to show that any $e \in E \setminus I$ is spanned by K_S in M_S or by K_T in M_T . Suppose $e \in \text{cl}_{M_S}(J_S)$, i.e. $J_S + e$ contains a circuit of M_S . Hence, $G[J_S]$ either contains an e -path R or two disjoint e - S paths R_1 and R_2 . We show that $E(R) \subseteq K_S$ or $E(R) \subseteq K_T$ in the former case and $E(R_1) \cup E(R_2) \subseteq K_S$ in the latter.

The path R is contained in some component C of $G[I]$. Suppose $C \in \mathcal{C}$ and v_C is an inner vertex of R . By assumption, the edges preceding and succeeding v_C on R are both in J_S and hence the edges of both components of $C - v_C$ which are met by R plus their edges to v_C got included into K_S , showing $E(R) \subseteq K_S$. Otherwise $C \notin \mathcal{C}$ or $C \in \mathcal{C}$ but v_C is no inner vertex of R . In both cases the whole set $E(R)$ got included into K_S or K_T .

We apply the same argument to R_1 and R_2 except for one difference. If $C \notin \mathcal{C}$ or $C \in \mathcal{C}$ but v_C is no inner vertex of R_i , then $E(R_i)$ got included into K_S as R_i meets S .

Although the definitions of K_S and K_T are not symmetrical, a similar argument shows $e \in \text{cl}_{M_S}(K_S) \cup \text{cl}_{M_T}(K_T)$ if e is spanned by J_T in M_T . \square

Note that the above proof requires only that Conjecture 1.1 holds for finite-cycle matroids.

4 From infinite matroid union to infinite matroid intersection

In this section, we prove Theorem 1.3.

Proof of Theorem 1.3. Our starting point is the well-known proof from finite matroid theory that matroid union implies a solution to the matroid intersection problem. With that said, let $B_1 \cup B_2^* \in \mathcal{B}(M_1 \vee M_2^*)$ where $B_1 \in \mathcal{B}(M_1)$ and $B_2^* \in \mathcal{B}(M_2^*)$, and let $B_2 = E \setminus B_2^* \in \mathcal{B}(M_2)$. Then, put $I = B_1 \cap B_2$ and note that $I \in \mathcal{I}(M_1) \cap \mathcal{I}(M_2)$. We show that I admits the required partition.

For an element $x \notin B_i$, $i = 1, 2$, we write $C_i(x)$ to denote the fundamental circuit of x into B_i in M_i . For an element $x \notin B_2^*$, we write $C_2^*(x)$ to denote the fundamental circuit of x into B_2^* in M_2^* . Put $X = B_1 \cap B_2^*$, $Y = B_2 \setminus I$, and $Z = B_2^* \setminus X$, see Figure 1.

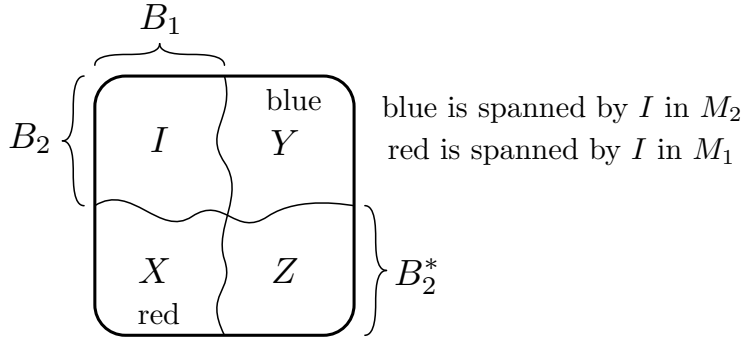


Figure 1: The sets X , Y , and Z and their colorings.

Observe that

$$cl_{M_1}(I) \cup cl_{M_2}(I) = E = I \cup X \cup Y \cup Z. \quad (6)$$

To see (6), note first that

$$X \subseteq cl_{M_2}(I). \quad (7)$$

Clearly, no member of X is spanned by I in M_1 . Assume then that $x \in X$ is not spanned by I in M_2 so that there exists a $y \in C_2(x) \cap Y$. Then, $x \in C_2^*(y)$, by Lemma 2.1. Consequently, $B_1 \cup B_2^* \subsetneq B_1 \cup (B_2^* + y - x) \in \mathcal{I}(M_1 \vee M_2^*)$; contradiction to the maximality of $B_1 \cup B_2^*$, implying (7).

By a similar argument, it holds that

$$Y \subseteq cl_{M_1}(I). \quad (8)$$

To see that

$$Z \subseteq cl_{M_1}(I) \cup cl_{M_2}(I), \quad (9)$$

assume, towards contradiction, that some $z \in Z$ is not spanned by I neither in M_1 nor in M_2 so that there exist an $x \in C_1(z) \cap X$ and a $y \in C_2(z) \cap Y$. Then $B_1 - x + z$ and $B_2 - y + z$ are bases and thus $B_1 \cup B_2^* \subsetneq (B_1 - x + z) \cup (B_2^* - z + y)$; contradiction to the maximality of $B_1 \cup B_2^*$. Assertion (6) is proved.

The problem of finding a suitable partition $I = J_1 \cup J_2$ can be phrased as a (directed) graph coloring problem. By (6), each $x \in E \setminus I$ satisfies $C_1(x) - x \subseteq I$ or $C_2(x) - x \subseteq I$. Define $G = (V, E)$ to be the directed graph whose vertex set is $V = E \setminus I$ and whose edge set is given by

$$E = \{(x, y) : C_1(x) \cap C_2(y) \cap I \neq \emptyset\}. \quad (10)$$

Recall that a *source* is a vertex with no incoming edges and a *sink* is a vertex with no outgoing edges. As $C_1(x)$ does not exist for any $x \in X$ and $C_2(y)$ does not exist for any $y \in Y$, it follows that

$$\text{the members of } X \text{ are sinks and those of } Y \text{ are sources in } G. \quad (11)$$

A 2-coloring of the vertices of G , by say blue and red, is called *divisive* if it satisfies the following:

- (D.1) I spans all the blue elements in M_1 ;
- (D.2) I spans all the red elements in M_2 ; and
- (D.3) $J_1 \cap J_2 = \emptyset$ where $J_1 := (\bigcup_{x \text{ blue}} C_1(x)) \cap I$ and $J_2 := (\bigcup_{x \text{ red}} C_2(x)) \cap I$.

Clearly, if G has a divisive coloring, then I admits the required partition.

We show then that G admits a divisive coloring. Color with blue all the sources. These are the vertices that can only be spanned by I in M_1 . Color with red all the sinks, that is, all the vertices that can only be spanned by I in M_2 . This defines a partial coloring of G in which all members of X are red and those of Y are blue. Such a partial coloring can clearly be extended into a divisive coloring of G provided that

$$G \text{ has no } (y, x)\text{-path with } y \text{ blue and } x \text{ red.} \quad (12)$$

Indeed, given (12) and (11), color all vertices reachable by a path from a blue vertex with the color blue, color all vertices from which a red vertex is reachable by a path with red, and color all remaining vertices with, say, blue. The resulting coloring is divisive.

It remains to prove (12). We show that the existence of a path as in (12) contradicts the following property:

Suppose that M and N are matroids and $B \cup B'$ is maximal in $\mathcal{I}(M \vee N)$. Let $y \notin B \cup B'$ and let $x \in B \cap B'$. Then, (by Lemma 2.2)

$$\text{there exists no } (B, B', y, x)\text{-chain}; \quad (13)$$

(in fact, the contradiction in the proofs of (7),(8), and (9) arose from simple instances of such forbidden chains).

Assume, towards contradiction, that P is a (y, x) -path with y blue and x red; the intermediate vertices of such a path are not colored since they are not a sink nor a source. In what follows we use P to construct a $(B_1, B_2^*, y_0, y_{2|P|})$ -chain $(y_0, y_1, \dots, y_{2|P|})$ such that $y_0 \in Y$, $y_{2|P|} \in X$, all odd indexed members of the chain are in $V(P) \cap Z$, and all even indexed elements of the chain other than y_0 and $y_{2|P|}$ are in I . Existence of such a chain would contradict (13).

Definition of y_0 . As y is pre-colored blue then either $y \in Y$ or $C_2(y) \cap Y \neq \emptyset$. In the former case set $y_0 = y$ and in the latter choose $y_0 \in C_2(y) \cap Y$.

Definition of $y_{2|P|}$. In a similar manner, x is pre-colored red since either $x \in X$ or $C_1(x) \cap X \neq \emptyset$. In the former case, set $y_{2|P|} = x$ and in the latter case choose $y_{2|P|} \in C_1(x) \cap X$.

The remainder of the chain. Enumerate $V(P) \cap Z = \{y_1, y_3, \dots, y_{2|P|-1}\}$ where the enumeration is with respect to the order of the vertices defined by P . Next, for an edge $(y_{2i-1}, y_{2i+1}) \in E(P)$, let $y_{2i} \in C_1(y_{2i-1}) \cap C_2(y_{2i+1}) \cap I$; such exists by the assumption that $(y_{2i-1}, y_{2i+1}) \in E$. As $y_{2i+1} \in C_2^*(y_{2i})$ for all relevant i , by Lemma 2.1, the sequence $(y_0, y_1, y_2, \dots, y_{2|P|})$ is a $(B_1, B_2^*, y_0, y_{2|P|})$ -chain in $\mathcal{I}(M_1 \vee M_2^*)$.

This completes our proof of Theorem 1.3. \square

Note that in the above proof, we do not use the assumption that $M_1 \vee M_2^*$ is a matroid; in fact, we only need that $\mathcal{I}(M_1 \vee M_2^*)$ has a maximal element.

5 The graphic nearly finitary matroids

In this section we prove Propositions 1.4 and 1.5 yielding a characterization of the graphic nearly finitary matroids.

For a connected graph G , a maximal set of edges containing no finite cycles is called an *ordinary spanning tree*. A maximal set of edges containing no finite cycles nor any double ray is called an *algebraic spanning tree*.

These are the bases of $M_F(G)$ and $M_A(G)$, respectively. We postpone the discussion about $M_C(G)$ to Section 5.2.

To prove Propositions 1.4 and 1.5, we require the following theorem of Halin [7, Theorem 8.2.5].

Theorem 5.1 (Halin 1965). *If an infinite graph G contains k disjoint rays for every $k \in \mathbb{N}$, then G contains infinitely many disjoint rays.*

5.1 The nearly finitary algebraic-cycle matroids

The purpose of this subsection is to prove Proposition 1.4.

Proof of Proposition 1.4. Suppose that G has k disjoint rays for every integer k ; so that G has a set \mathcal{R} of infinitely many disjoint rays by Theorem 5.1. We show that $M_A(G)$ is not nearly finitary.

The edge set of $\bigcup \mathcal{R} = \bigcup_{R \in \mathcal{R}} R$ is independent in $M_A(G)^{\text{fin}}$ as it induces no finite cycle of G . Therefore there is a base of $M_A(G)^{\text{fin}}$ containing it; such induces an ordinary spanning tree, say T , of G . We show that

$$T - F \text{ contains a double ray for any finite edge set } F \subseteq E(T). \quad (14)$$

This implies that $E(T) \setminus I$ is infinite for every independent set I of $M_A(G)$ and hence $M_A(G)$ is not nearly finitary. To see (14), note that $T - F$ has $|F| + 1$ components for any finite edge set $F \subseteq E(T)$ as T is a tree and successively removing edges always splits one component into two. So one of these components contains infinitely many disjoint rays from \mathcal{R} and consequently a double ray.

Suppose next, that G has at most k disjoint rays for some integer k and let T be an ordinary spanning tree of G , that is, $E(T)$ is maximal in $M_A(G)^{\text{fin}}$. To prove that $M_A(G)$ is nearly finitary, we need to find a finite set $F \subseteq E(T)$ such that $E(T) \setminus F$ is independent in $M_A(G)$, i.e. it induces no double ray of G . Let \mathcal{R} be a maximal set of disjoint rays in T ; such exists by assumption and $|\mathcal{R}| \leq k$. As T is a tree and the rays of \mathcal{R} are vertex-disjoint, it is easy to see that T contains a set F of $|\mathcal{R}| - 1$ edges such that $T - F$ has $|\mathcal{R}|$ components which each contain one ray of \mathcal{R} . By maximality of \mathcal{R} no component of $T - F$ contains two disjoint rays, or equivalently, a double ray. \square

5.2 The nearly finitary topological-cycle matroids

In this section we prove Proposition 1.5 that characterizes the nearly finitary topological-cycle matroids. Prior to that, we first define these matroids. To

that end we shall require some additional notation and terminology on which more details can be found in [4].

An *end of G* is an equivalence class of rays, where two rays are *equivalent* if they cannot be separated by a finite edge set. In particular, two rays meeting infinitely often are equivalent. Let the *degree* of an end ω be the size of a maximal set of vertex-disjoint rays belonging to ω , which is well-defined [7]. We say that a double ray *belongs to* an end if the two rays which arise from the removal of one edge from the double ray belong to that end; this does not depend on the choice of the edge. Such a double ray is an example of a *topological cycle*⁶

For a graph G the topological-cycle matroid of G , namely $M_C(G)$, has $E(G)$ as its ground set and its set of circuits consists of the finite and topological cycles. In fact, every infinite circuit of $M_C(G)$ induces at least one double ray; provided that G is locally finite [7].

A graph G has only finitely many disjoint rays if and only if G has only finitely many ends, each with finite degree. Also, note that

$$\text{every end of a 2-connected locally finite graph has degree at least 2.} \quad (15)$$

Indeed, applying Menger's theorem inductively, it is easy to construct in any k -connected graph for any end ω a set of k disjoint rays of ω .

Now we are in a position to start the proof of Proposition 1.5.

Proof of Proposition 1.5. If G has only a finite number of vertex-disjoint rays then $M_A(G)$ is nearly finitary by Proposition 1.4. Since $M_A(G)^{\text{fin}} = M_C(G)^{\text{fin}}$ and $\mathcal{I}(M_A(G)) \subseteq \mathcal{I}(M_C(G))$, we can conclude that $M_C(G)$ is nearly finitary as well.

Now, suppose that G contains k vertex-disjoint rays for every $k \in \mathbb{N}$. If G has an end ω of infinite degree, then there is an infinite set \mathcal{R} of vertex-disjoint rays belonging to ω . As any double ray containing two rays of \mathcal{R} forms a circuit of $M_C(G)$, the argument from the proof of Proposition 1.4 shows that $M_C(G)$ is not nearly finitary.

Assume then that G has no end of infinite degree. There are infinitely many disjoint rays, by Theorem 5.1. Hence, there is a countable set of ends $\Omega = \{\omega_1, \omega_2, \dots\}$.

We inductively construct a set \mathcal{R} of infinitely many vertex-disjoint double rays, one belonging to each end of Ω . Suppose that for any integer $n \geq 0$ we

⁶Formally, the topological cycles of G are those subgraphs of G which are homeomorphic images of S^1 in the Freudenthal compactification $|G|$ of G . However, the given example is the only type of topological cycle which shall be needed for the proof.

have constructed a set \mathcal{R}_n of n disjoint double rays, one belonging to each of the ends $\omega_1, \dots, \omega_n$. Different ends can be separated by finitely many vertices so there is a finite set S of vertices such that $\bigcup \mathcal{R}_n$ has no vertex in the component C of $G - S$ which contains ω_{n+1} . Since ω_{n+1} has degree 2 by (15), there are two disjoint rays from ω_{n+1} in C and thus also a double ray D belonging to ω_{n+1} . Set $\mathcal{R}_{n+1} := \mathcal{R}_n \cup \{D\}$ and $\mathcal{R} := \bigcup_{n \in \mathbb{N}} \mathcal{R}_n$.

As $\bigcup \mathcal{R}$ contains no finite cycle of G , it can be extended to an ordinary spanning tree of G . Removing finitely many edges from this tree clearly leaves an element of \mathcal{R} intact. Hence, the edge set of the resulting graph still contains a circuit of $M_C(G)$. Thus, $M_C(G)$ is not nearly finitary in this case as well. \square

In the following we shall propose a possible common generalization of Propositions 1.4 and 1.5 to all infinite matroids. We call a matroid M *k-nearly finitary* if every base of its finitarization contains a base of M such that their difference has size at most k . Note that saying ‘at most k ’ is not equivalent to saying ‘equal to k ’, consider for example the algebraic-cycle matroid of the infinite ladder. In terms of this new definition, Propositions 1.4 and 1.5 both state for a certain class of infinite matroids that each member of this class is *k-nearly finitary* for some k . In fact, for all known examples of nearly finitary matroids, there is such a k . This raises the following open question.

Question 5.2. *Is every nearly finitary matroid k-nearly finitary for some k?*

5.3 Graphic matroids and the intersection conjecture

By Corollary 1.7, the intersection conjecture is true for $M_C(G)$ and $M_{FC}(H)$ for any two graphs G and H since the first is co-finitary and the second is finitary. Using also Proposition 1.5, we obtain the following.

Corollary 5.3. *Suppose that G and H are 2-connected, locally finite graphs with only a finite number of vertex-disjoint rays. Then, $M_C(G)$ and $M_C(H)$ satisfy the intersection conjecture.* \square

Using Proposition 1.4 instead of Proposition 1.5, we obtain the following.

Corollary 5.4. *Suppose that G and H are graphs with only a finite number of vertex-disjoint rays. Then, $M_A(G)$ and $M_A(H)$ satisfy the intersection conjecture if both are matroids.* \square

With a little more work, the same is also true for $M_{FC}(G)$, see Corollary 1.8.

Proof of Corollary 1.8. First we show that $((M_C(G)^{\text{fin}})^*)^{\text{fin}} = M_C(G)$ if G is locally finite. Indeed, then $M_C(G)^{\text{fin}} = M_{FC}(G)$, $M_{FC}(G)^*$ is the matroids whose circuits are the finite and infinite bonds of G , and its finitization has as its circuits the finite bonds of G . And the dual of this matroid is $M_C(G)$, see [5] for example.

Having showed that $((M_C(G)^{\text{fin}})^*)^{\text{fin}} = M_C(G)$ if G is locally finite, we next show that if $M_C(G)$ is nearly finitary, then so is $M_{FC}(G)^*$. For this let B be a base of $M_{FC}(G)^*$ and B' be a base of $(M_{FC}(G)^*)^{\text{fin}}$. Then $B' \setminus B = (E \setminus B) \setminus (E \setminus B')$. Now $E \setminus B$ is a base of $M_{FC}(G) = M_C(G)^{\text{fin}}$ and by the above $E \setminus B'$ is a base of $M_C(G)$. Since $M_C(G)$ is nearly finitary, $B' \setminus B$ is finite, yielding that $M_{FC}(G)^*$ is nearly finitary.

As $M_{FC}(G)^*$ is nearly finitary and $M_{FC}(H)$ is finitary, $M_{FC}(H)$ and $M_{FC}(G)$ satisfy the intersection conjecture by Corollary 1.7. \square

A similar argument shows that if G and H are 2-connected, locally finite graphs with only a finite number of vertex-disjoint rays, then one can also prove that $M_{FC}(G)^*$ and $M_{FC}(H)^*$ satisfy the intersection conjecture. Similar results are true for $M_C(G)^*$ or $M_A(G)^*$ in place of $M_{FC}(G)^*$.

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