

**Hints and solutions for problem sheet # 05**  
**Advanced Algebra — Winter term 2016/17**  
(Ingo Runkel)

---

**Problem 17**

1.  $R\text{-Mod}$  has kernels and co-kernels, so purely as maps of  $k$ -vector spaces, we know that every  $g : A \rightarrow B$  will have a kernel and cokernel. Now to produce filtrations of these.

**For the kernel:** As a  $k$ -linear map,  $g$  has a kernel the sub-vector space of  $A$  given by  $a \in A$  which are sent to 0 under  $g$ ;  $(\ker(g), K : \ker(g) \hookrightarrow A)$ . As a result, the induced associated filtration of the kernel is defined as  $\ker(g)_i := \ker(g) \cap A_i$ .

**Proof that this is the kernel:** Let  $\alpha : U \rightarrow A$  be a  $\mathbb{Z}$ -filtered linear map such that  $g \circ \alpha = 0$ . We need to show that there is a unique  $\mathbb{Z}$ -filtered linear map  $\tilde{\alpha} : U \rightarrow \ker(g)$  such that  $K \circ \tilde{\alpha} = \alpha$ . But this is clear from the definitions as in this case  $\tilde{\alpha} = \alpha$ , understood as a filtered linear map  $U \rightarrow \ker(g)$ .

**For the cokernel:** Recall that as an  $k$ -vector space, the cokernel is of the form  $(B/(\text{im}(g)), c : B \rightarrow B/(\text{im}(g)))$ . Then the filtration on  $B_i$  induces one on the cokernel of  $g$ , via  $c$ ;  $(\text{coker}(g))_i := c(B_i)$ .

**Proof that this is the cokernel:** (is similar to kernel).

2. As on the underlying  $k$ -vector space  $k$ ,  $f$  is the identity map, and mono and epi are immediate.

All filtered linear maps  $W \rightarrow V$  are zero, since to preserve filtration, they have to map  $W_0 = k$  into  $V_0 = 0$ . The category of  $\mathbb{Z}$ -filtered vector spaces does not contain an inverse to  $f$ .

**Problem 18**

1. Write  $R := \mathcal{C}(A, A)$ . Since  $\mathcal{C}$  is an Ab-category,  $R$  is an abelian group. The composition of  $\mathcal{C}$  defines an associative composition  $\circ : R \times R \rightarrow R$ . By additivity of  $\mathcal{C}$ , the composition is bilinear, i.e. the distributive law holds. The identity  $1_A \in \mathcal{C}(A, A)$  is the unit of  $R$ . (A one-object category can be additive only if its unique object is the zero object, so that  $R = \{0\}$ .)
2. Let  $F : \mathcal{C} \rightarrow \mathbf{Ab}$  be an additive functor. Thus  $F(\bullet) =: M$  is an abelian group, and  $F : \mathcal{C}(\bullet, \bullet) \rightarrow \text{End}(M)$  is a homomorphism of abelian groups. Since  $F$  preserves unit and composition, we have  $F(1_R) = \text{id}_M$  and  $F(f \circ g) = F(f) \circ F(g)$ . Thus  $F : R \rightarrow \text{End}(M)$  is a ring homomorphism. By Problem 3, this amounts to turning  $M$  into an  $R$ -module.

Conversely, given an  $R$ -module  $M$  we can set  $F(\bullet) = M$  and use the group homomorphism  $R \rightarrow \text{End}(M)$  from Problem 3 to define  $F$  on morphisms.

3. Let  $F, G : \mathcal{C} \rightarrow \mathbf{Ab}$  be the two additive functors. Write  $M = F(\bullet)$  and  $N = G(\bullet)$  for the corresponding  $R$ -modules. The natural transformation consists of a single map  $\eta_\bullet =: f : M \rightarrow N$  (a group homomorphism, as it is a morphism in  $\mathbf{Ab}$ ). The naturality square reads, for  $r \in R = \mathcal{C}(\bullet, \bullet)$ ,

$$\begin{array}{ccc} M & \xrightarrow{F(r)} & M \\ f \downarrow & & \downarrow f \\ N & \xrightarrow{G(r)} & N \end{array}$$

This commutes for all  $r$  if and only if  $f$  is an  $R$ -module homomorphism.

### Problem 19

1. The morphism sets are

$$(\mathcal{C} \times \mathcal{D})((C_1, D_1), (C_2, D_2)) := \{(f, g) \in \mathcal{C}(C_1, C_2) \times \mathcal{D}(D_1, D_2)\}$$

Composition is coordinate-wise. Associativity and identity properties follow from those of  $\mathcal{C}$  and  $\mathcal{D}$ .

2. Let  $F$  be a functor from the product-category  $\mathcal{C} \times \mathcal{D}$  to  $\mathcal{E}$ . Then  $F_C : \mathcal{C} \rightarrow \mathcal{E}$  is given by, for  $D, D' \in \mathcal{D}$  and  $f : D \rightarrow D'$ :

$$F_C(D) = F((C, D)) \quad , \quad F_C(f) = F((1_C, f)) .$$

The functor  $F_D : \mathcal{D} \rightarrow \mathcal{E}$  is defined analogously. From functoriality of  $F$  one immediately obtains that  $F_C, F_D$  are functors. For the commuting square note

$$\begin{aligned} F_{C'}(g) \circ F_D(f) &= F((1_{C'}, g)) \circ F((f, 1_D)) = F((1_{C'}, g) \circ (f, 1_D)) \\ &= F((f, g)) = F((f, 1_{D'}) \circ (1_C, g)) = F_{D'}(f) \circ F_C(g) . \end{aligned}$$

3. We define  $F$  on objects  $(C, D) \in \mathcal{C} \times \mathcal{D}$  as

$$F((C, D)) := F_D(C) = F_C(D) .$$

On morphisms  $(f, g) : (C, D) \rightarrow (C', D')$  we define

$$F((f, g)) := F_{D'}(f) \circ F_C(g) = F_{C'}(g) \circ F_D(f) ,$$

where the equality is the commuting diagram we assume.

Now we need to show that  $F$  is a functor. We have  $F((1_C, 1_D)) = F_D(1_C) \circ F_C(1_D) = 1_{(C, D)} \circ 1_{(C, D)} = 1_{(C, D)}$ . For morphisms  $(f, g) : (C, D) \rightarrow (C', D')$

and  $(h, k) : (C', D') \rightarrow (C'', D'')$  we have

$$\begin{aligned}
F((h, k) \circ (f, g)) &= F((h \circ f, k \circ g)) \\
&= F_{D''}(h \circ f) \circ F_C(k \circ g) \\
&= F_{D''}(h) \circ F_{D''}(f) \circ F_C(k) \circ F_C(g) \\
&\stackrel{(*)}{=} F_{D''}(h) \circ F_{C'}(k) \circ F_{D'}(f) \circ F_C(g) \\
&= F((h, k)) \circ F((f, g)) ,
\end{aligned}$$

where in (\*) the commuting square was used.

### Problem 20

1. Simpler than 2, so let's just look at 2:
2. For  $f : A \rightarrow A'$ ,  $x \in \mathcal{C}(A', B)$  the definition of  $f^* : \mathcal{C}(A', B) \rightarrow \mathcal{C}(A, B)$  was  $f^*(x) = x \circ f$ .  $\mathcal{C}(-, B)$  being a contravariant functor amounts to checking  $id^*(x) = x$  and, for  $g : A'' \rightarrow A$ ,  $(f \circ g)^*(x) = g^*(f^*(x))$ . The first identity is immediate, and for the second note that both sides are equal to  $x \circ f \circ g$ .
3. Given  $f : A \rightarrow A'$  and  $h : B \rightarrow B'$ , and any  $g \in \mathcal{C}(A, B)$ , then  $\mathcal{C}(id_A, h) : \mathcal{C}(A, B) \rightarrow \mathcal{C}(A, B')$  via  $g \mapsto h \circ g$ .  $\mathcal{C}(f, id_B) : \mathcal{C}(A', B) \rightarrow \mathcal{C}(A, B)$  via  $g \mapsto g \circ f$ . And so on, so that we map  $g \in \mathcal{C}(A, B)$  to  $\mathcal{C}(A', B')$  via  $g \mapsto h \circ g \mapsto (h \circ g) \circ f$  in one direction and  $g \mapsto g \circ f \mapsto h \circ (g \circ f)$  in the other. Associativity of composition of morphisms means these two are equivalent, and we have the square as desired.